

## 15F-1 Bookkeeping

- 0 pts Correct

SF-2

I pick the first option.

do<sub>inv</sub> c while b

which is equivalent to c; while<sub>inv</sub> b do C

$VC(\text{do}_{\text{inv}} c \text{ while } b, B)$

$= VC(c; \text{while}_{\text{inv}} b \text{ do } c, B)$

$= VC(c, VC(\text{while}_{\text{inv}} b \text{ do } c, B))$

$= VC(c, \text{inv}) \wedge (\forall x_1 \dots x_n \text{ inv} \Rightarrow (b \Rightarrow VC(c, \text{inv}) \wedge b \Rightarrow B))$

## 2 5F-2 VCGen Do-While

- 0 pts Correct

5F-3.

$\frac{\vdash \{X \wedge b\} c \{X\}}{\vdash \{X\} \text{while } b \text{ do } c \{X\}}$  stark.

1. stark
2. A:  $n < 2$
3. B:  $n \leq 0$
4.  $\sigma$ :  $\sigma(n) = 1$
5.  $\sigma'$ :  $\sigma'(n) = 0$
6. c: while  $n > 0$  do  $n = n - 1$
7.  $\langle \langle \sigma \rangle \vee \sigma' \rangle$ :  $n = 0$  after  $n = n - 1$   
then jump out of while loop
8.  $\sigma \models A$ :  $\sigma \models n < 2$
9.  $\sigma' \models B$ :  $\sigma' \models n \leq 0$
10. It's not possible to prove  $\{n < 2\} \text{while } n > 0 \text{ do } n = n - 1 \{n \leq 0\}$   
using stark rule.  
According to stark rule:  $\vdash \{X\} \text{while } b \text{ do } c \{X\}$   
but there is no X satisfies the (A) c (B) above  
 $n < 2 \Rightarrow X \wedge X \Rightarrow n \leq 0$  is unsatisfiable.

$\frac{\vdash \{X\} c \{X\}}{\vdash \{X\} \text{while } b \text{ do } c \{X \wedge \neg b\}}$  tergyaryen

1. tergyaryen
2. A:  $n \geq 0$
3. B:  $n = 0$
4.  $\sigma$ :  $\sigma(n) = 1$
5.  $\sigma'$ :  $\sigma'(n) = 0$
6. c: while  $n > 0$  do  $n = n - 1$
7.  $\langle \langle \sigma \rangle \vee \sigma' \rangle$ :  $n = 0$  after  $n = n - 1$   
then jump out of while loop
8.  $\sigma \models A$ :  $\sigma \models n \geq 0$
9.  $\sigma' \models B$ :  $\sigma' \models n = 0$
10. It is not possible to prove  $\vdash \{n \geq 0\} \text{while } n > 0 \text{ do } n = n - 1 \{n = 0\}$   
X should satisfy:  $\{X\} n = n - 1 \{X\}$  ①  
 $n \geq 0 \Rightarrow X$  ②  
 $X \wedge \neg(n > 0) \Rightarrow n = 0$  ③  
from ③, we get a constrain:  $X \Rightarrow n > 0$   $\left( \begin{array}{l} X \wedge n = 0 \Rightarrow n = 0 \\ \exists X \vee n > 0 \vee n = 0 \\ \exists X \Rightarrow n > 0 \end{array} \right)$   
Combined with ①, X must be  $n > 0$   
 $\{n > 0\} \text{while } n > 0 \{n = 0\}$  is wrong.  
So, there is no X with this rule can prove  $\{A\} c \{B\}$ .

### 3 5F-3 VCGen Mistakes

- 0 pts Correct