

1 4F-1 Bookkeeping

- 0 pts Correct

**Exercise 4F-2. VCGen for Let [6 points].** In class we gave the following rules for the (backward) verification condition generation of assignment and let:

$$\begin{aligned} \text{VC}(c_1; c_2, B) &= \text{VC}(c_1, \text{VC}(c_2, B)) \\ \text{VC}(x := e, B) &= [e/x] B \\ \text{VC}(\text{let } x = e \text{ in } c, B) &= [e/x] \text{VC}(c, B) \end{aligned}$$

That rule for let has a bug. Give a correct rule for let.

First unwind let as follow

$$\text{let } x = e \text{ in } c = t := x; x := e; c; x := t$$

Therefore

$$\begin{aligned} \text{VC}(\text{let } x = e \text{ in } c, B) &= \text{VC}(t := x; x := e; c; x := t, B) \\ &= [x/t][e/x]\text{VC}(c; x := t, B) \\ &= [x/t][e/x]\text{VC}(c; \text{VC}(x := t, B)) \\ &= [x/t][e/x]\text{VC}(c; [t/x]B) = \text{VC}([e/x]c, B) \end{aligned}$$

Hence the correct rule is

$$\text{VC}(\text{let } x = e \text{ in } c, B) = \text{VC}([e/x]c, B)$$

2 4F-2 VCGen for Let

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**Exercise 4F-3. VCGen Mistakes [6 points].** Given  $\{A\}c\{B\}$  we desire that  $A \implies \text{VC}(c, B) \implies \text{WP}(c, B)$ . We say that our VC rules are *sound* if  $\models \{\text{VC}(c, B)\} c \{B\}$ . Demonstrate the unsoundness of the buggy let rule by giving the following six things:

1. a command  $c$  and
2. a post-condition  $B$  and
3. a state  $\sigma$  such that
4.  $\sigma \models \text{VC}(c, B)$  and
5.  $\langle c, \sigma \rangle \Downarrow \sigma'$  but
6.  $\sigma' \not\models B$ .

Demonstration:

1.  $c = \text{let } x = 5 \text{ in skip}$
2.  $B$  is  $x = 5$
3.  $\sigma(x) = 0$
4.  $\sigma \models \text{VC}(c, B)$  because  $\text{VC}(c, B)$  is  $5 = 5$
5.  $\langle c, \sigma \rangle \Downarrow \sigma'$ , and we know  $\sigma'(x) = 0$  based on the property of let
6.  $\sigma' \not\models B$  since  $\sigma'(x) = 0$  and  $\sigma' \not\models x = 5$

### 3 4F-3 VCGen Mistakes

- 0 pts Correct

**Exercise 4F-4. Axiomatic Do-While [6 points].** Write a sound and complete Hoare rule for `do c while b`. This statement has the standard semantics (e.g.,  $c$  is executed at least once, before  $b$  is tested).

We know that

`do c while b = c; while do c`

Therefore, we can obtain the Hoare rule for `do c while b`

$$\frac{\vdash \{A\}c\{B\} \quad \vdash \{B \wedge b\}c\{B\}}{\vdash \{A\}\text{do } c \text{ while } b\{B \wedge \neg b\}}$$

#### 4 4F-4 Axiomatic Do-While

- 0 pts Correct