

14F-1 Bookkeeping

- 0 pts Correct

4F-2. VCgen for let.

It has a bug because of the ignorance of the original value of  $x$  in the let rule. Assume there is a fresh variable called  $tmp$  storing the original value of  $x$ . Then, if we use the  $tmp$  variable to express  $VC(\text{let } x=e \text{ in } c, B)$ , we have.

$$\begin{aligned} & VC(\text{let } x=e \text{ in } c, B) \\ &= VC(tmp=x; x=e; c; x=tmp, B) \\ &= VC(\neg tmp=x; VC(x=e; c; x=tmp, B)) \\ &= [x/tmp] VC(x=e; c; x=tmp, B) \\ &= [x/tmp] VC(x=e, VC(c; x=tmp, B)) \\ &= [x/tmp] [e/x] VC(c; x=tmp, B) \\ &= [x/tmp] [e/x] VC(c, VC(x=tmp, B)) \\ &= [x/tmp] [e/x] VC(c, [tmp/x] B) \end{aligned}$$

Therefore, the correct VCgen for let should be

$$[x/tmp] [e/x] VC(c, [tmp/x] B)$$

2 4F-2 VCGen for Let

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### 4F-3 VCien Mistakes.

We could show the bugginess of let rules by doing the following:

1. Set the command  $c$  as `let  $x=j$  in skip`, where  $j$  is an arbitrary number
2. Set the precondition  $B$  as  $x=j$ .
3. Set a state  $b$  as  $b(x)=i$ , where  $i$  is an arbitrary integer not equal to  $j$ .
4.  $b \models VC(c, B)$  because  $VC(c, B) = j$ .
5.  $\langle c, b \rangle \Downarrow b'$ , but  $b'(x) = i$  because we need to reassign back  $x$ 's original value.

6. Therefore,  $b' \not\models B$  because  $b'(x) = i$  but the post-condition  $B$  is  $x=j$ .

We could also provide another detailed example.

let  $x=y$   `$x=x+1$` .

$b(x) = 1$ ,  $b(y) = 1$ .

$B: x=2$

$VC(c, B) = VC(x=y+1, x=2) = \{y=1\}$ .

$b \models VC(c, B)$

while  $\langle c, b \rangle \Downarrow b'$   $b'(x)=1$   $b'(y)=1$   $b' \not\models B$ .

Because it doesn't recover  $x$  with the original value after the let rule, it is problematic.

### 3 4F-3 VCGen Mistakes

- 0 pts Correct

4F-4

Because  $c$  is executed at least once, before  $b$  is tested as in the 4F-4, we have

do  $c$  while  $b = c$ ; while  $b$  do  $c$  (And we could take into the  $\{B\}$ .)

Then, we have

$$\frac{\vdash \{A\} c \{B\} \quad \vdash \{B \wedge b\} c \{B\}}{\vdash \{A\} \text{do } c \text{ while } b \{B \wedge \neg b\}}$$

$A$  is arbitrary precondition, and  $B$  represents loop invariant.

~~AbT/p and q and~~

~~AbT/p~~

#### 4 4F-4 Axiomatic Do-While

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