

15F-1 Bookkeeping

- 0 pts Correct

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Exercise 5F-2:

$$VC(\text{do}_{\text{Inv}} c \text{ while } b, P)$$

$$= VC(c; \text{while}_{\text{Inv}} b \text{ do } c, P)$$

$$= VC(c, VC(\text{while}_{\text{Inv}} b \text{ do } c, P))$$

$$= VC(c, \text{Inv} \wedge (\exists X_1 \dots X_n. \text{Inv} \rightarrow c b \rightarrow VC(c, \text{Inv}) \wedge \neg b \rightarrow P)).$$

2 5F-2 VCGen Do-While

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Exercise 5F-3:

⊖. I select the stark rule. Compared with the original while rule, stark rule is incomplete because it is missing the " $\neg b$ " check when the loop finishes.

1. stark

2. Let A be $\{x=0\}$

3. Let B be $\{x=1\}$

4. $\sigma(x) = 0$

5. $\sigma'(x) = 1$

6. c is while $x \neq 1$ do $x := 1$

7. $\langle c, \sigma \rangle \Downarrow \sigma'$ because $\sigma(x) = 0$, c will assign 1 to x, and thus $\sigma'(x) = 1$

8. $\sigma \models A$ because $\sigma(x) = 0$ and A is $\{x=0\}$

9. $\sigma' \models B$ because $\sigma'(x) = 1$ and B is $\{x=1\}$

10. it is not possible to prove $\vdash \{A\} c \{B\}$ using the stark rule.
we have

$$\frac{\vdash \{x \wedge b\} c \{x\}}{\vdash \{x\} \text{while } b \text{ do } c \{x\}}$$

where b is the assumption $x \neq 1$ and c is the loop body $x := 1$

We cannot assume $\neg b$, which is $x=1$, when the loop finishes in the stark rule. Then, given A, or $\{x=0\}$ which satisfies $x \neq 1$, we cannot make a derivation that when the loop terminates, we have B, or $\{x=1\}$.

②. I select the Hoare rule. The problem of this variant is that it doesn't let you assume b when executing the loop body c in $\text{while } b \text{ do } c$.

1. Hoare

2. A is $\{x \leq 0\}$

3. B is $\{x = 6\}$

4. $\sigma(x) = 0$

5. $\sigma'(x) = 6$

6. c is $\text{while } x \leq 5 \text{ do } x := x + 1$

7. $\langle c, \sigma \rangle \Downarrow \sigma'$ by apply the while loop on $x = 0$.

8. $\sigma \models A$ because $\sigma(x) = 0$ and A is $\{x \leq 0\}$.

9. $\sigma' \models B$ because $\sigma'(x) = 6$ and B is $\{x = 6\}$.

10. it is not possible to prove $\vdash \{A\} c \{B\}$ using the Hoare rule.

We first use invariant $x \leq 6$.

$$\frac{\vdash x \leq 6 \wedge x \leq 5 \rightarrow x + 1 \leq 6 \quad \overline{\vdash \{x + 1 \leq 6\} x := x + 1 \{x \leq 6\}}}{\text{fail} \Rightarrow \vdash \{x \leq 6 \wedge x \leq 5\} x := x + 1 \{x \leq 6\}}$$

$$\vdash \{x \leq 6\} c \{x \leq 6 \wedge x > 5\}$$

Then we finish off with consequence: $\vdash x \leq 0 \rightarrow x \leq 6$

$$\frac{\vdash x \leq 6 \wedge x > 5 \rightarrow x = 6 \quad \vdash \{x \leq 6\} c \{x \leq 6 \wedge x > 5\}}{\vdash \{x \leq 0\} c \{x = 6\}}$$

Compared with the correct while loop, the part written in red and enclosed by a red box is missing. Thus, the left part fails. There is no way to make a derivation for the $\vdash \{x \leq 6\} x := x + 1 \{x \leq 6\}$. This line fails when $x = 6$.

So Hoare is incomplete.

3 5F-3 VCGen Mistakes

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