

14F-1 Bookkeeping

- 0 pts Correct

Exercise 4F-2. VCGen for Let [6 points]. In class we gave the following rules for the (backward) verification condition generation of assignment and let:

$$\begin{aligned} \text{VC}(c_1; c_2, B) &= \text{VC}(c_1, \text{VC}(c_2, B)) \\ \text{VC}(x := e, B) &= [e/x] B \\ \text{VC}(\text{let } x = e \text{ in } c, B) &= [e/x] \text{VC}(c, B) \end{aligned}$$

That rule for let has a bug. Give a correct rule for let.

Answer: The new let rule is as follows:

$$\text{VC}(\text{let } x = e \text{ in } c, B) = [x/e] \text{VC}(c, [e/x]B)$$

This rule replaces the bound variable x in let with e in B . This allows the substitution to be scoped to the let expression only. After the VC has been generated, we must replace e with x to reverse this replacement. This is because e does not have any meaning outside of the let expression and therefore needs to be substituted back.

2 4F-2 VCGen for Let

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Exercise 4F-3. VCGen Mistakes [6 points]. Given $\{A\}c\{B\}$ we desire that $A \implies \text{VC}(c, B) \implies \text{WP}(c, B)$. We say that our VC rules are *sound* if $\models \{\text{VC}(c, B)\} c \{B\}$. Demonstrate the unsoundness of the buggy let rule by giving the following six things:

1. a command c and
2. a post-condition B and
3. a state σ such that
4. $\sigma \models \text{VC}(c, B)$ and
5. $\langle c, \sigma \rangle \Downarrow \sigma'$ but
6. $\sigma' \not\models B$.

Answer:

1. $c = \text{let } x = 2 \text{ in } x := x + 2$
 2. $B = \{x \geq 4\}$
 3. $\sigma = \{x = 3\}$
 4. We can compute $\text{VC}(c, B) = [2/x]\text{VC}(x := x + 2, B) = [2/x][x + 2/x]B = [2 + 2/x]B = [4/x]\{x \geq 4\} = \{4 \geq 4\} = \{\text{true}\}$. We know that for any $\sigma \models \text{true}$ always. Therefore, we can conclude that $\sigma \models \text{VC}(c, B)$.
 5. We have that $\langle c, \sigma \rangle \Downarrow \sigma'$ where $\sigma' = \sigma = \{x = 3\}$. σ is unaffected by c because x in let only affects x within the let command.
 6. We can see that $\sigma' \not\models B$ as $3 \not\geq 4$.
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3 4F-3 VCGen Mistakes

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Exercise 4F-4. Axiomatic Do-While [6 points]. Write a sound and complete Hoare rule for `do c while b`. This statement has the standard semantics (e.g., c is executed at least once, before b is tested).

Answer: This new rule uses the `while` rule found on slide #23 of lecture 8 (Introduction to Axiomatic Semantics):

$$\frac{\{A\} c \{C\} \quad \{C\} \text{ while } b \text{ do } c \{B\}}{\{A\} \text{ do } c \text{ while } b \{B\}}$$

This rule leverages the fact that a do-while loop is identical to a while-do loop, except it executes c once before checking the b condition. We can therefore compose a rule for it (like the one above) by executing c once and then using the same semantics as a while-do loop.

4 4F-4 Axiomatic Do-While

- 0 pts Correct