IronFleet: Proving Practical Distributed Systems Correct

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Presenter: Armin Vakil



Abstract

The Paxos algorithm, when presented in plain English, is very simple.













Examples of bug disruptions



Examples of bug disr **Facebook Explains October 4**

Outage

October 8, 2021 6:13 PM EDT Last Updated a month ago Technology

Facebook apologizes for second outage in a week, services back up

2 minute read

By Subrat Patnaik and Sheila Dang



Amazon: Here's what caused the major AWS outage last week

AWS explains how adding a small amount of capacity to Kinesis servers knocked out dozens of services for hours.



By Liam Tung | November 30, 2020 | Topic: Cloud

TECH • AMAZON WEB SERVICES

BY CHRIS MORRIS October 12, 2021 12:27 PM EDT

Twitter's massive outage may be over, company says 'no evidence' of hack

The cause is unclear

By Nick Statt | @nickstatt | Oct 15, 2020, 5:58pm EDT





Amazon Web Services dashboard goes temporarily offline







Verification to the Rescue



Verification to the Rescue

• Automated proof for correctness



Verification to the Rescue

• Automated proof for correctness







• First mechanically-checked proof of implementation



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- First mechanically-checked proof of implementation
- IronRSL Replicated state library







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- IronRSL Replicated state library
- IronKV- Sharded Key-Value Store







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- IronRSL Replicated state library
- IronKV- Sharded Key-Value Store
- Liveness proof along with safety







- First mechanically-checked proof of implementation
- IronRSL Replicated state library
- IronKV- Sharded Key-Value Store
- Liveness proof along with safety
- Reasonable proof-to-code ratio







Outline

- Safety
- Parallelism
- Liveness
- Evaluation
- Conclusion



Implementation







Specification

Implementation







Specification

Implementation















K	V
_	_
_	_
	_





















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Sharded Hash Table







Sharded Hash Table











































Host 1			
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Base



Base





Base

V





V

Base

V'





V

Base

 \mathbf{V}'









function Abstraction(L:Variables) : H.Variables







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lemma RefinementInit(v:Variables) requires **Init**(v) ensures **Safe**(v) // Safe base case





function Abstraction(L:Variables) : H.Variables predicate Safe(v:Variables)

lemma RefinementInit(v:Variables) requires **Init**(v) ensures **Safe**(v) // Safe base case ensures MapSpec.Init(Abstraction(v)) // Refinement base cas











lemma RefinementNext(v:Variables, v':Variables)



lemma RefinementNext(v:Variables, v':Variables) requires Next(v, v')



lemma RefinementNext(v:Variables, v':Variables) requires Next(v, v') requires **Safe**(v)



lemma RefinementNext(v:Variables, v':Variables)

requires **Next**(v, v')

requires **Safe**(v)

ensures **Safe**(v') // Safe inductive step



lemma RefinementNext(v:Variables, v':Variables)

requires Next(v, v')

requires **Safe**(v)

- ensures **Safe**(v') // Safe inductive step
- ensures MapSpec.Next(Abstraction(v), Abstraction(v'))



lemma RefinementNext(v:Variables, v':Variables)

requires **Next**(v, v')

requires **Safe**(v)

ensures **Safe**(v') // Safe inductive step

|| Abstraction(v) == Abstraction(v')




Subtleties of distributed protocols





Subtleties of distributed protocols

Implementation difficulties







Subtleties of distributed protocols

Implementation difficulties



Maintaining safety invariants





Subtleties of distributed protocols

Concurrent Hosts

Implementation difficulties





Maintaining safety invariants



Subtleties of distributed protocols

Concurrent Hosts

Implementation difficulties







Subtleties of distributed protocols

Concurrent Hosts

Implementation difficulties



Liveness

Memory management



Subtleties of distributed protocols

Concurrent Hosts

Implementation difficulties

Efficient Data structures

Maintaining safety invariants

Liveness

Memory management



Subtleties of distributed protocols

Concurrent Hosts

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Liveness

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Integer overflow



Subtleties of distributed protocols

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Host A





Host A





Host A





















Enforce all receives precede sends













Enforce all receives precede sends













Enforce all receives precede sends





















Host A





Host A





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Main Event Handler

```
method Main()
{
  var s:ImplState;
  s := ImplInit();
  while (true) {
    s := EventHandler(s);
  }
}
```


Main Event Handler

```
method Main()
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 var s:ImplState;
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•
```

• Event Handler runs infinitely often





Main Event Handler

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method Main()
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}
```

• Event Handler runs infinitely often







Main Event Handler

```
method Main()
 var s:ImplState;
 s := ImplInit();
 while (true) {
   s := EventHandler(s);
```

- Event Handler runs infinitely often
- Each action runs infinitely often, too





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Safety

	20,000				
	18,000				
	16,000				
Lines of	14,000				
Code	12,000				
Couc	10,000				
	8,000				
	6,000			_	
	4,000			_	
	2,000			_	
	0				
		Spec	Impl	Proof	
		Common Libraries			







Safety































IronRSL Performance

Maximum throughput (RPS)

	IronRSL		
45000			
40000			
35000			
30000			
25000			
20000			
15000			
10000			
5000			
0			

With Batching

SL Baseline (EPaxos)



Without Batching



IronRSL Performance

Maximum throughput (RPS)

	IronKS		
15000	 		
10000	 		
35000	 		
30000	 2.4x		
25000			
20000	 		
15000			
0000			
5000			
0			

With Batching

RSL Baseline (EPaxos)



Without Batching



IronKV Performance

IronKV 40 30 Peak throughput (kRPS) 20 10 0 128B 1KB 8KB Get





IronKV Performance

IronKV 40 **1.5**x 30 Peak throughput (kRPS) 20 10 0 128B 1KB 8KB

Get





IronKV Performance

IronKV 40 1.5x 30 Peak throughput (kRPS) 20 10 0 128B 1KB 8KB

Get



Set



Conclusion

- First mechanically-checked proof of implementation
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- Reasonable proof-to-code ratio
- Comparable performance with state-of-the-art



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Thanks,

