

15F-1 Bookkeeping

- 0 pts Correct

5F-2

$$\text{VC}(\text{do } c \text{ while } b, B) = \text{Inv} \wedge \text{VC}(c, \text{Inv}) \wedge (\forall x_1 \dots x_n. \text{Inv} \rightarrow (b \Rightarrow \text{VC}(c, \text{Inv}) \wedge \neg b \Rightarrow B))$$

5F-3

For the rule stark:

$$\begin{aligned} A &= (x = 0) \\ B &= (x \neq 0) \\ \sigma &= [x := 0] \\ \sigma' &= [x := 1] \\ c &= \text{while } x = 0 \text{ do } x := 1 \end{aligned}$$

This evaluates as shown in the following example with big-step semantics (Aexp and Bexp evaluation omitted for brevity)

$$\frac{\frac{\dots}{\langle x = 0, [x := 0] \rangle \Downarrow \text{true}} \quad \frac{\frac{\dots}{\langle x := 1, [x := 0] \rangle \Downarrow [x := 1]} \quad \frac{\frac{\dots}{\langle x = 0, [x := 1] \rangle \Downarrow \text{false}}}{\langle \text{while } x = 0 \text{ do } x := 1, [x := 1] \rangle \Downarrow [x := 1]}}{\langle x := 1; \text{while } x = 0 \text{ do } x := 1, [x := 1] \rangle \Downarrow [x := 1]}}{\langle \text{while } x = 0 \text{ do } x := 1, [x := 0] \rangle \Downarrow [x := 1]}$$

But with the condition X as both a precondition and postcondition in the rule, it is not possible for this to be proven using the rule stark unless A and B are equivalent, which is not the case here. The expression $\{x = 0 \wedge x = 0\} x := 1 \{x = 0\}$ is not provable.

The above example also works for the rule targaryen as the expression $\{x = 0\} x := 1 \{x = 0\}$ is not provable.

Neither of these rules account for the condition X no longer being satisfied after evaluation of c .

2 5F-2 VCGen Do-While

- 0 pts Correct

5F-2

$$\text{VC}(\text{do } c \text{ while } b, B) = \text{Inv} \wedge \text{VC}(c, \text{Inv}) \wedge (\forall x_1 \dots x_n. \text{Inv} \rightarrow (b \Rightarrow \text{VC}(c, \text{Inv}) \wedge \neg b \Rightarrow B))$$

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Neither of these rules account for the condition X no longer being satisfied after evaluation of c .

3 5F-3 VCGen Mistakes

- 0 pts Correct