

## 14F-1 Bookkeeping

- 0 pts Correct

Peer Review ID: 70884333 — enter this when you fill out your peer evaluation via gradescope

#### Exercise 4F-2

The original *let* rule doesn't consider the scope where  $x$  is assigned to  $A \text{exp } e$ . We can deduct a correct *let* rule by creating a new variable  $temp$ , assigning  $temp$  to the original value of  $x$  first, and reassigning  $x$  back to its original value  $temp$  after executing  $c$ . This means evaluating  $VC(\text{let } x = e \text{ in } c, B)$  is equivalent to evaluating  $VC(temp := x; x := e; c; x := temp, B)$ . Applying other given rules, then we have:

$$\begin{aligned} VC(temp := x; x := e; c; x := temp, B) &= [x/temp] VC(x := e; c; x := temp, B) \\ &= [x/temp] [e/x] VC(c; x := temp, B) \\ &= [x/temp] [e/x] VC(c, VC(x := temp, B)) \\ &= [x/temp] [e/x] VC(c, [temp/x] B) \end{aligned}$$

$$\text{Thus, } VC(\text{let } x = e \text{ in } c, B) = [x/temp] [e/x] VC(c, [temp/x] B)$$

#### Exercise 4F-3

To show that the buggy *let* rule is unsound, we set:

1.  $c == \text{let } x = 3 \text{ in } y := y - x$

2.  $B == x = y$

3.  $\sigma(x) = 1$  and  $\sigma(y) = 6$

4. We know that  $\sigma \models VC(c, B)$  because

$$\begin{aligned} VC(\text{let } x = 3 \text{ in } y := y - x, x = y) &= [3/x] VC(y := y - x, x = y) \\ &= [3/x] [y - x/y] x = y \\ &= [3/x] x = y - x \\ &= (3 = y - 3) \end{aligned}$$

since  $\sigma(y) = 6$ , we have  $3 = 3$ . Thus,  $VC(c, B)$  is true in  $\sigma$ .

5. According to  $\langle c, \sigma \rangle \Downarrow \sigma'$ , we have  $\sigma'(x) = 1$  and  $\sigma'(y) = 3$

6.  $\sigma' \not\models B$  because in  $\sigma'$  we have  $\sigma'(x) = 1$ ,  $\sigma'(y) = 3$ , but  $1 \neq 3$ .

Hence, we demonstrate that the buggy *let* rule can prove a false thing, and thus it is unsound.

2 4F-2 VCGen for Let

- 0 pts Correct

#### Exercise 4F-2

The original *let* rule doesn't consider the scope where  $x$  is assigned to  $A \text{exp } e$ . We can deduct a correct *let* rule by creating a new variable  $temp$ , assigning  $temp$  to the original value of  $x$  first, and reassigning  $x$  back to its original value  $temp$  after executing  $c$ . This means evaluating  $VC(\text{let } x = e \text{ in } c, B)$  is equivalent to evaluating  $VC(temp := x; x := e; c; x := temp, B)$ . Applying other given rules, then we have:

$$\begin{aligned} VC(temp := x; x := e; c; x := temp, B) &= [x/temp] VC(x := e; c; x := temp, B) \\ &= [x/temp] [e/x] VC(c; x := temp, B) \\ &= [x/temp] [e/x] VC(c, VC(x := temp, B)) \\ &= [x/temp] [e/x] VC(c, [temp/x] B) \end{aligned}$$

$$\text{Thus, } VC(\text{let } x = e \text{ in } c, B) = [x/temp] [e/x] VC(c, [temp/x] B)$$

#### Exercise 4F-3

To show that the buggy *let* rule is unsound, we set:

1.  $c == \text{let } x = 3 \text{ in } y := y - x$

2.  $B == x = y$

3.  $\sigma(x) = 1$  and  $\sigma(y) = 6$

4. We know that  $\sigma \models VC(c, B)$  because

$$\begin{aligned} VC(\text{let } x = 3 \text{ in } y := y - x, x = y) &= [3/x] VC(y := y - x, x = y) \\ &= [3/x] [y - x/y] x = y \\ &= [3/x] x = y - x \\ &= (3 = y - 3) \end{aligned}$$

since  $\sigma(y) = 6$ , we have  $3 = 3$ . Thus,  $VC(c, B)$  is true in  $\sigma$ .

5. According to  $\langle c, \sigma \rangle \Downarrow \sigma'$ , we have  $\sigma'(x) = 1$  and  $\sigma'(y) = 3$

6.  $\sigma' \not\models B$  because in  $\sigma'$  we have  $\sigma'(x) = 1$ ,  $\sigma'(y) = 3$ , but  $1 \neq 3$ .

Hence, we demonstrate that the buggy *let* rule can prove a false thing, and thus it is unsound.

### 3 4F-3 VCGen Mistakes

- 0 pts Correct

Exercise 4F-4

We know  $c$  is executed once before  $b$  is evaluated. Thus, we can derive Hoare rule for *do while* according to the *while* rule:

$$\vdash \{A\} c \{B\} \quad \vdash \{B\} \text{ while } b \text{ do } c \{C\}$$

---

$$\vdash \{A\} \text{ do } c \text{ while } b \{C\}$$

Then, we can further derive it to:

$$\vdash \{A\} c \{B\} \quad \vdash \{B\} \text{ while } b \text{ do } c \{B \wedge \neg b\}$$

---

$$\vdash \{A\} \text{ do } c \text{ while } b \{B \wedge \neg b\}$$

Finally, we will have:

$$\vdash \{A\} c \{B\} \quad \vdash \{B \wedge b\} c \{B\}$$

---

$$\vdash \{A\} \text{ do } c \text{ while } b \{B \wedge \neg b\}$$

#### 4 4F-4 Axiomatic Do-While

- 0 pts Correct