

14F-1 Bookkeeping

- 0 pts Correct

**Exercise 4F-2. VCGen for Let [6 points].** The bug in the `let` rule is that it does not reassign the initial value of  $x$  back to  $x$  after the command is finished. In order to ensure the original value is assigned back to  $x$ , we can provide the rule for `let` as a sequence of commands where we:

1. Store the initial value of  $x$  to a new variable  $a$  that hasn't been used yet and is not used in  $e$  and  $c$ .
2. Assign  $e$  to  $x$ .
3. Execute  $c$ .
4. Reassign  $a$ , which stores the original value of  $x$ , to  $x$ .

Below is a correct rule for `let`, where  $a$  is a variable that is not in scope before the `let` command is executed and is not used in  $e$  and  $c$ :

$$\begin{aligned}
 \text{VC}(\text{let } x = e \text{ in } c, B) &= \text{VC}(a := x ; x := e ; c ; x := a, B) \\
 &= \text{VC}(a := x, \text{VC}(x := e, \text{VC}(c, \text{VC}(x := a, B)))) \\
 &= \text{VC}(a := x, \text{VC}(x := e, \text{VC}(c, [a/x] B))) \\
 &= \text{VC}(a := x, [e/x] \text{VC}(c, [a/x] B)) \\
 &= [x/a] [e/x] \text{VC}(c, [a/x] B)
 \end{aligned}$$

2 4F-2 VCGen for Let

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**Exercise 4F-3. VCGen Mistakes [6 points].**

1. Command:  $c = \text{let } x = y + 1 \text{ in skip}$
2. Post-condition:  $B = \{x > 0\}$
3. State:  $\sigma = [x := 0, y := 0]$
4. We see that  $\sigma \models VC(c, B)$ :

$$\begin{aligned}
 & VC(\text{let } x = y + 1 \text{ in skip}, \{x > 0\}) \\
 &= [y + 1/x] VC(\text{skip}, \{x > 0\}) \\
 &= [y + 1/x] \{x > 0\} \\
 &= \{y + 1 > 0\},
 \end{aligned}$$

and because  $\sigma(y) = 0$ ,  $\sigma(y) + 1 = 1 > 0$ .

5. We apply the operational semantics rule for **let** to obtain  $\sigma'$ . We have  $\sigma' = \sigma = [x := 0, y := 0]$  because the original value of  $\sigma(x) = 0$  is restored after the **let** command is finished, and no variables were modified inside the body of the **let** command.

$$\frac{\frac{\frac{\langle y, \sigma \rangle \Downarrow 0 \quad \langle 1, \sigma \rangle \Downarrow 1}{\langle y + 1, \sigma \rangle \Downarrow 1}}{\langle x, \sigma \rangle \Downarrow 0} \quad \frac{\langle x := y + 1, \sigma \rangle \Downarrow \sigma[x := 1]}{\langle \text{skip}, \sigma[x := 1] \rangle \Downarrow \sigma[x := 1]}}{\langle \text{let } x = y + 1 \text{ in skip}, \sigma \rangle \Downarrow \sigma}$$

6.  $\sigma' \not\models B = \{x > 0\}$  because  $\sigma'(x) = 0 \not> 0$ .

### 3 4F-3 VCGen Mistakes

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**Exercise 4F-4. Axiomatic Do-While [6 points].** We present the following Hoare rule for `do c while b`:

$$\frac{\vdash \{A\} c \{B\} \quad \vdash \{B\} \text{ while } c \text{ do } b \{B \wedge \neg b\}}{\vdash \{A\} \text{ do } c \text{ while } b \{B \wedge \neg b\}}$$

In this rule, we execute  $c$  with pre-condition  $\{A\}$  and post-condition  $\{B\}$ . After executing  $c$  once, the rest of the `do c while b` command reduces to a normal `while b do c` command. Thus, the post-condition  $\{B\}$  after executing  $c$  serves as the loop invariant for the rest of the `while` loop. Using the Hoare rule for `while`, we know that the post-condition of the rest of the loop should be  $\{B \wedge \neg b\}$ , the invariant  $B$  with the additional condition that the boolean expression  $b$  is false. Overall, we denote this sequence of operations to have precondition  $\{A\}$  with the post-condition after executing the rest of the loop  $\{B \wedge \neg b\}$ .

#### 4 4F-4 Axiomatic Do-While

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