# Optimistic Recovery for High-Availability Software via Partial Process State Preservation

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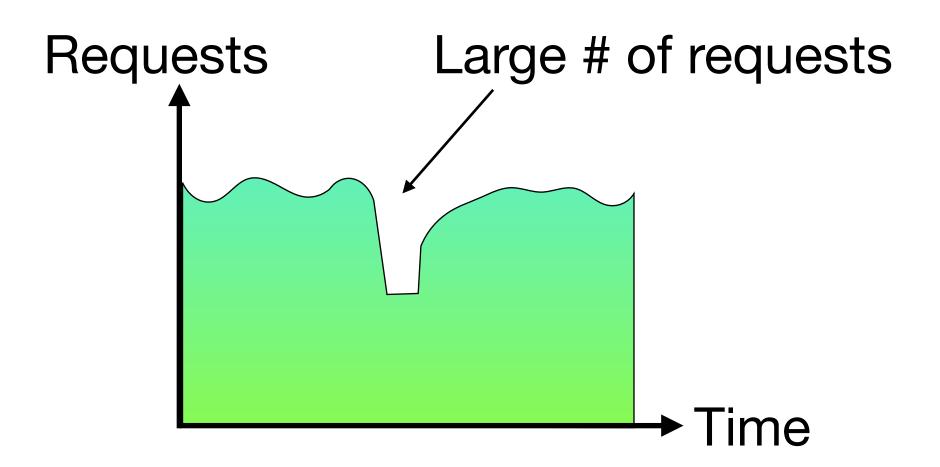
University of Michigan



## Availability Is Crucial

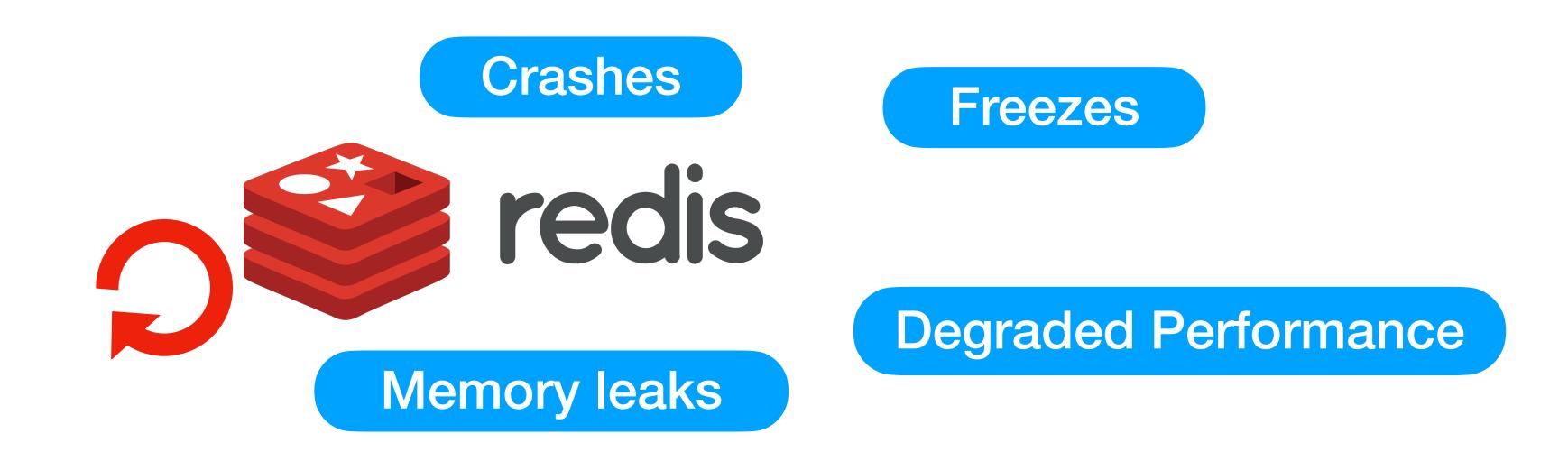


High SLO Requirement



Brief downtime but high impact

#### Software Faults Are Inevitable

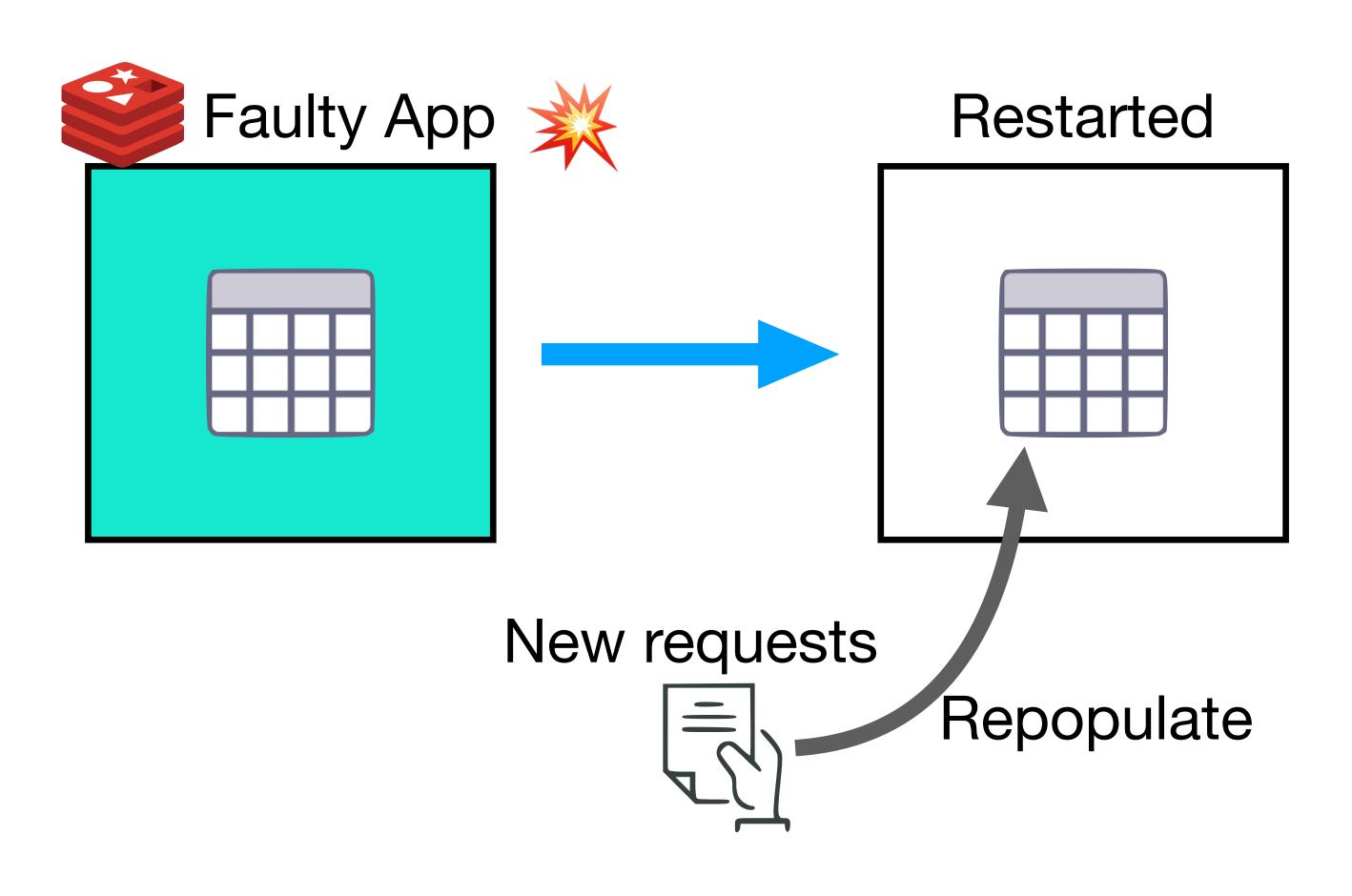


#### Recovery requirements:

- 1. High availability: must recover both quickly and correctly
- 2. Meaningful availability: service not only being up but also high performance

### Recovery In Practice

#### **Practice 1: Restart with Empty State**



#### Pros:

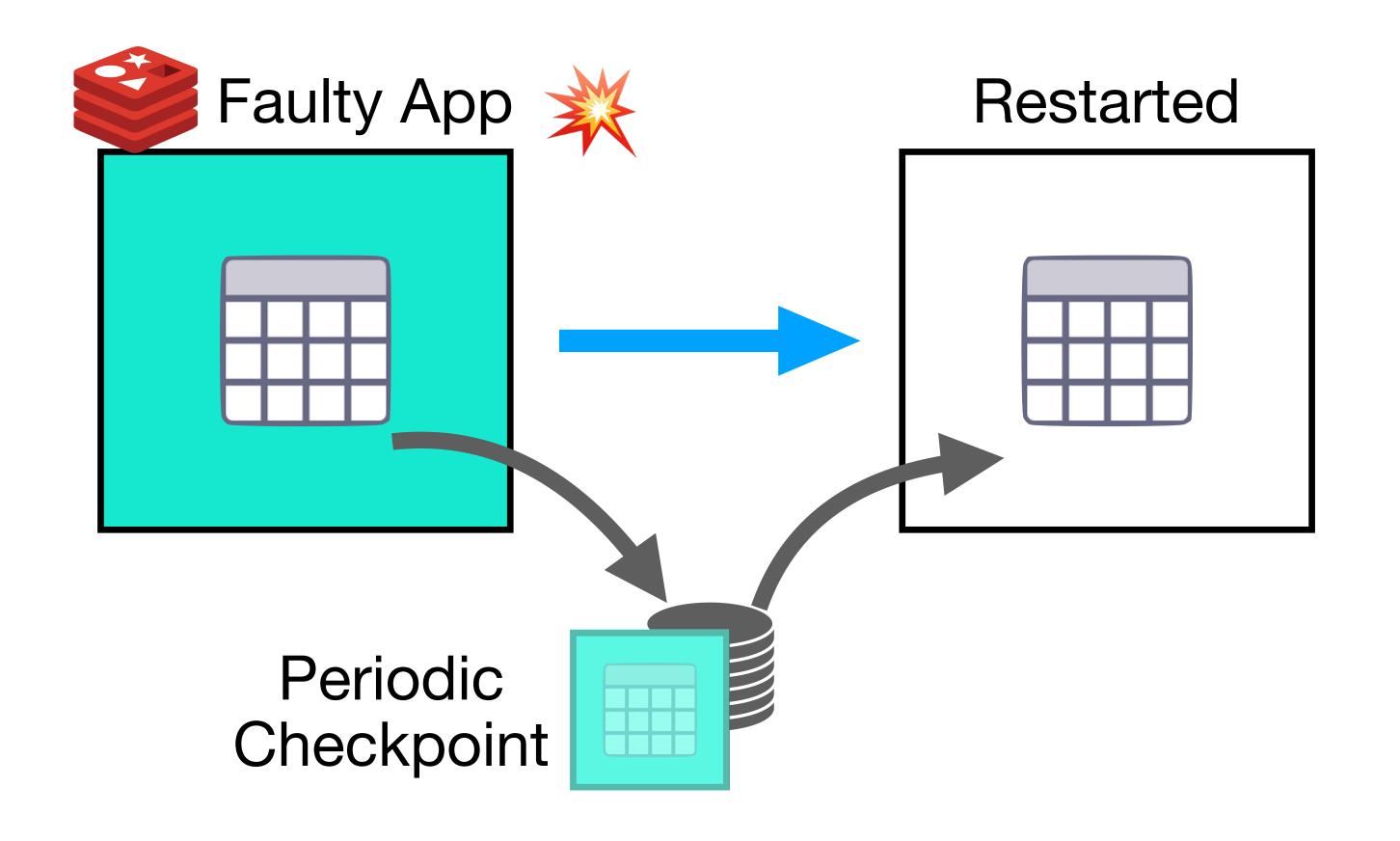
- Simplicity
- Eliminate bad state

#### Cons:

Poor availability post-restart

### Recovery In Practice

#### Practice 2: Process Checkpointing



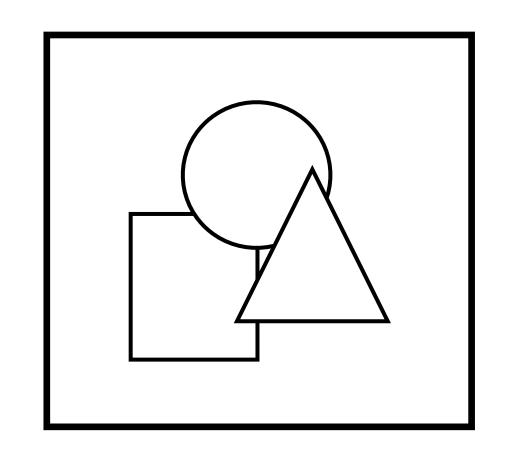
#### **Pros:**

 Better post-restart performance

#### Cons:

Risks persisting faulty state

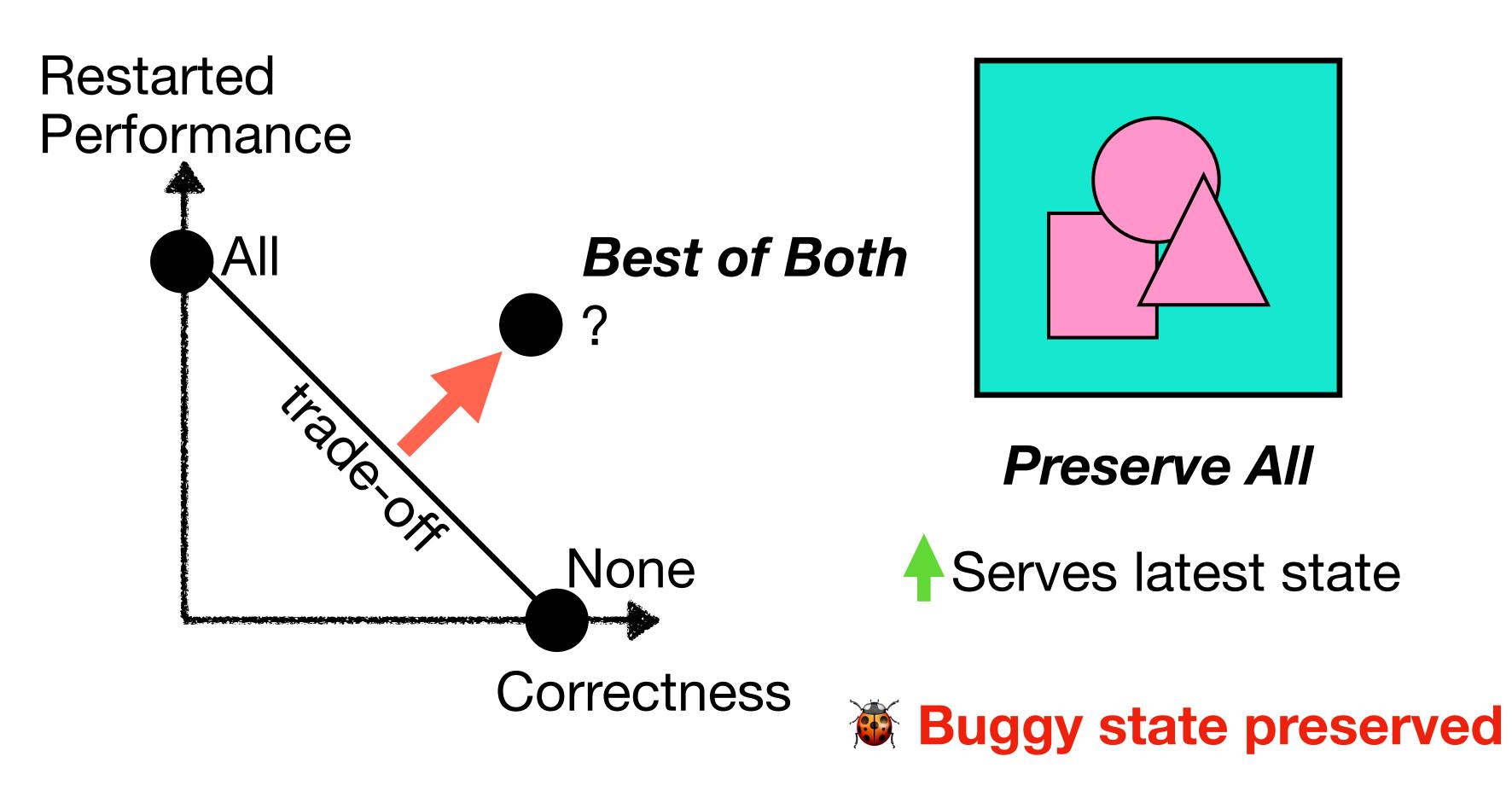
#### Dilemma



Preserve None



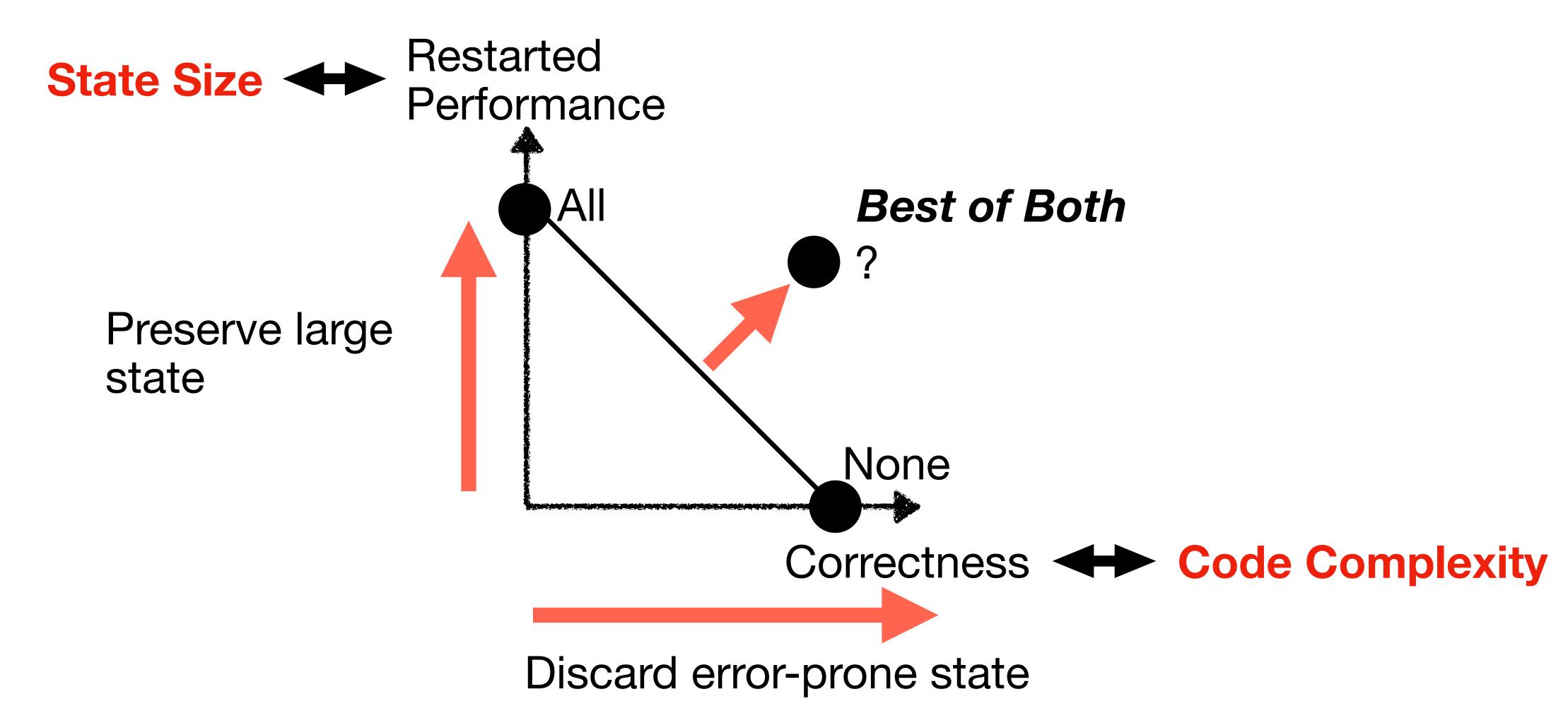
Poor availability during warmup



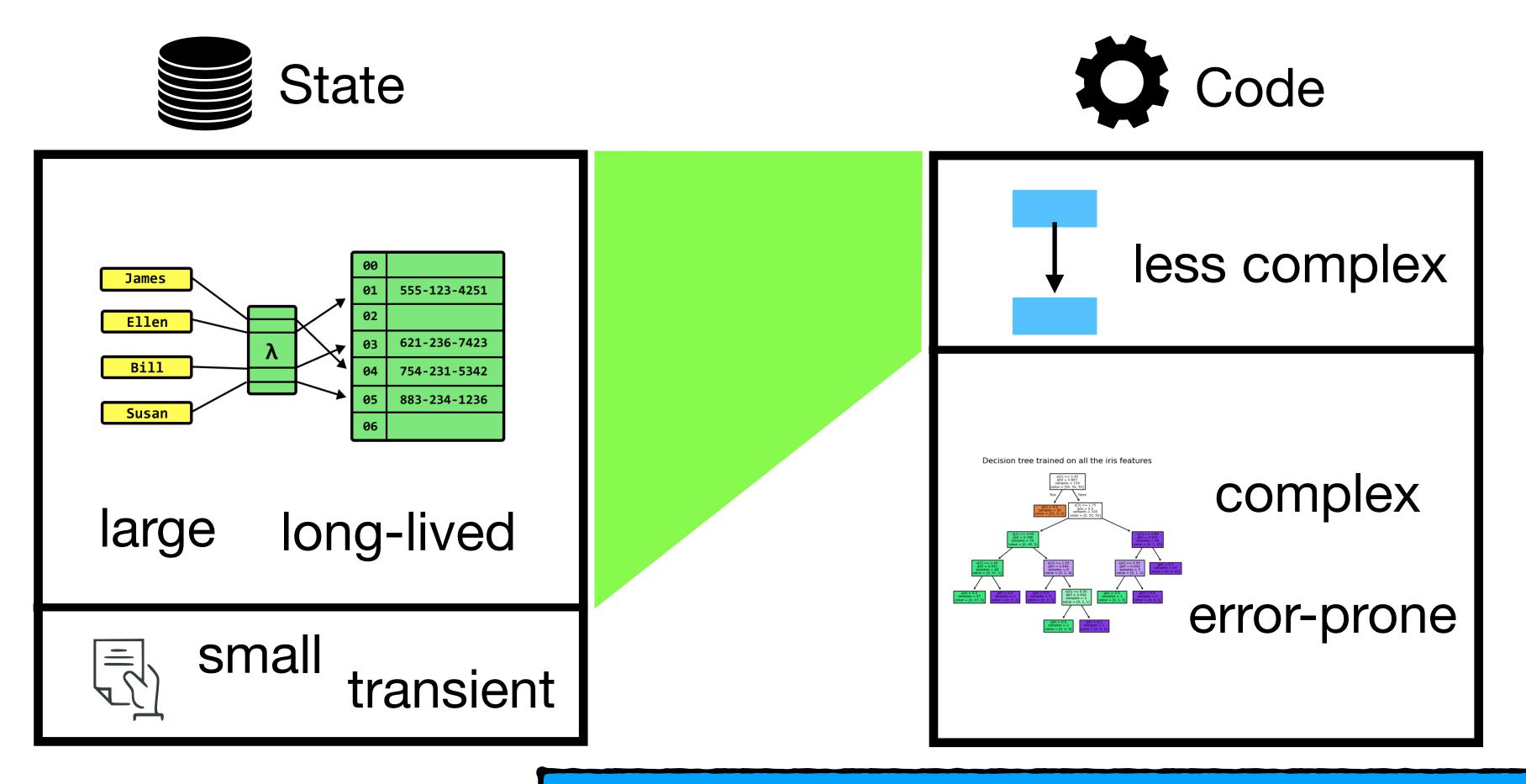
## Challenges

- How to determine what state to preserve?
  - Partial Process State Preservation
- How to get high availability without compromising correctness?
  - Optimistic Recovery

## What state to preserve?

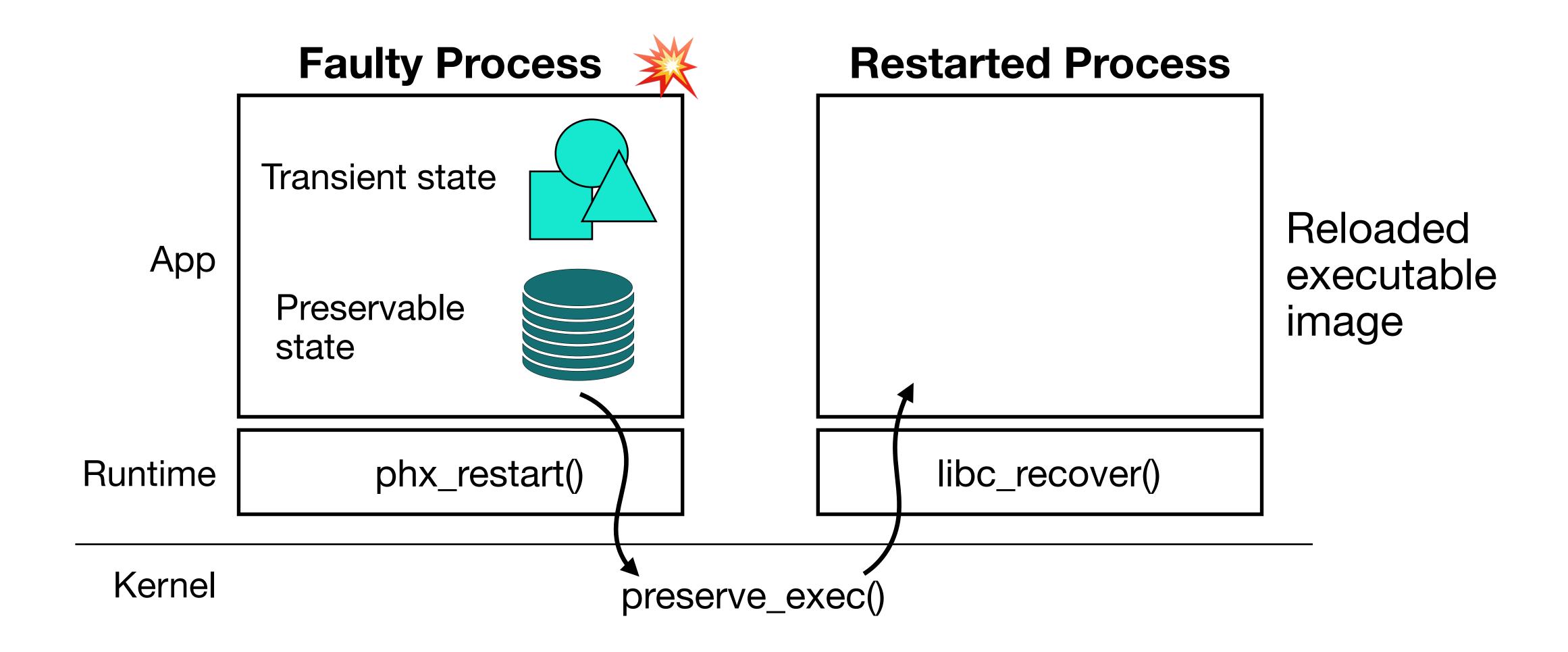


## Insight: State Size & Complexity Imbalance

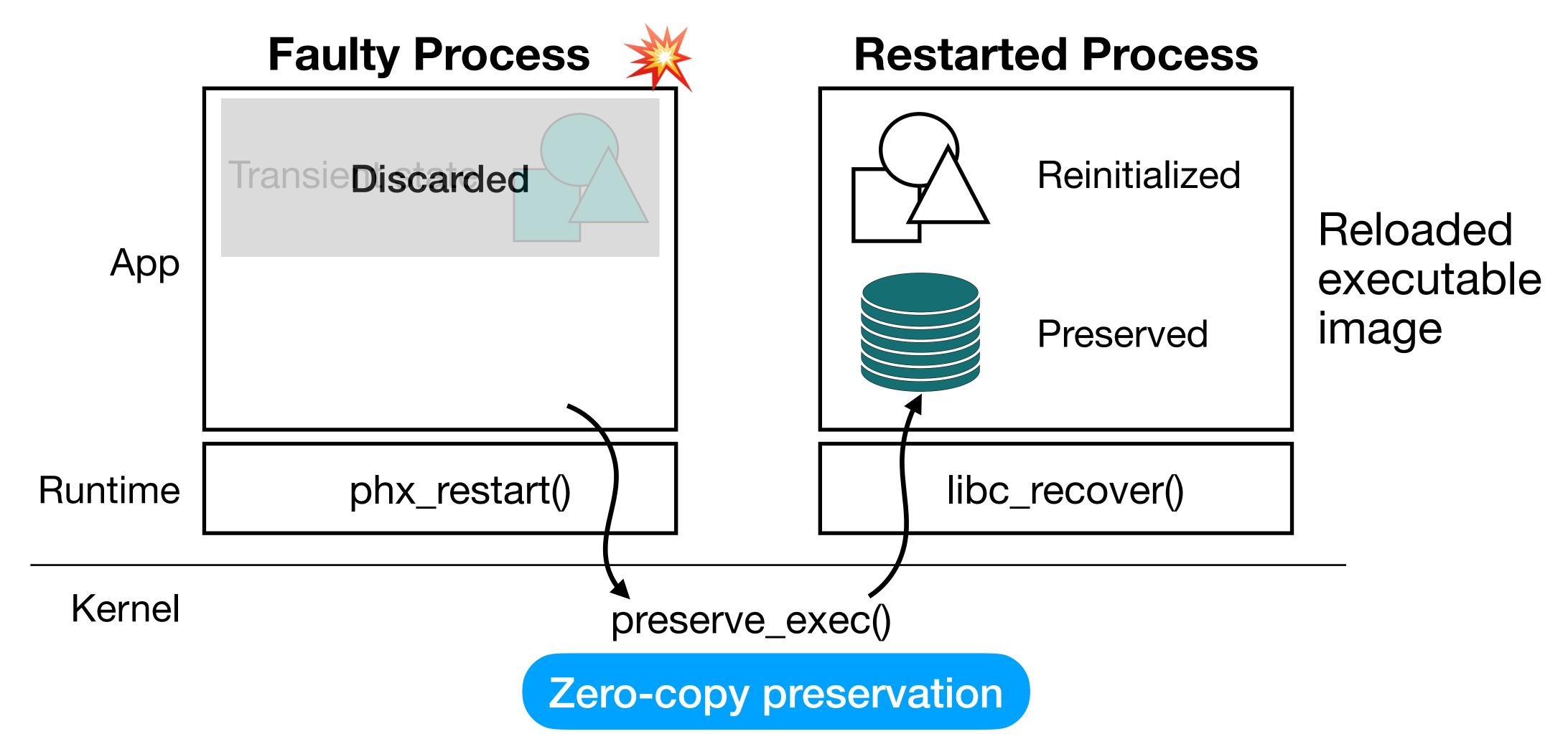


Opportunity: preserve large state with low chance of incorrectness

### **Phoenix: Partial State Preservation**



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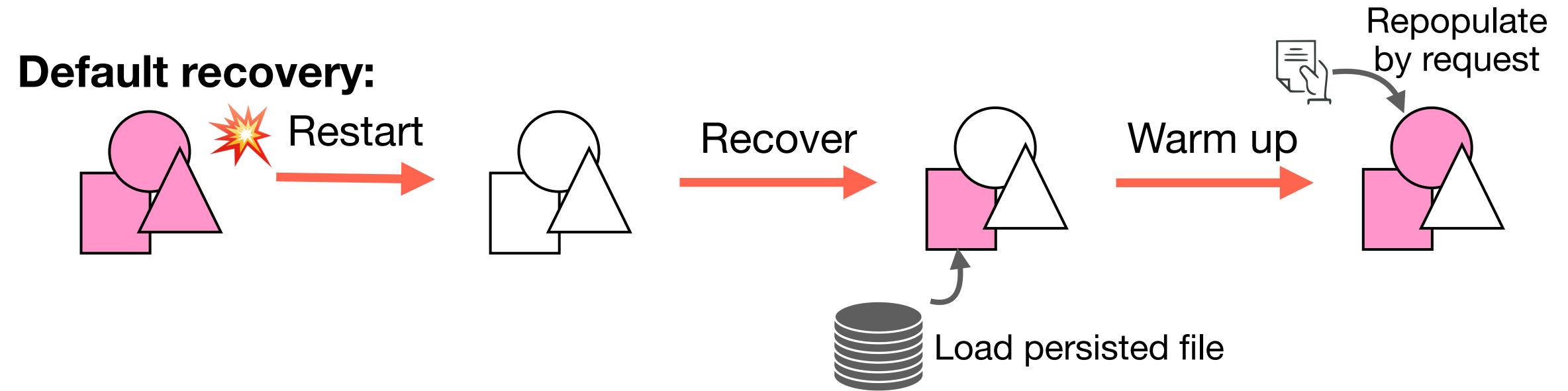
#### Phoenix

Partial Process State Preservation

- Optimistic Recovery
  - Unsafe Region: efficient consistency check
  - Cross-Check Validation: long-term assurance

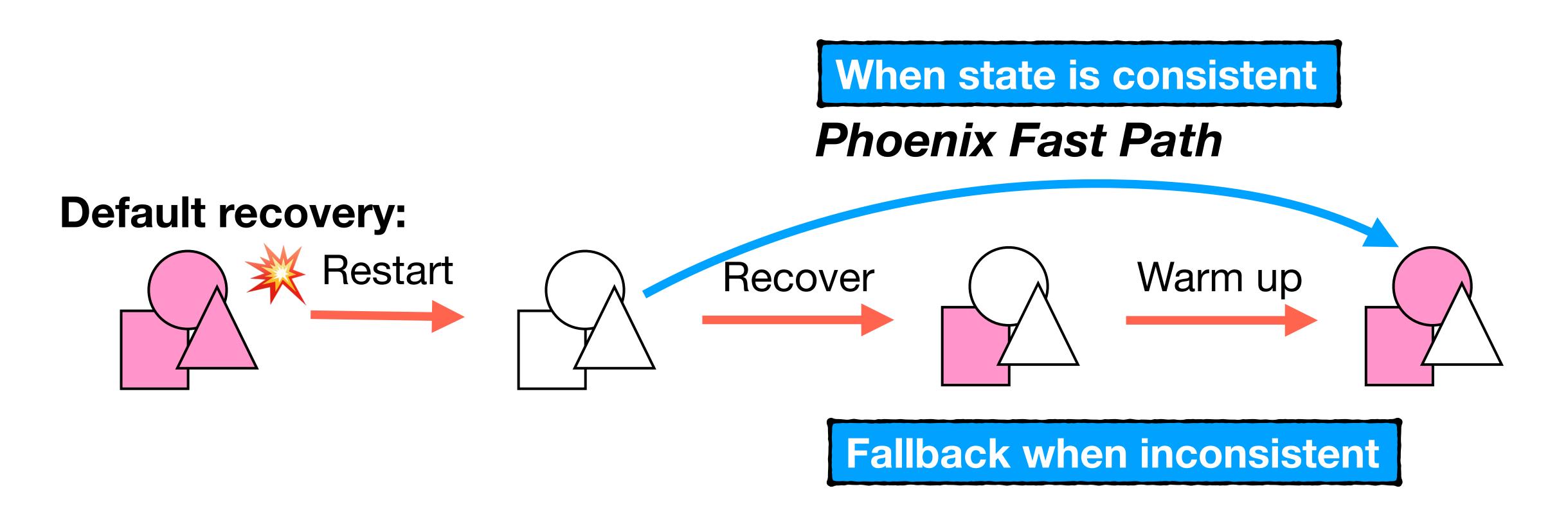
## **Optimistic Recovery**

Designed to provide high chance of quick recovery



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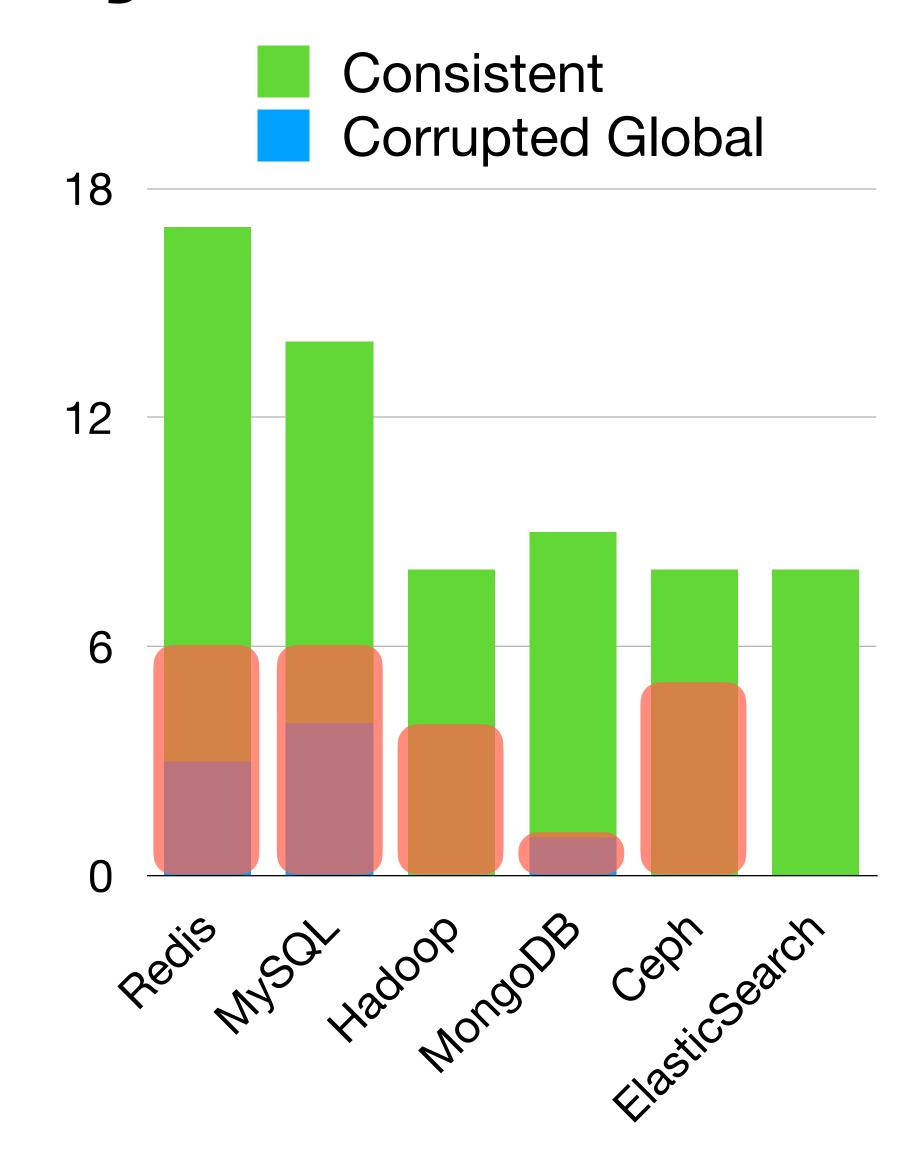
## How to detect inconsistency?

Insight from Real-World Bug Study

Collect 64 case from 6 systems

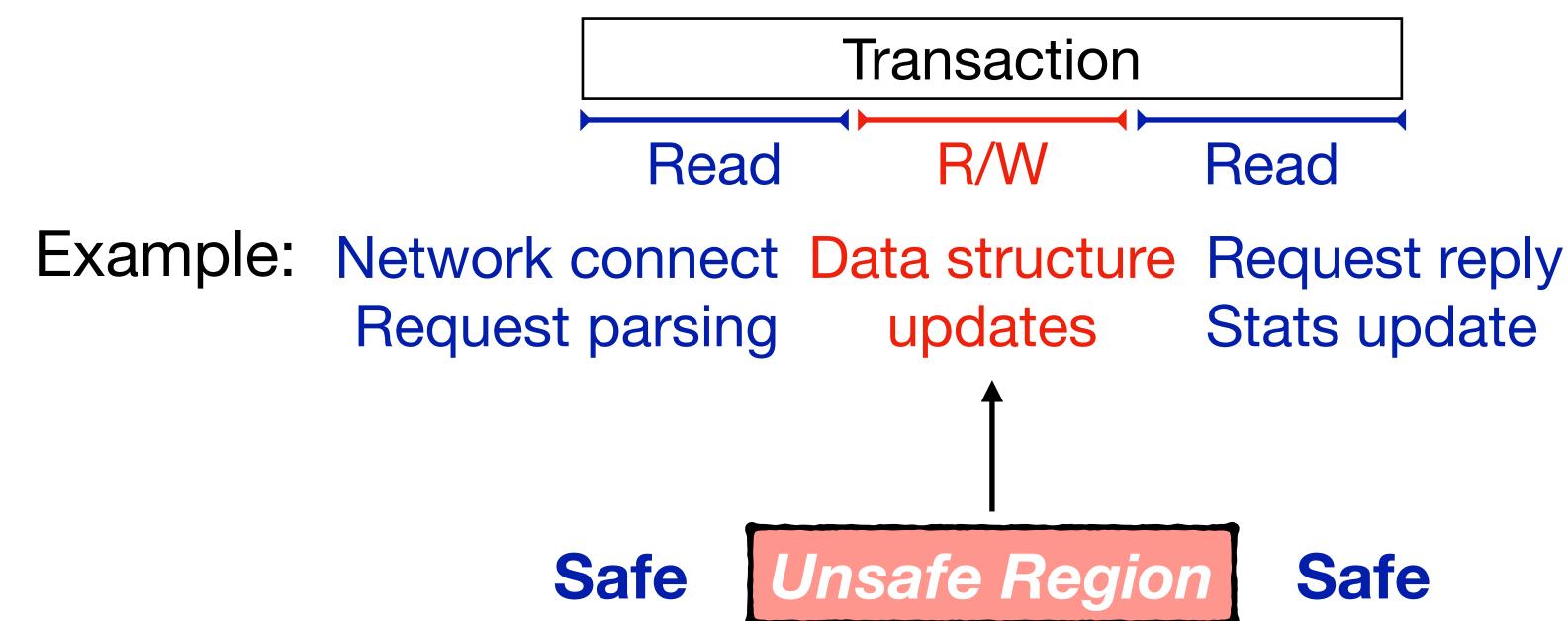
- 56 transient-only
  - Discarding transient state already give correctness
- 8 corruption in long-lived state
  - High-correlation: all happen during modification operations

Solution: fallback on inconsistent write



# Unsafe Region

Write portion of preservable state in one transaction

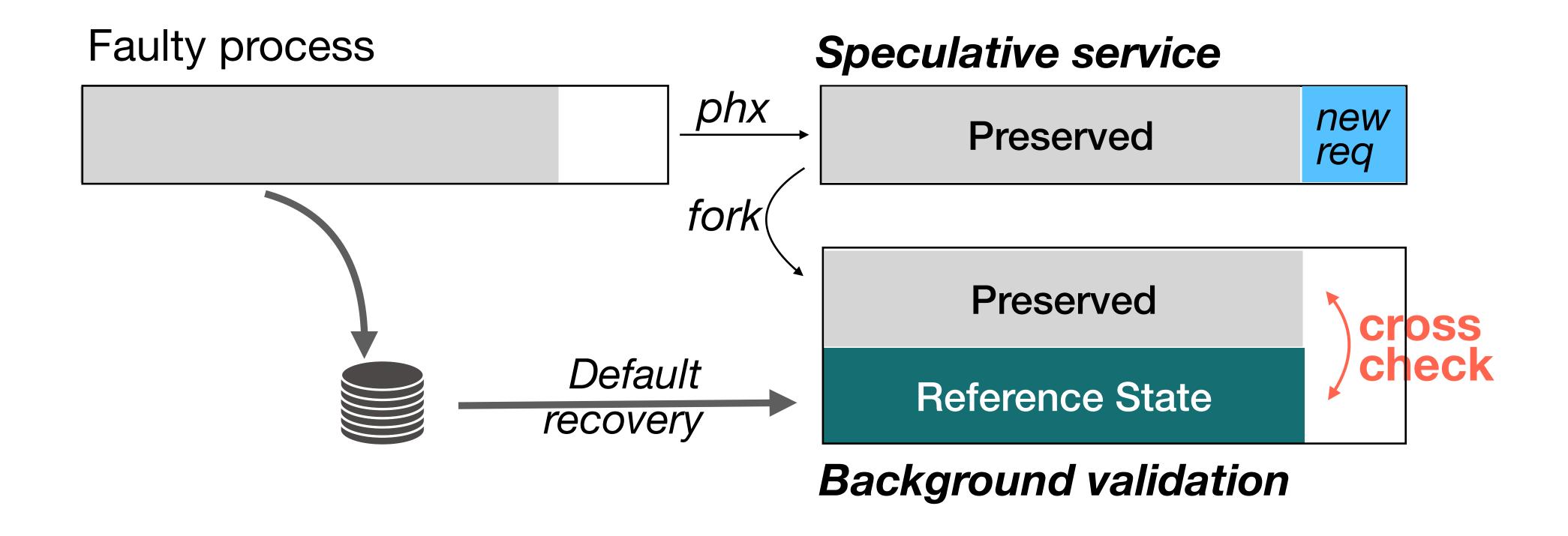


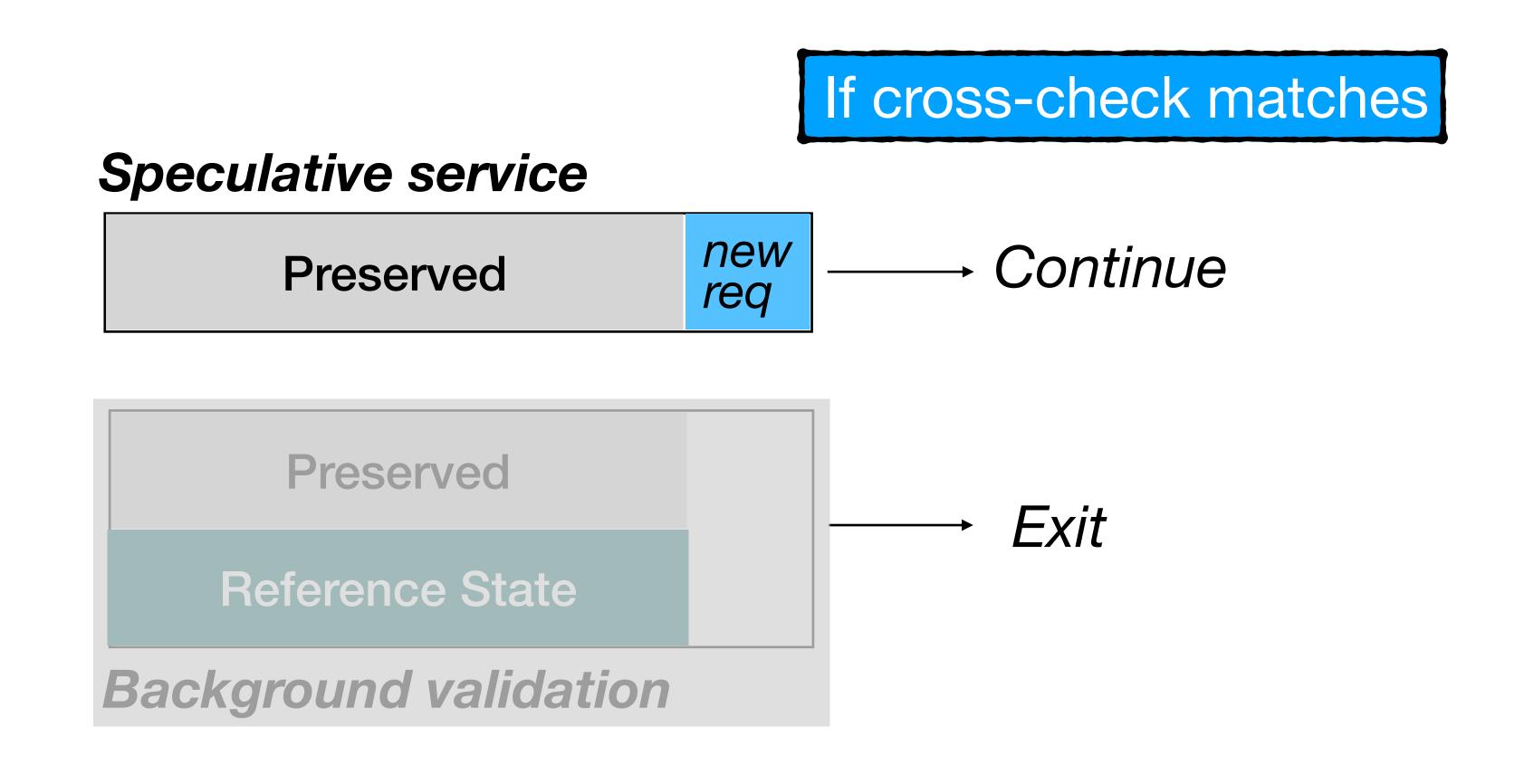
Either manual marking or using Phoenix compiler auto-instrumentation

## Unsafe Region

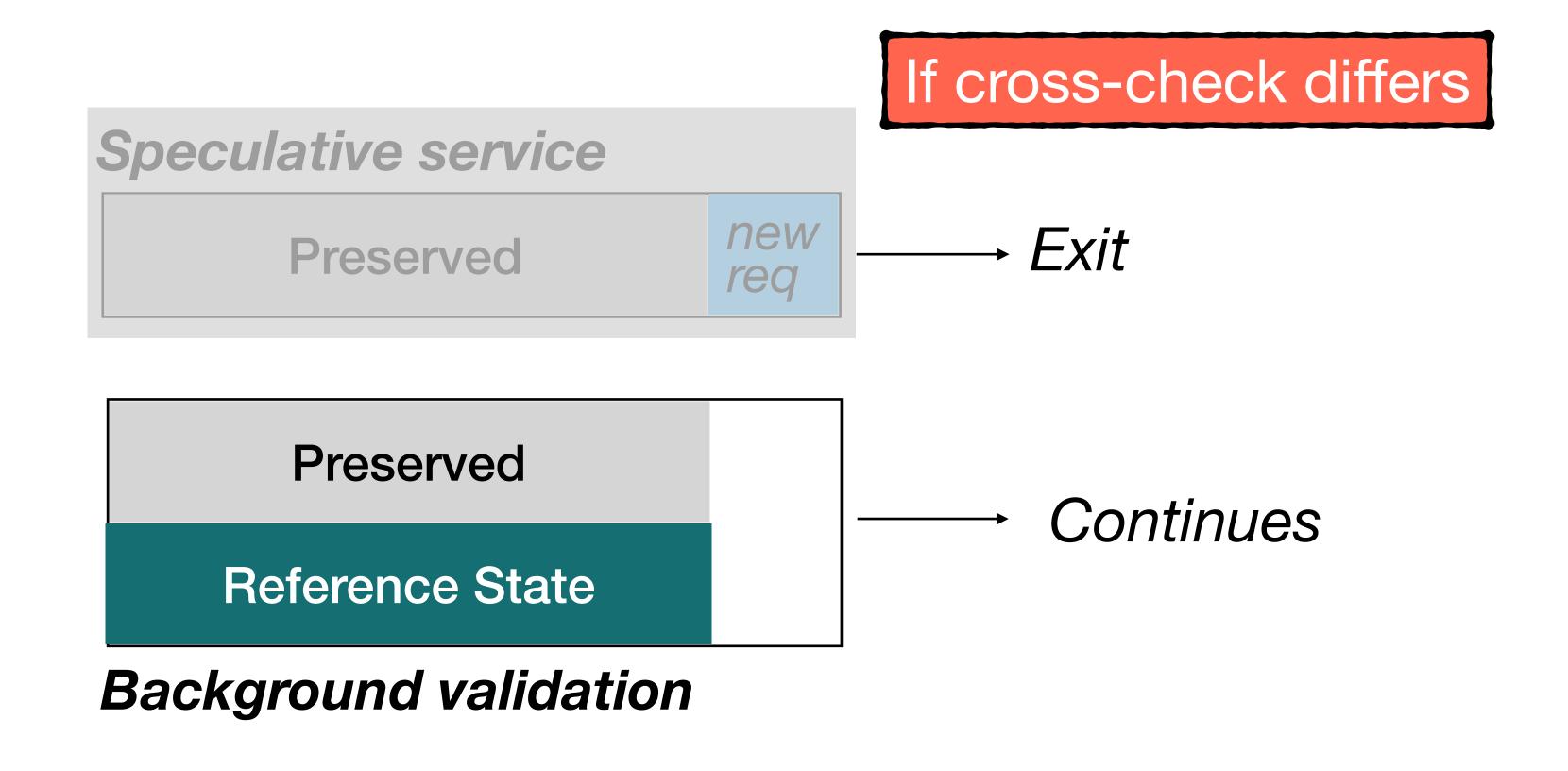
Safe and effective indication of inconsistency in experiment

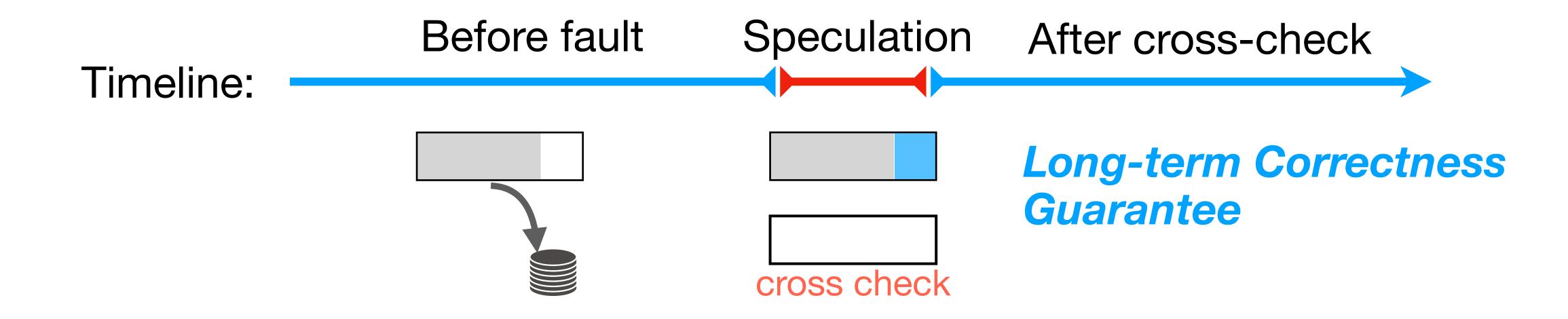
- However, indirect indication is not assurance
  - Background Cross-Check Validation





Availability gain from Phoenix restart!





Potential inconsistency bounded to cross-check window

#### Evaluation

- Implemented Phoenix in Linux kernel 5.15
- Ported 6 applications without knowing the bugs
- We try to answer questions:
  - Effectiveness: How much availability improvement does Phoenix achieve?
  - Correctness: Does Phoenix recover correctly?
  - Effort: What effort and overhead are required to employ Phoenix?

# Ported Applications

App	Preserved States	Codebase LoC	Changes
Redis	In-mem KV hash table	140,996	348 (0.25%)
LevelDB	Skiplist memory tables	21,192	312 (1.50%)
Varnish	Web page cache objects	109,564	281 (0.26%)
Squid	Web page cache objects	186,219	190 (0.10%)
XGBoost	Gradients and model	38,906	158 (0.41%)
VPIC	Particles and physical fields	44,773	272 (0.61%)

Only moderate usage efforts

## Availability

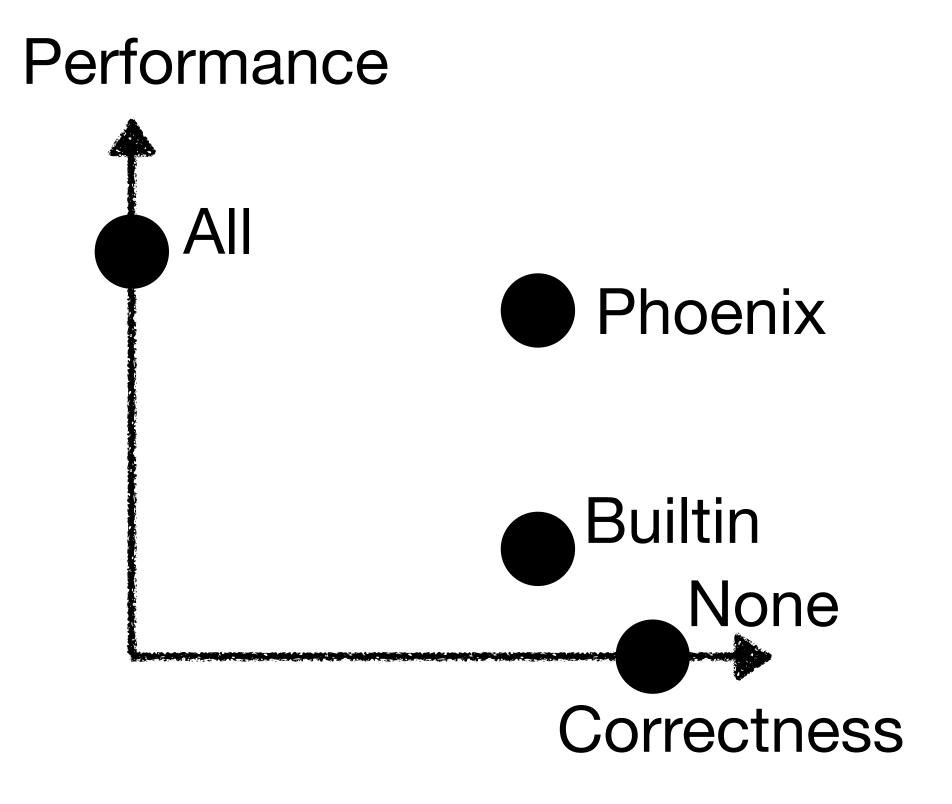
Reproduce 17 real-world bugs from various categories.

Compare between 4 setups

- None: Original application without persistence
- All: Whole process checkpointing (CRIU)
- Builtin: Application default recovery method
- Phoenix: Phoenix with Builtin as fallback

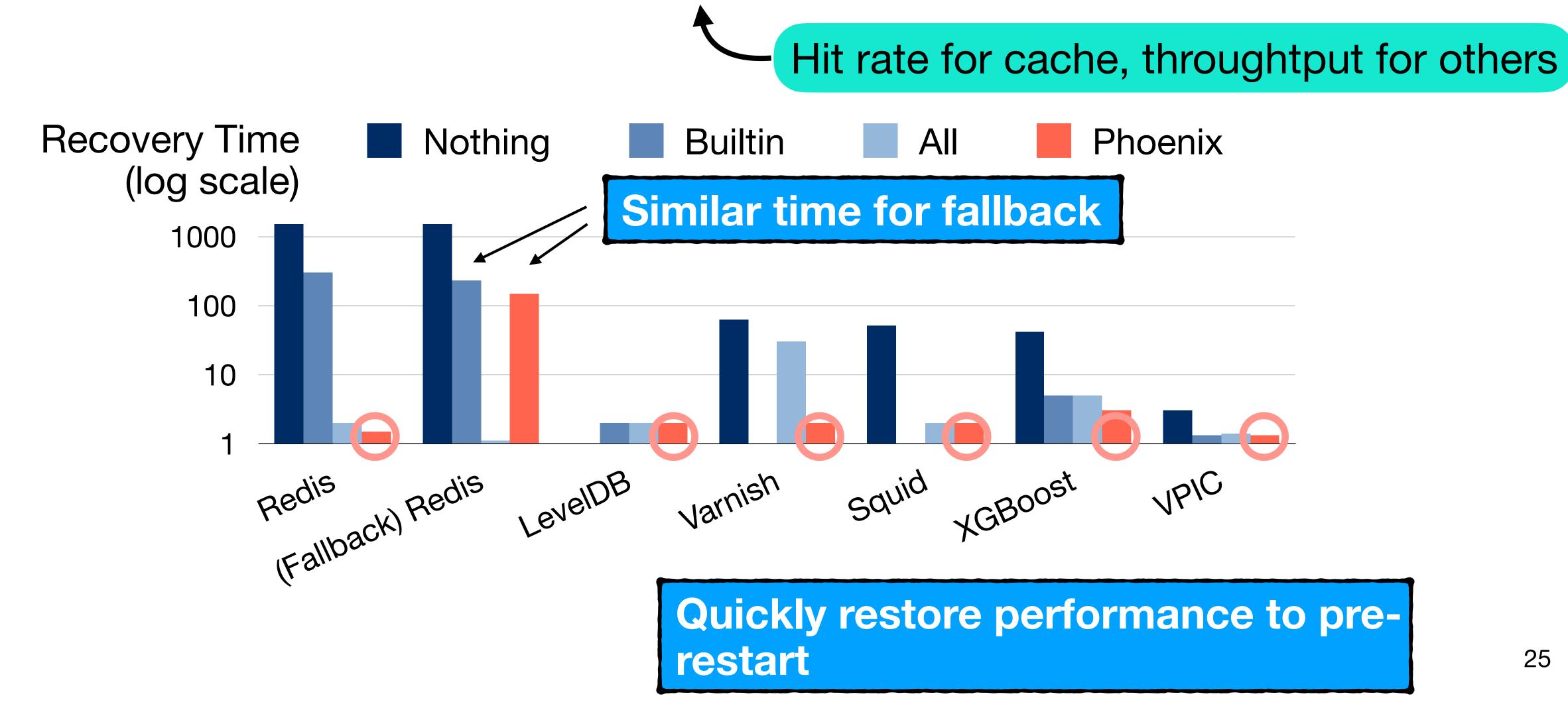
16 recovered by Phoenix

1 fallback to default recovery



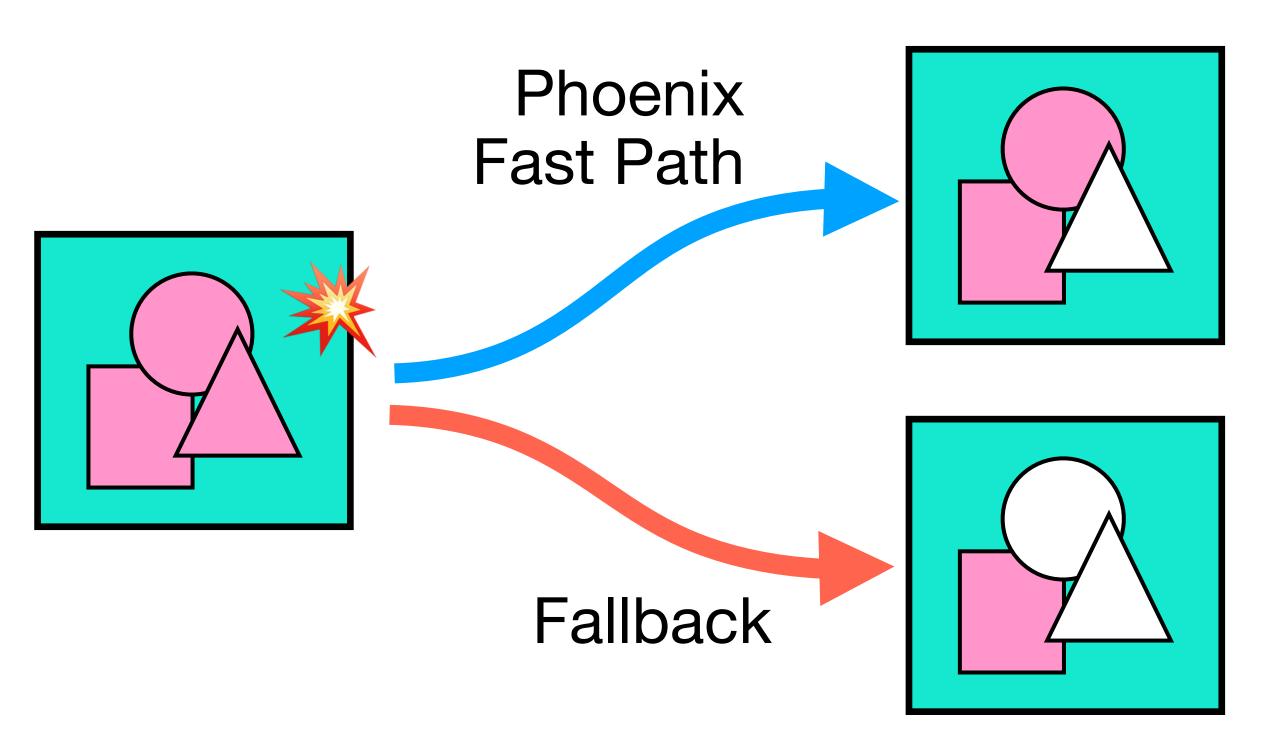
## Availability

#### Time to Recover to 90% of Effective Availability



# Correctness: Large Scale Injection

#### Only Unsafe Region



Total: 4,000 runs

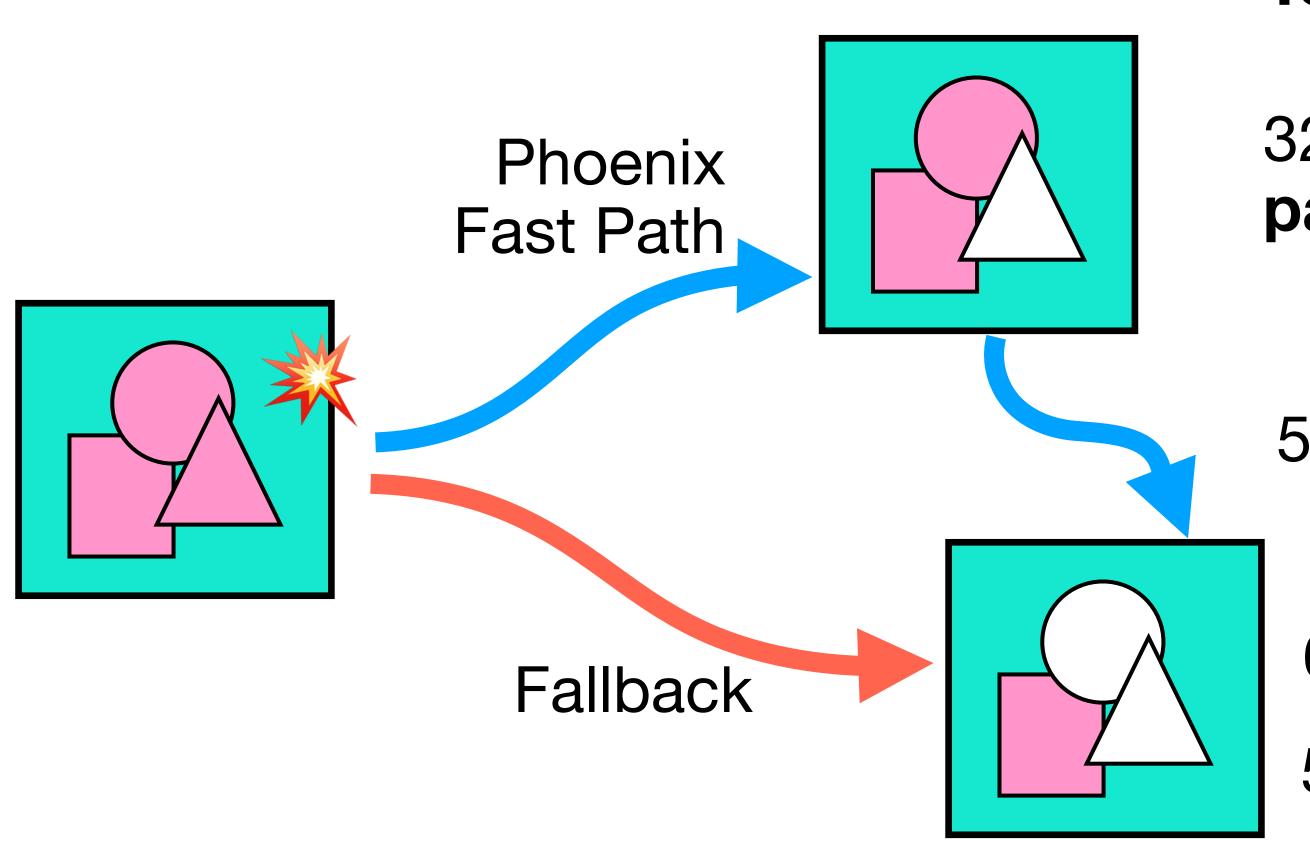
3,287 successful Phoenix recovery

82% fast path taken

662 unsafe region caught 51 faulty restart caught

# Correctness: Large Scale Injection

**Cross-Check Mode** 



Total: 400 runs

325 successful Phoenix recovery passing cross-check

5 cross-check early fallbacks

65 unsafe region caught

5 faulty restart caught

#### Conclusion

- Partial state preservation
  - Effectively preserves most state while discarding most faults
- Optimistic recovery
  - High availability with long-term correctness guarantee
- New OS-level mechanism *Phoenix* 
  - Significant availability and performance improvements

