

## 15F-1 Bookkeeping

- 0 pts Correct

SF-2

$$\text{Inv} \wedge \forall C (c, \text{Inv} \wedge (\forall x_i \text{ s.t. } \text{Inv} \Rightarrow (b \Rightarrow \forall C (c, \text{Inv} \wedge \neg b \Rightarrow P)))$$

where  $i \in [1, n]$

SF-3. select stark & targaryen

stark

$A = \{y \leq 10\}$

$B = \{y = 10\}$

$\sigma = \{y = 0\}$

$\sigma' = \{y = 10\}$

$c: \text{while } y < 10 \text{ do } y += 2$

easy to show  $\{c, \sigma\} \Downarrow \sigma'$   
 $\sigma \models A$   
 $\sigma' \models B$

Then it is not possible to prove  $\{A\} c \{B\}$ .

Because it should be

$$\frac{\vdash \{x \wedge b\} c \{x\}}{\vdash \{x\} \text{while } b \text{ do } c \{x \wedge b\}}$$
 stark

since only we have

$\{y \leq 10 \wedge \neg (y < 10)\}$   
 $= \{y \leq 10 \wedge y \geq 10\} = \{y = 10\}$

if we don't have that,  
then information we have  
can not prove  $\{A\} c \{B\}$ .

2 5F-2 VCGen Do-While

- 0 pts Correct

SF-2

$Inv \wedge \forall C (c, Inv \wedge (\forall x_i \text{ s.t. } Inv \Rightarrow (b \Rightarrow \forall C (c, Inv) \wedge b \Rightarrow P)))$

where  $i \in [1, n]$

SF-3. select stark & targaryen

stark

$A = \{y \leq 10\}$

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since only we have

$\{y \leq 10 \wedge \neg(y < 10)\}$   
 $= \{y \leq 10 \wedge y \geq 10\} = \{y = 10\}$

if we don't have that,  
then information we have  
can not prove  $\{A\} c \{B\}$ .

targaryen

$$A = \{y \leq 10\}$$

$$B = \{y = 10\}$$

$$\sigma = \{y = 0\}$$

$$\sigma' = \{y = 10\}$$

c: while  $y < 10$  do (if  $y < 10$  then  $y + 2$  else  $y = 11$ )

easy to show  $\{c, \sigma\} \Downarrow \sigma'$   
 $\sigma \models A$   
 $\sigma' \models B$

Then let's show it is impossible to show  $\vdash \{A\} c \{B\}$

since it should be

$$\frac{\vdash \{X \wedge b\} c \{X\}}{\vdash \{X\} \text{ while } b \text{ do } c \{X \wedge b\}} \text{ targaryen}$$

since if we not include  $\wedge b$ ,

then by the c, we may directly get  
 $y = 11$ , then it will break the while loop  
and we will not get to  $y = 10$  ( $\sigma'$ )

so, that is why it is impossible to  
show  $\vdash \{A\} c \{B\}$

### 3 5F-3 VCGen Mistakes

- 0 pts Correct