1 5F-1 Bookkeeping

- **0 pts** Correct

Exercise 5F-2. VCGen Do-While As I really never learned how to think about a do-while loop, I'll convert directly from "do c while b" to "c; while b do c".

So do_{inv} c while b becomes:
c; $while_{inv}$ b do c
Let's drop this directly into a VC and follow a similar method as we used for the "vanilla" while rule:

VC(c; $while_{inv}$ b do c,B)

To read this we might say in english: The post condition B holds given the two commands given.
We can drop the while VCGen from the slides in for $while_{inv}$ b do c:
VC(c; INV $\wedge(\forall$ x$_1$...x_2. INV $\rightarrow (b \rightarrow (VC(c, INV) \wedge \neg b \rightarrow B)))$

We have a nested VC, so we can simplify this by splitting on the "and" ($\wedge$) dividing our first and second commands:

$$VC(c; INV) \wedge (\forall x_1 ... x_2 . INV \rightarrow (b \rightarrow (VC(c, INV) \wedge \neg b \rightarrow B))$$

We realized that do-while is just c; while in disguise, so we made like Scooby-Doo and unmasked it!

Exercise 5F-3. VCGen Mistakes I choose Targaryen and Stark because these rules both fail to modify the post condition.

Exercise 5F-3: Targaryen

1. Targaryen
2. A = true
   We don't really care about A.
3. B = (x = 42)
   Some value that should occur when we exit the loop.
4. $\sigma(x) = 0$
   IMP style initialization.
5. $\sigma'(x) = 42$
   The expected final value.
6. c = while $x < 1$ do x := 42
   Some command to intentionally make this fail.
7. $$\frac{\langle x < 1, \sigma\rangle \Downarrow true \quad \langle x := 42, \sigma\rangle \Downarrow \sigma' \quad \langle while \;\; x < 1 \;\; do \;\; x := 42, \sigma'\rangle \Downarrow \sigma''}{\langle while \;\; x < 1 \;\; do \;\; x := 42, \sigma\rangle \Downarrow \sigma''}$$

2 5F-2 VCGen Do-While

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If the guard is false do nothing:

$$\frac{\langle x < 1, \sigma\rangle \Downarrow false}{\langle while\ \ x < 1\ \ do\ \ x := 42, \sigma\rangle \Downarrow \sigma}$$

σ'' is our σ'.

8. A is just true, so $\sigma \models A$
9. $\sigma' \models B$ Check this by noticing that the loop does one iteration and ends on x := 42. Therefore $\sigma'(x) = 42$
10. $\vdash$ {A} c {B} is not possible because the rule is assuming pre and post conditions are the same (X = X). We can just check A = B and find true = 42. This is clearly not correct. The rule only works in the case the loop body is never encountered (if b is false coming in).

The rule should consider a correct post condition, and should have the loop guard considered coming in.

Exercise 5F-3: Stark Stark, as usual, is missing some key information. A correct rule probably would have seen the Red Wedding coming. Let's see where the strategy could be improved:

1. Stark
2. A = true
   We don't really care about A, it will simply show the modification made in the loop isn't correctly tracked by the post condition.
3. B = (x = 42)
   Some value that should occur when we exit the loop.
4. $\sigma(x) = 0$
   IMP style initialization.
5. $\sigma'(x) = 42$
   The expected final value.
6. c = while $x < 1$ do x := 42
   Some command to intentionally make this fail.
7. $$\frac{\langle x < 1, \sigma\rangle \Downarrow true \quad \langle x := 42, \sigma\rangle \Downarrow \sigma' \quad \langle while\ \ x < 1\ \ do\ \ x := 42, \sigma'\rangle \Downarrow \sigma''}{\langle while\ \ x < 1\ \ do\ \ x := 42, \sigma\rangle \Downarrow \sigma''}$$

   If the guard is false do nothing:

   $$\frac{\langle x < 1, \sigma\rangle \Downarrow false}{\langle while\ \ x < 1\ \ do\ \ x := 42, \sigma\rangle \Downarrow \sigma}$$

   note that I add an extra step here to be explicit, but σ'' is our σ'.

8. A is just true, so $\sigma \models A$

9. $\sigma' \models B$ Check this by noticing that the loop does one iteration and ends on x := 42. Therefore $\sigma'(x) = 42$

10. $\vdash$ {A} c {B} is not possible because the rule is assuming pre and post conditions are the same (X = X). We can just check A = B and find true = 42. This is clearly not correct. The rule only works in the case the loop body is never encountered (if b is false coming in).

Lannister seems reasonable but maybe a little needlessly complex.

3 5F-3 VCGen Mistakes

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