

14F-1 Bookkeeping

- 0 pts Correct

Peer Review ID: 67844295 — enter this when you fill out your peer evaluation via gradescope

4F-2. A correct rule for let is given by

$$\text{VC}(\text{let } x = e \text{ in } c, B) = [x/y][e/x] \text{VC}(c, [y/x]B), \text{ where } y \text{ is fresh.}$$

4F-3. The unsoundness of the buggy let rule is apparent from the following:

1. (a command c) $\text{let } x = 1 \text{ in skip}$
2. (a post condition B) $x = 1$
3. (a state σ) $[x := 0]$
4. (that $\sigma \models \text{VC}(c, B)$) Using the buggy let rule and the rule for `skip`, we see that

$$\text{VC}(\text{let } x = 1 \text{ in skip}, x = 1) = [1/x]\text{VC}(\text{skip}, x = 1) = [1/x](x = 1) = (1 = 1) = \text{true},$$

so indeed

$$\sigma \models \text{VC}(\text{let } x = 1 \text{ in skip}, x = 1).$$

5. (that $\langle c, \sigma \rangle \Downarrow \sigma'$) Applying the operational semantics rule for `let` and `skip`, we have the following derivation tree

$$\frac{\overline{\langle 1, \sigma \rangle \Downarrow 1} \quad \overline{\langle \text{skip}, \sigma[x := 1] \rangle \Downarrow \sigma[x := 1]}}{\overline{\langle \text{let } x = 1 \text{ in skip}, \sigma \rangle \Downarrow \sigma[x := 1][x := \sigma(x)]}}.$$

Since $\sigma[x := 1][x := \sigma(x)] = \sigma$, this gives that

$$\langle \text{let } x = 1 \text{ in skip}, \sigma \rangle \Downarrow \sigma'$$

with $\sigma' = \sigma$.

6. (that $\sigma' \not\models B$) Since $\sigma'(x) = \sigma(x) = 0$, we see that $\langle x, \sigma' \rangle \Downarrow 0$ and $\langle 1, \sigma' \rangle \Downarrow 1$, thus $0 \neq 1$ implies $\sigma' \not\models (x = 1)$ as desired.

4F-4. We give the following sound and complete Hoare rule for `do c while b`:

$$\frac{\vdash \{A\} c \{B\} \quad \vdash \{B \wedge b\} c \{B\}}{\vdash \{A\} \text{do } c \text{ while } b \{B \wedge \neg b\}}.$$

2 4F-2 VCGen for Let

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3 4F-3 VCGen Mistakes

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4 4F-4 Axiomatic Do-While

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