

**Exercise 5F-2.**

$VC(\text{do}_{Inv1, Inv2} \ c \text{ while } b, P) = Inv1 \wedge (\forall x_1 \dots x_n. Inv1 \implies ((b \implies VC(c, Inv2)) \wedge (\neg b \implies P))) \wedge (\forall x_1 \dots x_n. Inv2 \implies ((b \implies VC(c, Inv2)) \wedge (\neg b \implies P)))$

Where  $x_1 \dots x_n$  are all variables modified in  $c$ .

**Exercise 5F-3.****Rule 1**

1. stark
2.  $A := (x = 0)$
3.  $B := (x = 1)$
4.  $\sigma := \{x = 0\}$
5.  $\sigma' := \{x = 1\}$
6.  $c := \text{while } x < 1 \text{ do } x := x + 1$
7.  $\langle \text{while } x < 1 \text{ do } x := x + 1, \{x = 0\} \rangle \Downarrow \{x = 1\}$
8.  $\{x = 0\} \models (x = 0)$
9.  $\{x = 1\} \models (x = 1)$
10. Since stark has the same conditions on each side of the conclusion about while, and since it's not true that  $x = 1$  before  $c$ , we cannot prove that  $x = 1$  after  $c$ .

**Rule 2**

1. targaryen
2.  $A := (x = 0)$
3.  $B := (x = 1)$
4.  $\sigma := \{x = 0\}$
5.  $\sigma' := \{x = 1\}$
6.  $c := \text{while } x < 1 \text{ do } x := x + 1$

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7.  $\langle \text{while } x < 1 \text{ do } x := x + 1, \{x = 0\} \rangle \Downarrow \{x = 1\}$
8.  $\{x = 0\} \models (x = 0)$
9.  $\{x = 1\} \models (x = 1)$
10. To prove  $\{x = 0\} \text{ while } x < 1 \text{ do } x := x + 1 \{x = 1\}$  using targaryen, we would need to find some predicate  $X$  such that  $\{X\} x := x + 1 \{X\}$  and  $(X \wedge \neg x = 0) \implies x = 1$ . From the Hoare rule for assignment, we get  $[x + 1/x]X = X$ , meaning  $X$  does not refer to  $x$ , so  $(X \wedge \neg x = 0)$  cannot prove  $x = 1$ .