Review Set 2

This Review Set asks you to prepare written answers to questions on context-free grammars. Each of the questions has a short answer. You may discuss this Review Set with other students and work on the problems together.

1. Let L_1 be the language consisting of all non-empty palindromes over the alphabet $\Sigma = \{a, b\}$. That is, L_1 consists of all sequences of a's and b's that read the same forward or backward. For example, $aba \in L_1$ and $abb \in L_1$ and $aabbbaa \in L_1$, but $abb \notin L_1$.

Let L_2 be the language over $\Sigma = \{a, b\}$ denotated by the regular expression a(a|b)*.

The language $L_3 = L_1 \cap L_2$ is context-free. A string s is in L_3 if $s \in L_1$ and $s \in L_2$. Write a context-free grammar for the language L_3 .

Optional Thing To Think About: Is the intersection of a context-free language and a regular language always context-free?

2. Consider the following grammar:

$$S \rightarrow aSb$$

$$S \rightarrow Sb$$

$$S \rightarrow \epsilon$$

- (a) Give a one-sentence description of the language generated by this grammar.
- (b) Show that this grammar is ambiguous by giving a single string that can be parsed in two different ways. Draw both parse trees.
- (c) Give an unambiguous grammar that accepts the same language as the grammar above.
- 3. Using the context-free grammar for Cool given in the Cool Reference Manual, draw a parse tree for the following expression.

while not
$$(x <- z <- 0)$$
 loop $y <- z + 2 * x + 1$ pool

Note that the context-free grammar by itself is ambiguous, so you will need to refer to the precedence and associativity rules to get the correct tree.

4. Give an example of a grammar that is LL(3) but not LL(2).