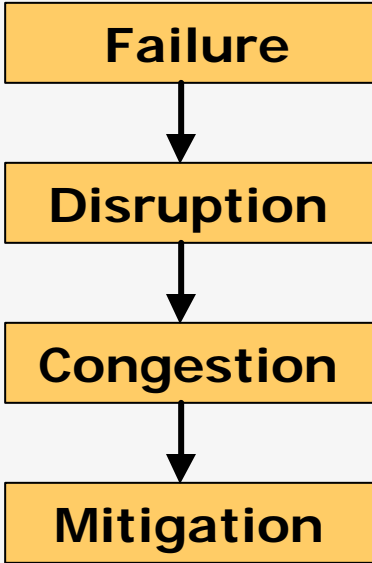
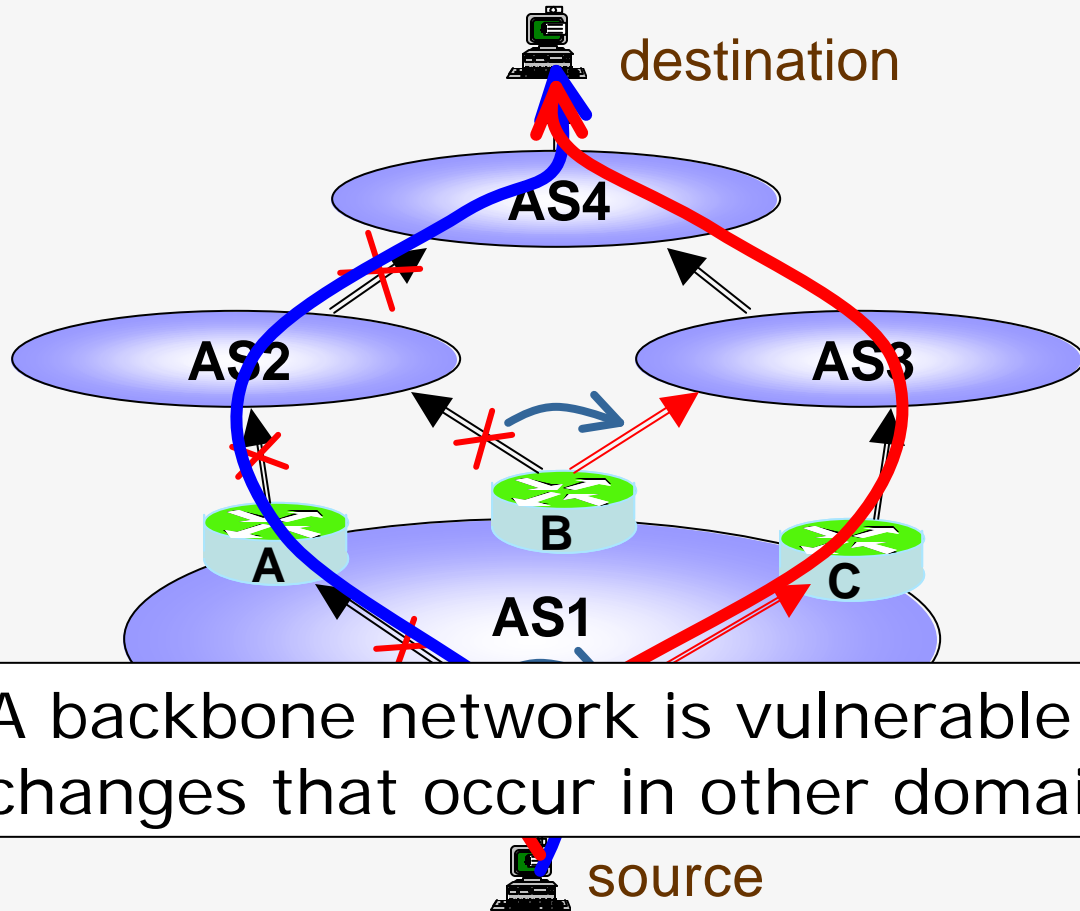


Finding a Needle in a Haystack: Pinpointing Significant BGP Routing Changes in an IP Network

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Z. Morley Mao (University of Michigan)
Jennifer Rexford (Princeton University)
Jia Wang (AT&T Labs Research)

Motivation



A backbone network is vulnerable to routing changes that occur in other domains.

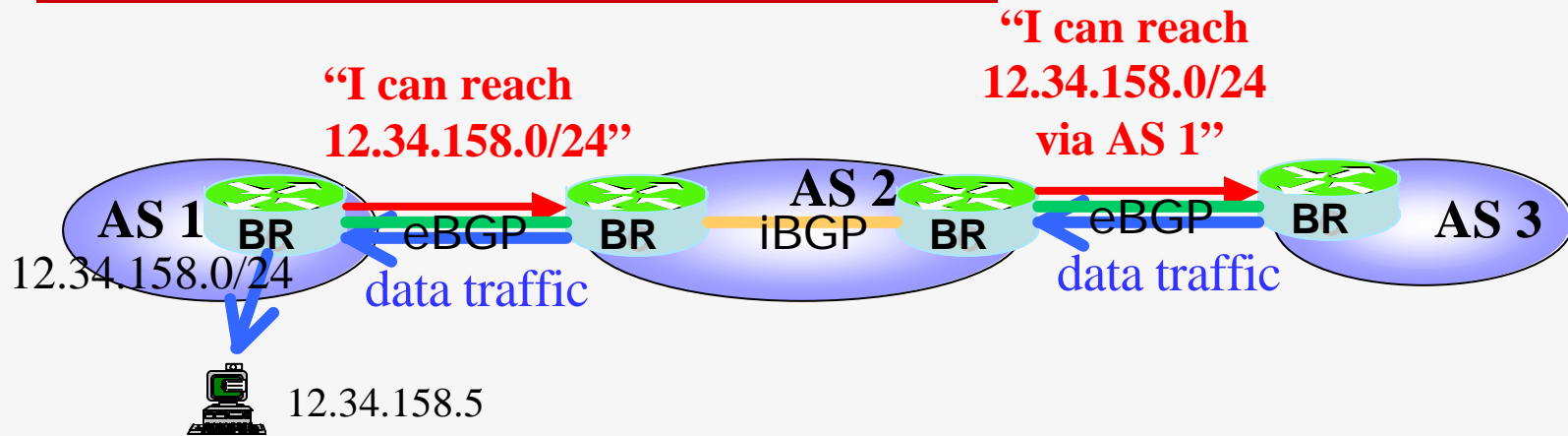
Goal

- Identify important routing anomalies
 - Lost reachability
 - Persistent flapping
 - Large traffic shifts

Contributions:

- Build a tool to identify a **small** number of **important** routing disruptions from a **large volume** of raw BGP updates in **real time**.
- Use the tool to characterize routing disruptions in **an operational network**

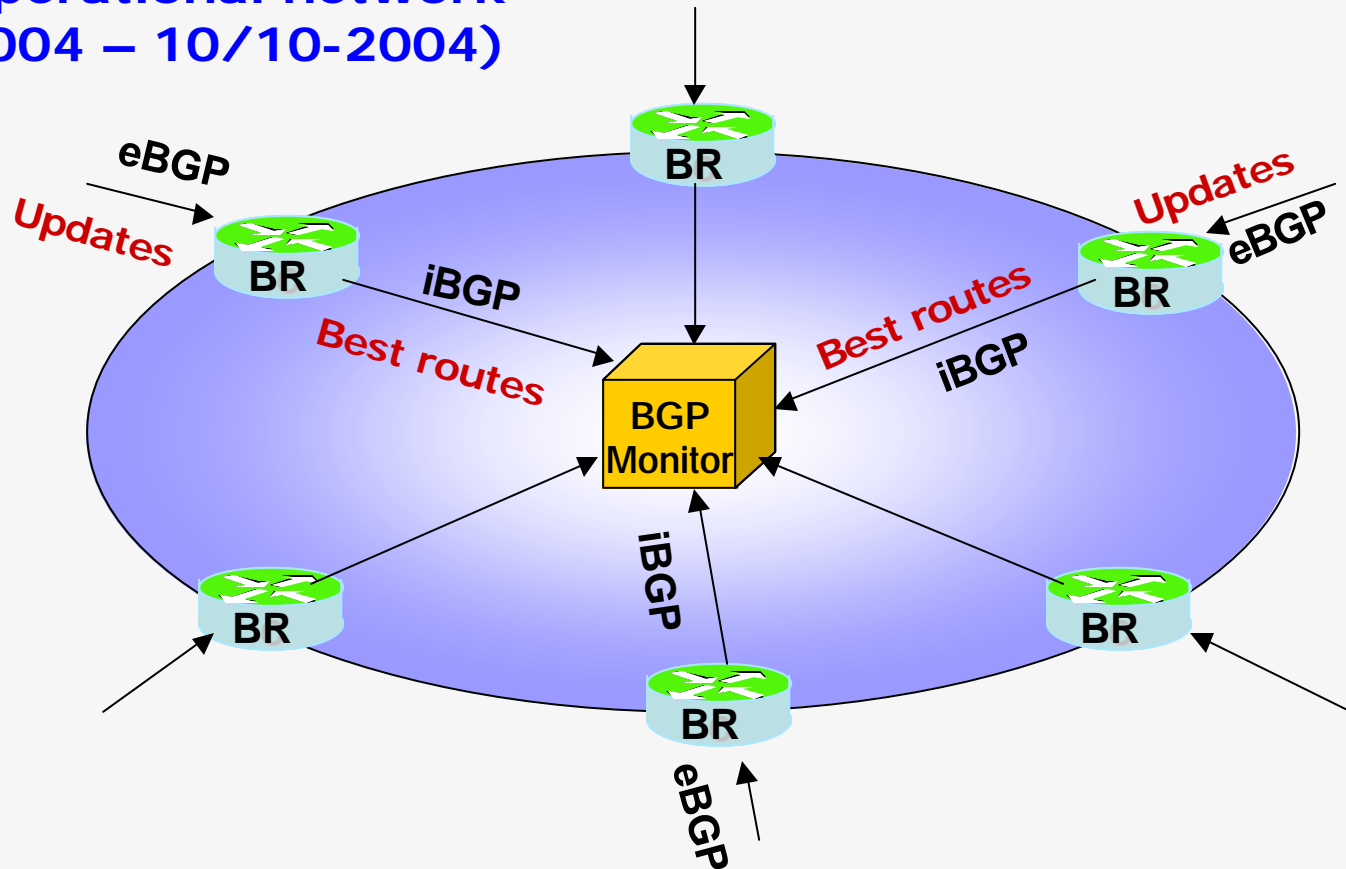
Interdomain Routing: Border Gateway Protocol



- Prefix-based: one route per prefix
- Path-vector: list of ASes in the path
- Incremental: every update indicates a change
- Policy-based: local ranking of routes

Capturing Routing Changes

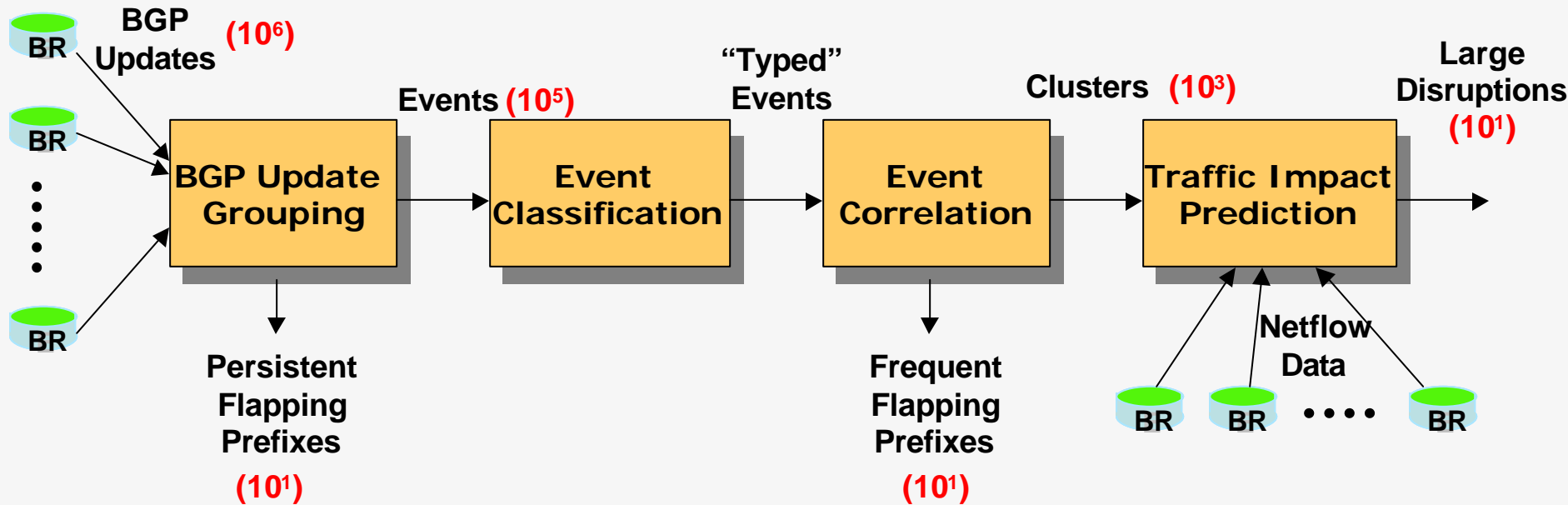
A large operational network
(8/16/2004 – 10/10-2004)



Challenges

- Large volume of BGP updates
 - Millions daily, very bursty
 - Too much for an operator to manage
- Different from root-cause analysis
 - Identify changes and their effects
 - Focus on actionable events rather than diagnosis
 - Diagnose causes in/near the AS

System Architecture

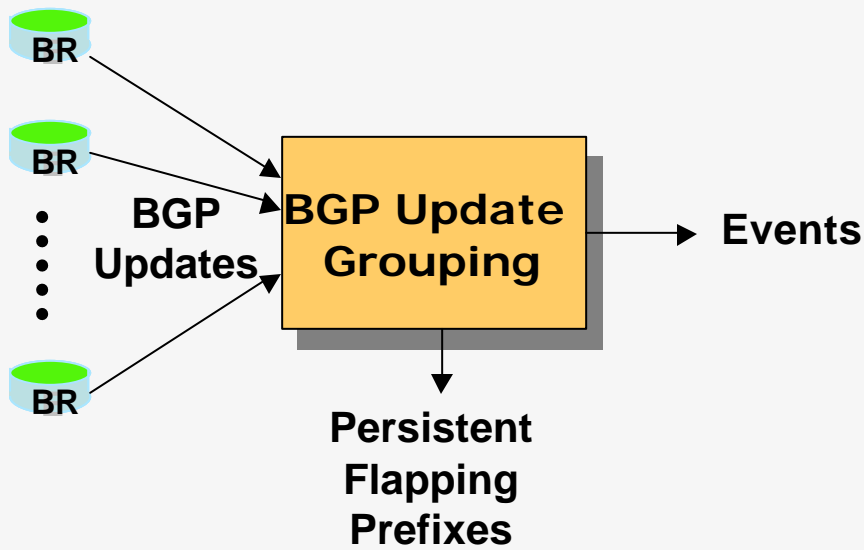


From millions of updates to a few dozen reports

Grouping BGP Update into Events

Challenge: A single routing change

- leads to multiple update messages
- affects routing decisions at multiple routers



Approach:

- Group together all updates for a prefix with inter-arrival < 70 seconds
- Flag prefixes with changes lasting > 10 minutes.

Grouping Thresholds

- Based on our understanding of BGP and data analysis
- Event timeout: 70 seconds
 - $2 * \text{MRAI timer} + 10 \text{ seconds}$
 - 98% inter-arrival time < 70 seconds
- Convergence timeout: 10 minutes
 - BGP usually converges within a few minutes
 - 99.9% events < 10 minutes

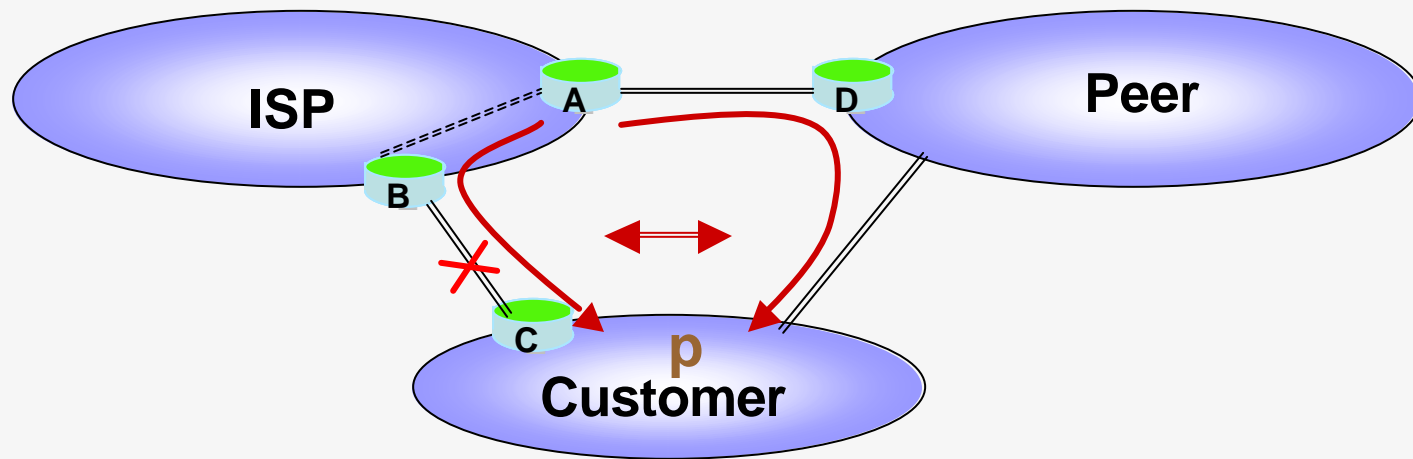
Persistent Flapping Prefixes

A surprising finding:

15.2% of updates were caused by persistent-flapping prefixes even though flap damping is enabled.

- Types of persistent flapping
 - Conservative damping parameters (78.6%)
 - Protocol oscillations due to MED (18.3%)
 - Unstable interfaces or BGP sessions (3.0%)

Example: Unstable eBGP Session

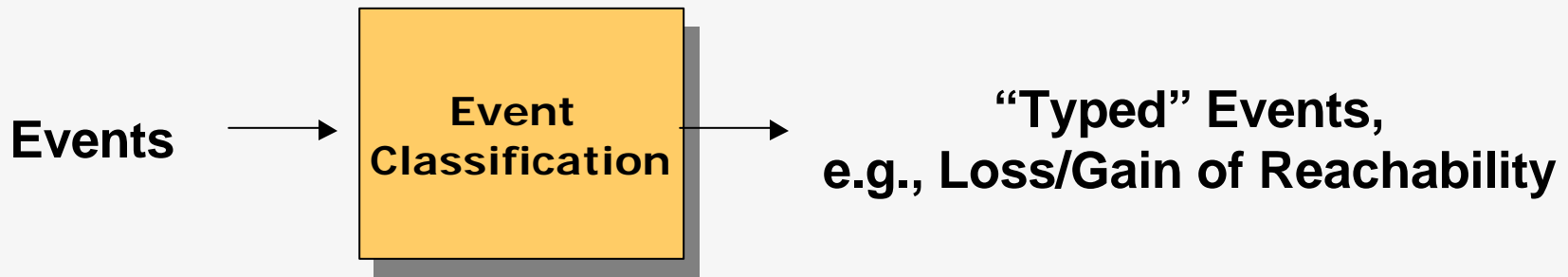


- Flap damping parameters is session-based
- Damping not implemented for iBGP sessions

Event Classification

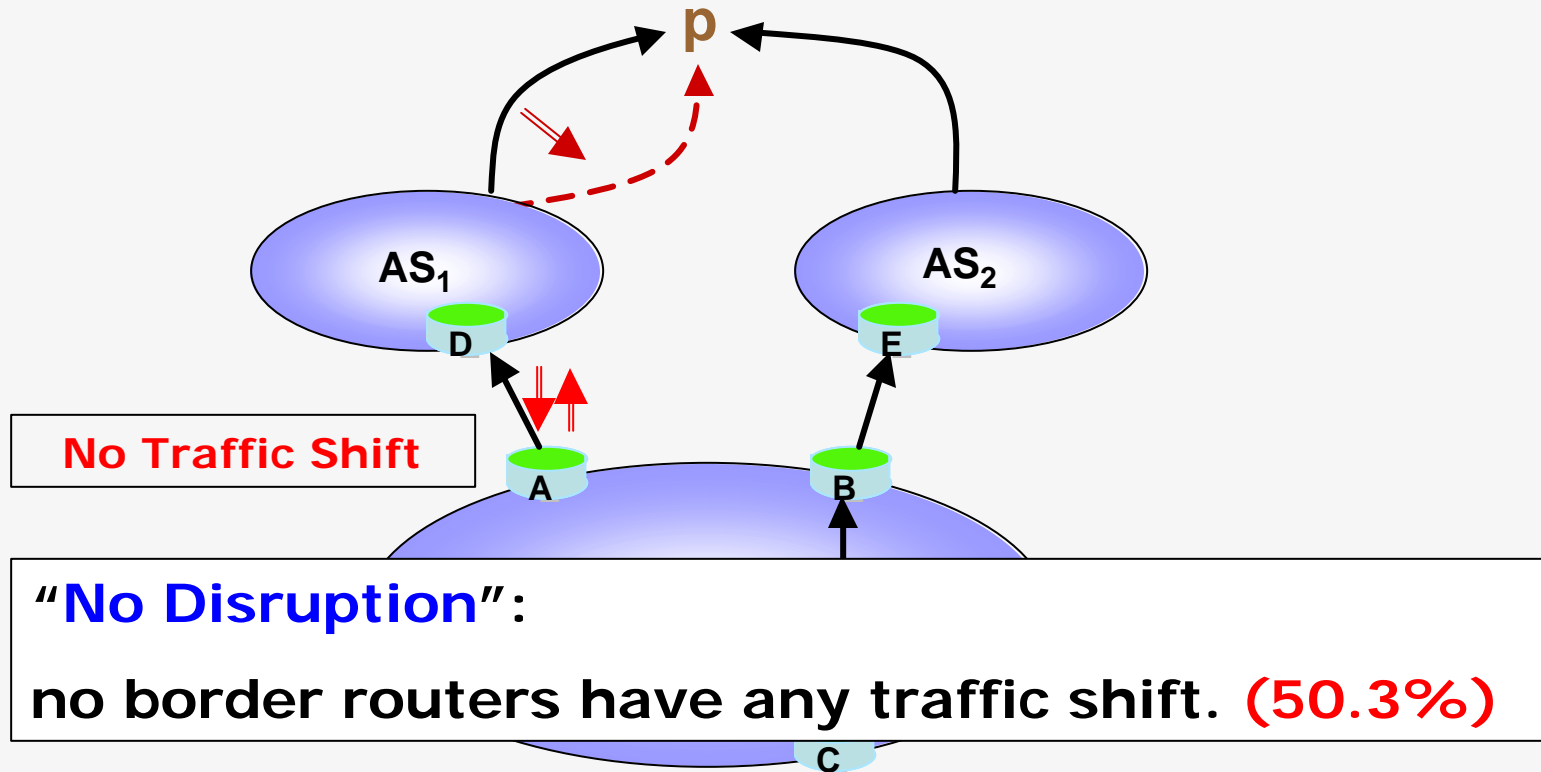
Challenge: Major concerns in network management

- Changes in reachability
- Heavy load of routing messages on the routers
- Change of flow of the traffic through the network

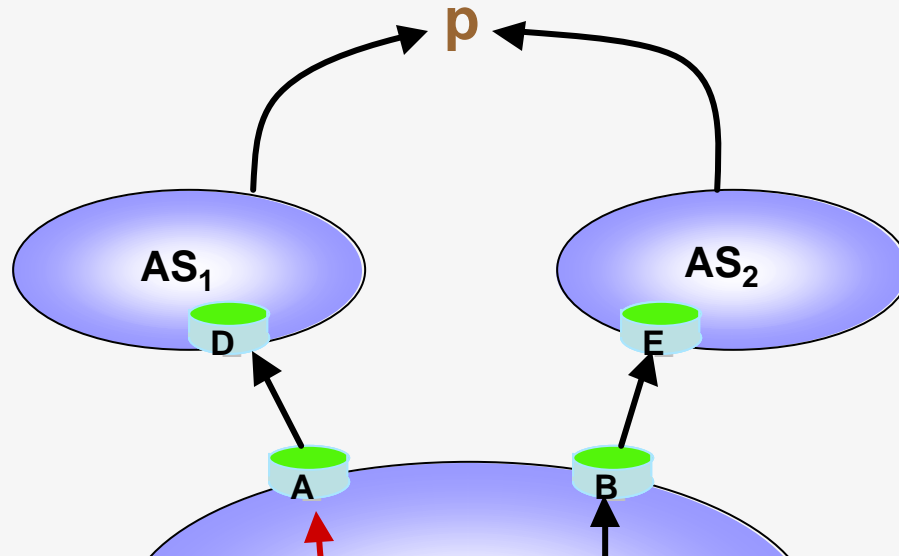


Solution: classify events by severity of their impact

Event Category – “No Disruption”



Event Category – “Internal Disruption”

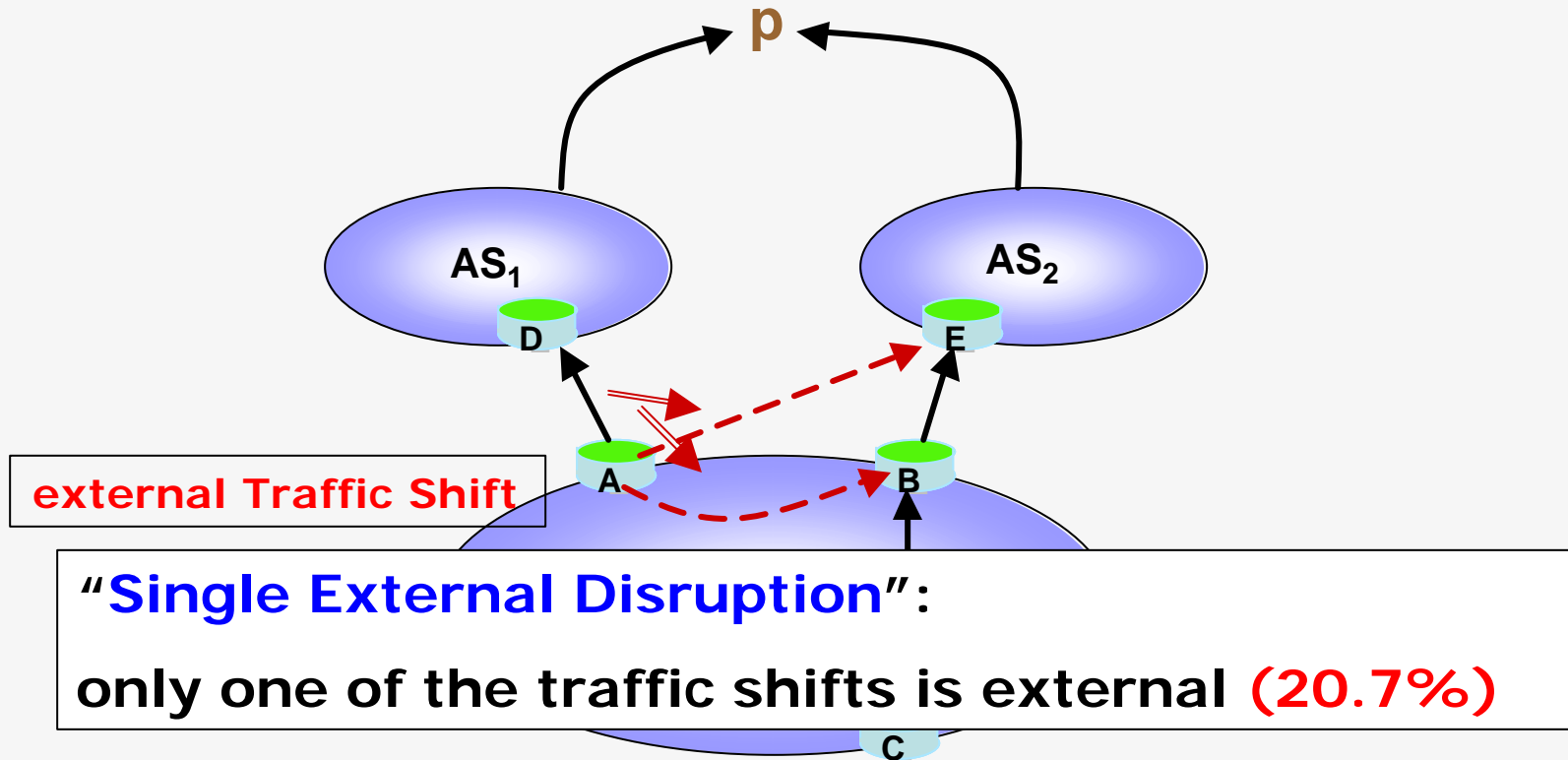


“Internal Disruption”:

all traffic shifts are internal. (15.6%)

Internal Traffic Shift

Event Category – “Single External Disruption”



Statistics on Event Classification

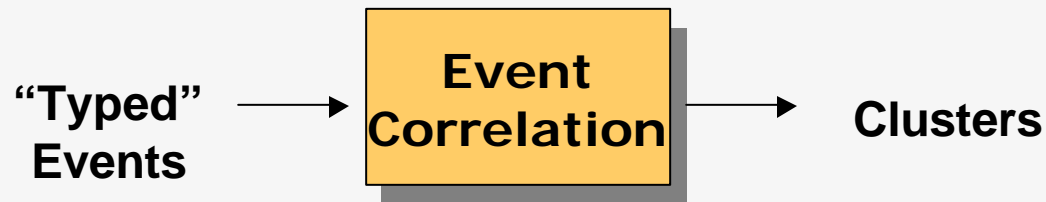
	Events	Updates
No Disruption	50.3%	48.6%
Internal Disruption	15.6%	3.4%
Single External Disruption	20.7%	7.9%
Multiple External Disruption	7.4%	18.2%
Loss/Gain of Reachability	6.0%	21.9%

- ❑ First 3 categories have significant day-to-day variations
- ❑ Updates per event depends on the type of events and the number of affected routers

Event Correlation

Challenge: A single routing change

- affects multiple destination prefixes

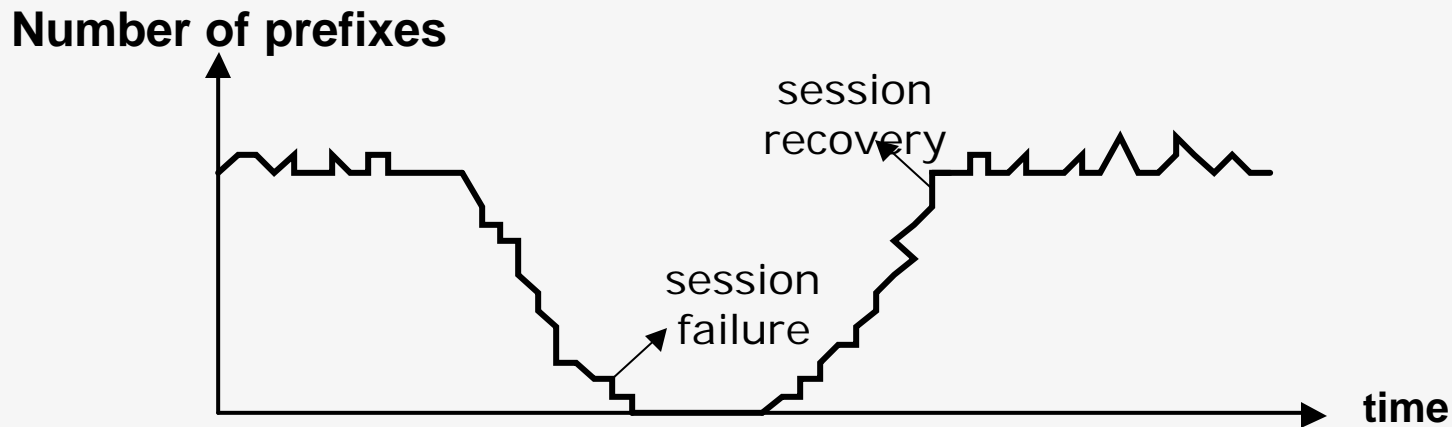


Solution:

group the same-type, close-occurring events

EBGP Session Reset

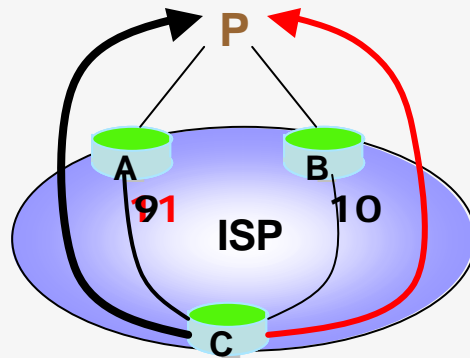
- ❑ Caused most of “single external disruption” events
- ❑ Check if the number of prefixes using that session as the best route changes dramatically



- ❑ Validation with Syslog router report (95%)

Hot-Potato Changes

□ Hot-Potato Changes

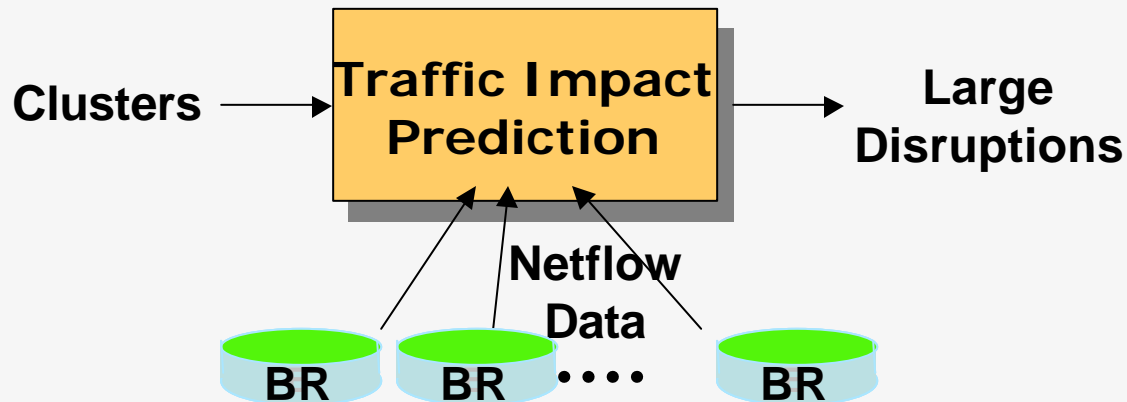


“Hot-potato routing” =
route to closest egress point

- Caused “internal disruption” events
- Validation with OSPF measurement (95%)
[Teixeira *et al* – SIGMETRICS’ 04]

Traffic Impact Prediction

Challenge: Routing changes have different impacts on the network which depends on the popularity of the destinations



Solution: weigh each cluster by traffic volume

Traffic Impact Prediction

- Traffic weight
 - Per-prefix measurement from netflow
 - 10% prefixes accounts for 90% of traffic
- Traffic weight of a cluster
 - the sum of “traffic weight” of the prefixes
 - A small number of large clusters have large traffic weight
 - Mostly session resets and hot-potato changes

Performance Evaluation

□ Memory

- Static memory: "current routes", 600 MB
- Dynamic memory: "clusters", 300 MB

□ Speed

- 99% of intervals of 1 second of updates can be process within 1 second
- Occasional execution lag
- Every interval of 70 seconds of updates can be processed within 70 seconds

Measurements were based on 900MHz CPU

Conclusion

- BGP troubleshooting system
 - Fast, online fashion
 - Operators' concerns (reachability, flapping, traffic)
 - Significant information reduction
 - millions of update → a few dozens of large disruptions
- Uncovered important network behavior
 - Hot-Potato changes
 - Session resets
 - Persistent-flapping prefixes