EECS 583 – Automatic Parallelization Via Decoupled Software Pipelining

University of Michigan

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Parallelization: Scientific vs Non-Scientific Codes

Scientific Codes (FORTRAN-like)

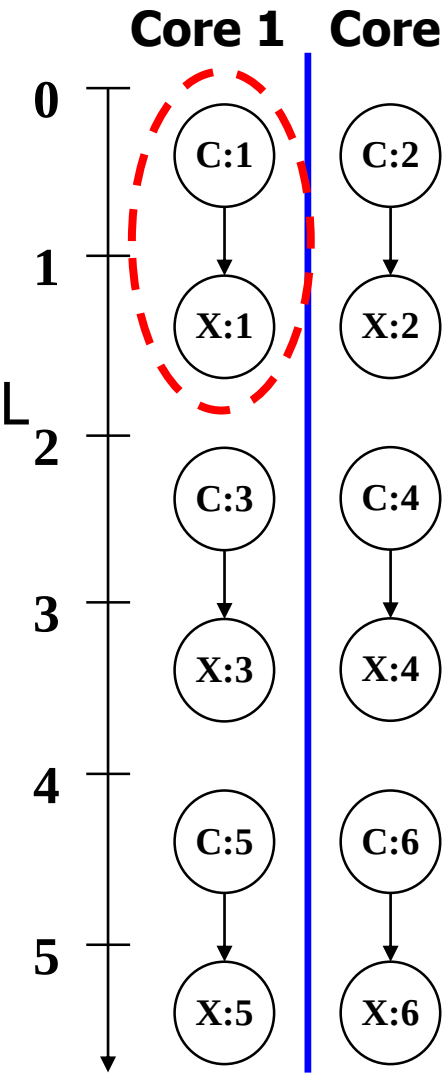
```
for(i=1; i<=N; i++) // C
  a[i] = a[i] + 1;  // X
```

General-purpose Codes (legacy C/C++)

```
while(ptr = ptr->next)    // LD
  ptr->val = ptr->val + 1; // X
```

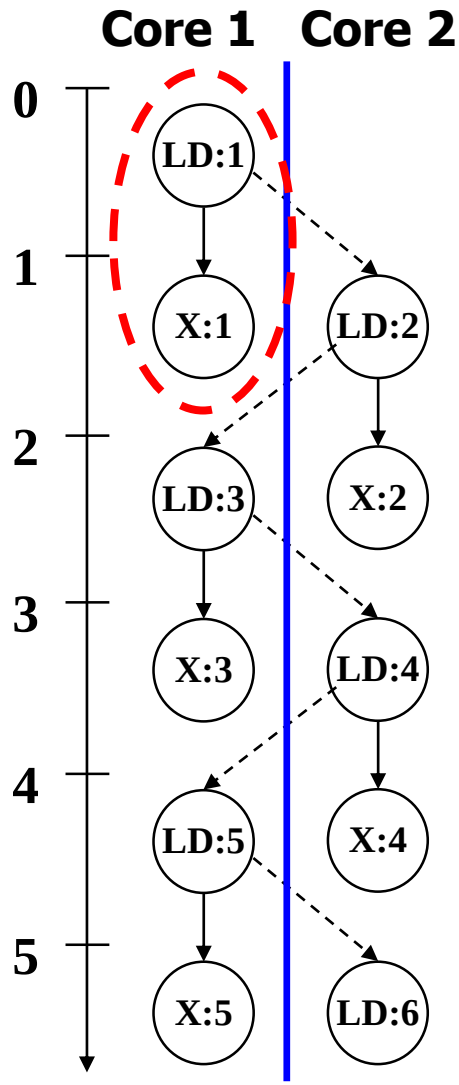
Independent Multithreading (IMT)

Example: DOALL parallelization



Cyclic Multithreading (CMT)

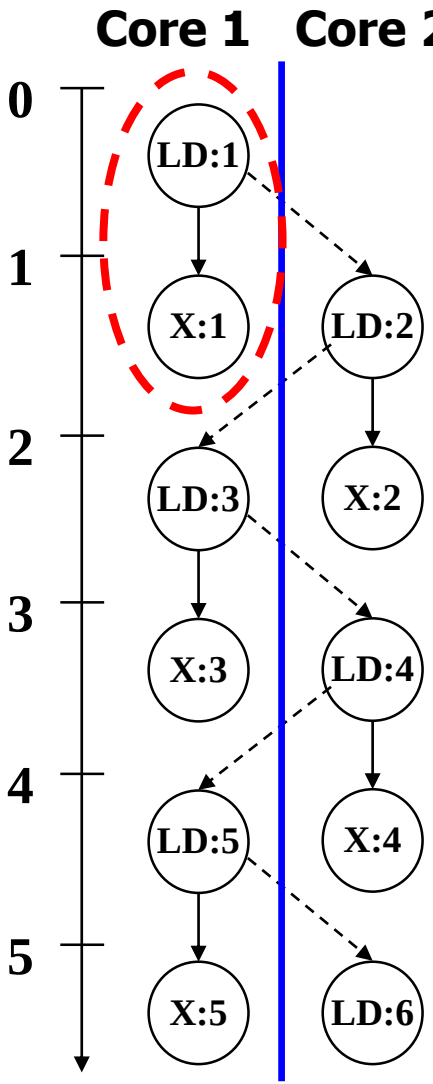
Example: DOACROSS [Cytron, ICPP 86]



Alternative Parallelization Approaches

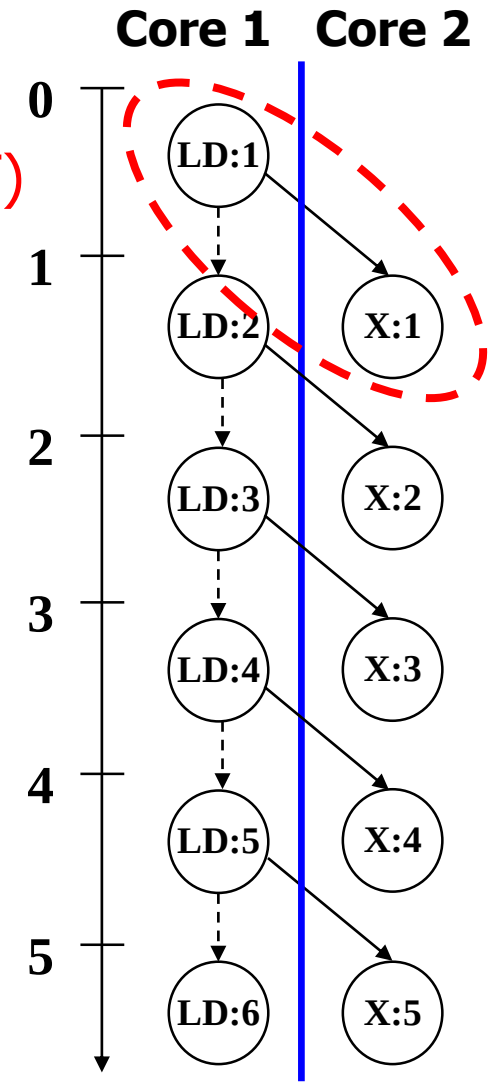
```
while(ptr = ptr->next)     // LD
  ptr->val = ptr->val + 1; // X
```

Cyclic Multithreading (CMT)

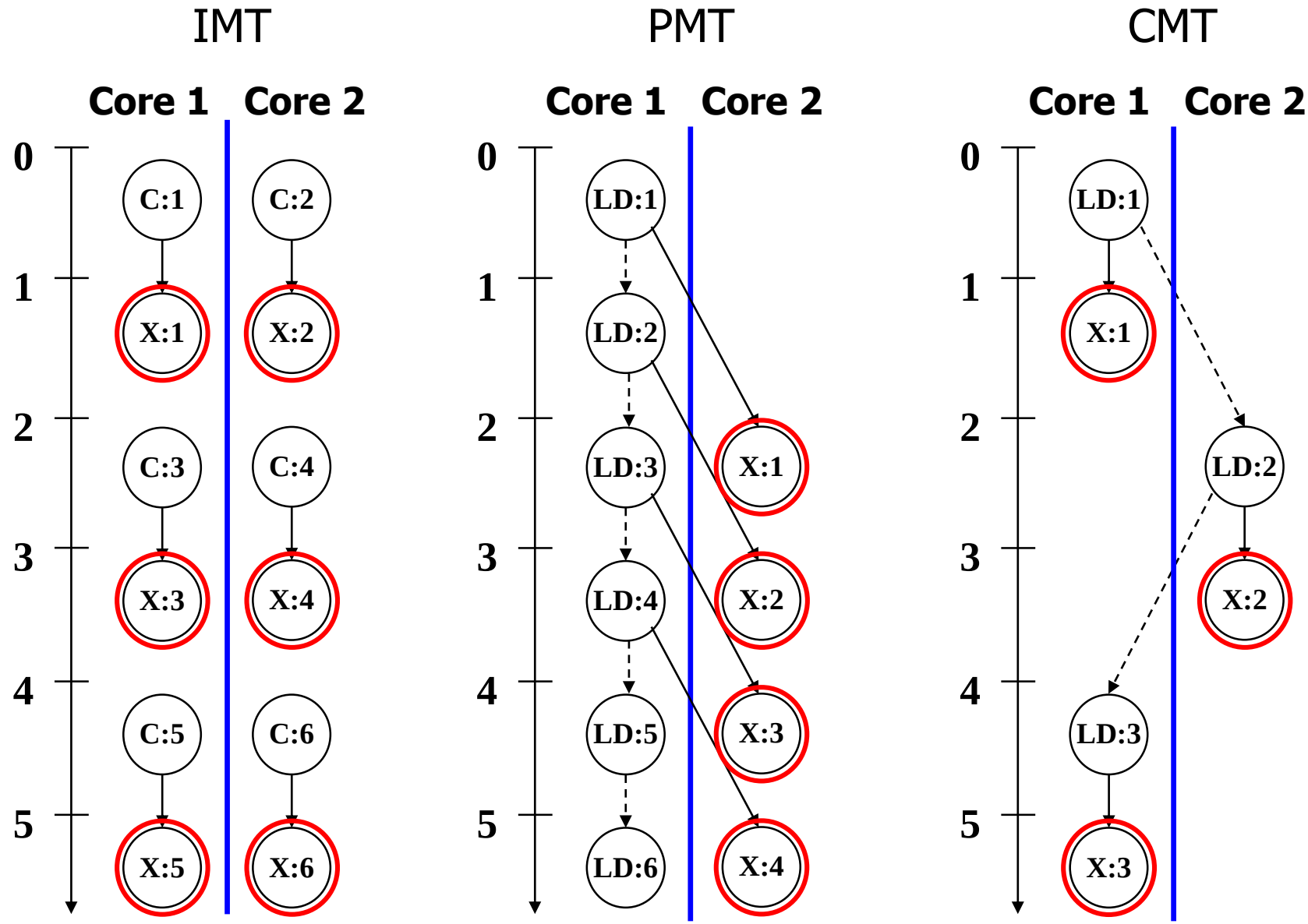


Pipelined Multithreading (PMT)

Example: DSWP **[PACT 2004]**



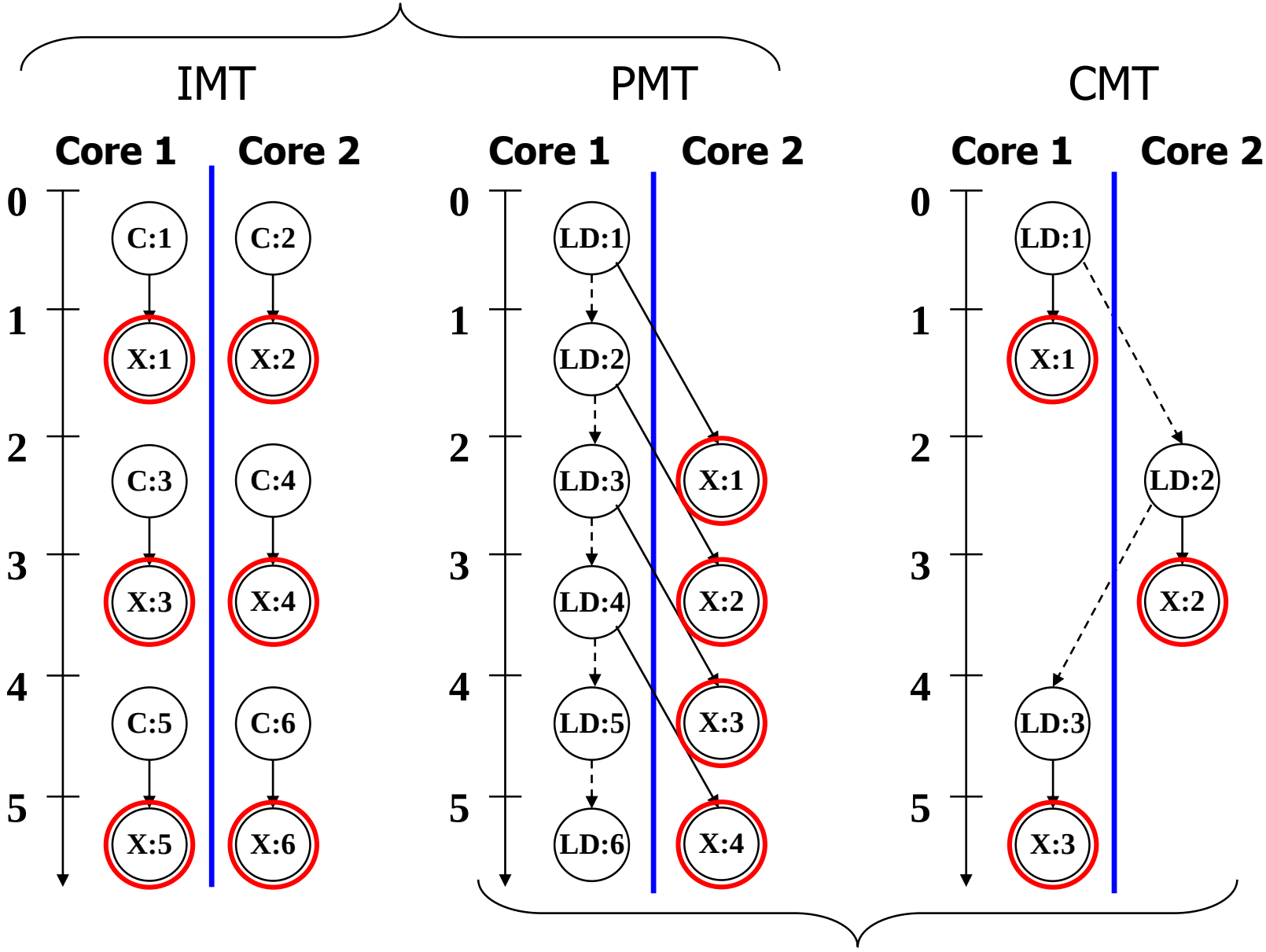
Comparison: IMT, PMT, CMT



	IMT	PMT	CMT
lat(comm) = 1:	1 iter/cycle	1 iter/cycle	1 iter/cycle
lat(comm) = 2:	1 iter/cycle	1 iter/cycle	0.5 iter/cycle

Comparison: IMT, PMT, CMT

Thread-local Recurrences ➔ Fast Execution

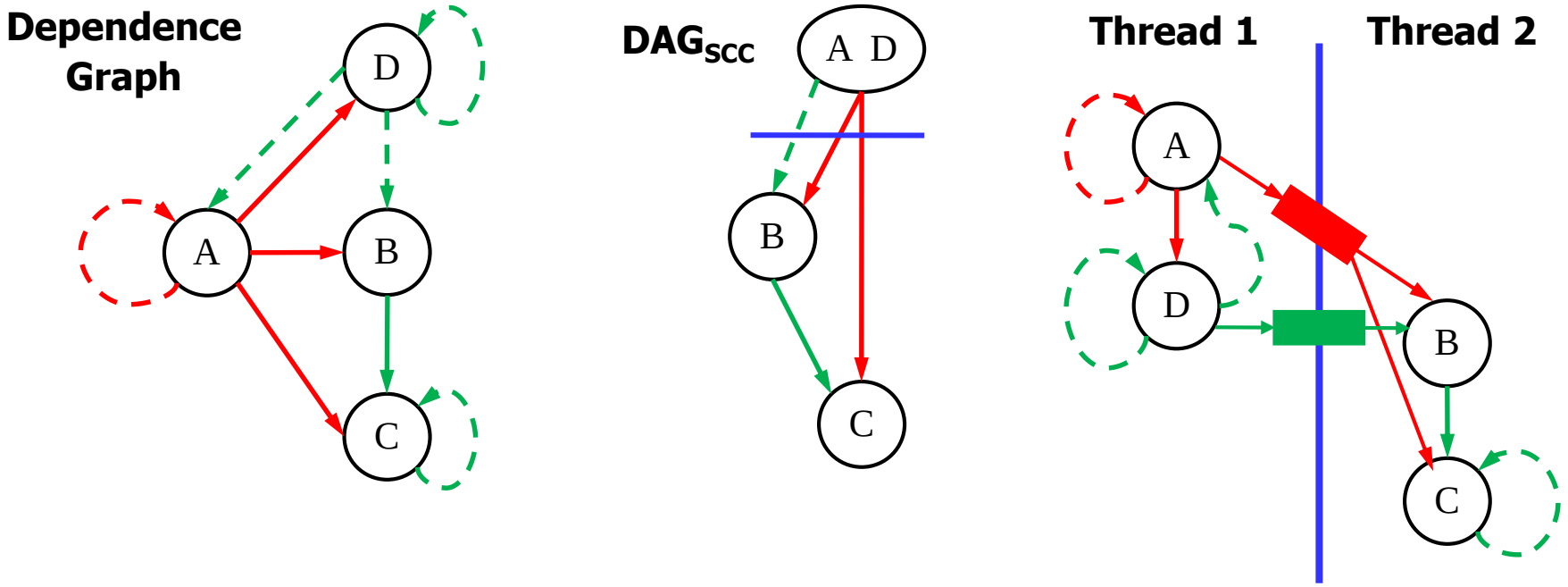


Cross-thread Dependences ➔ Wide Applicability

Decoupled Software Pipelining

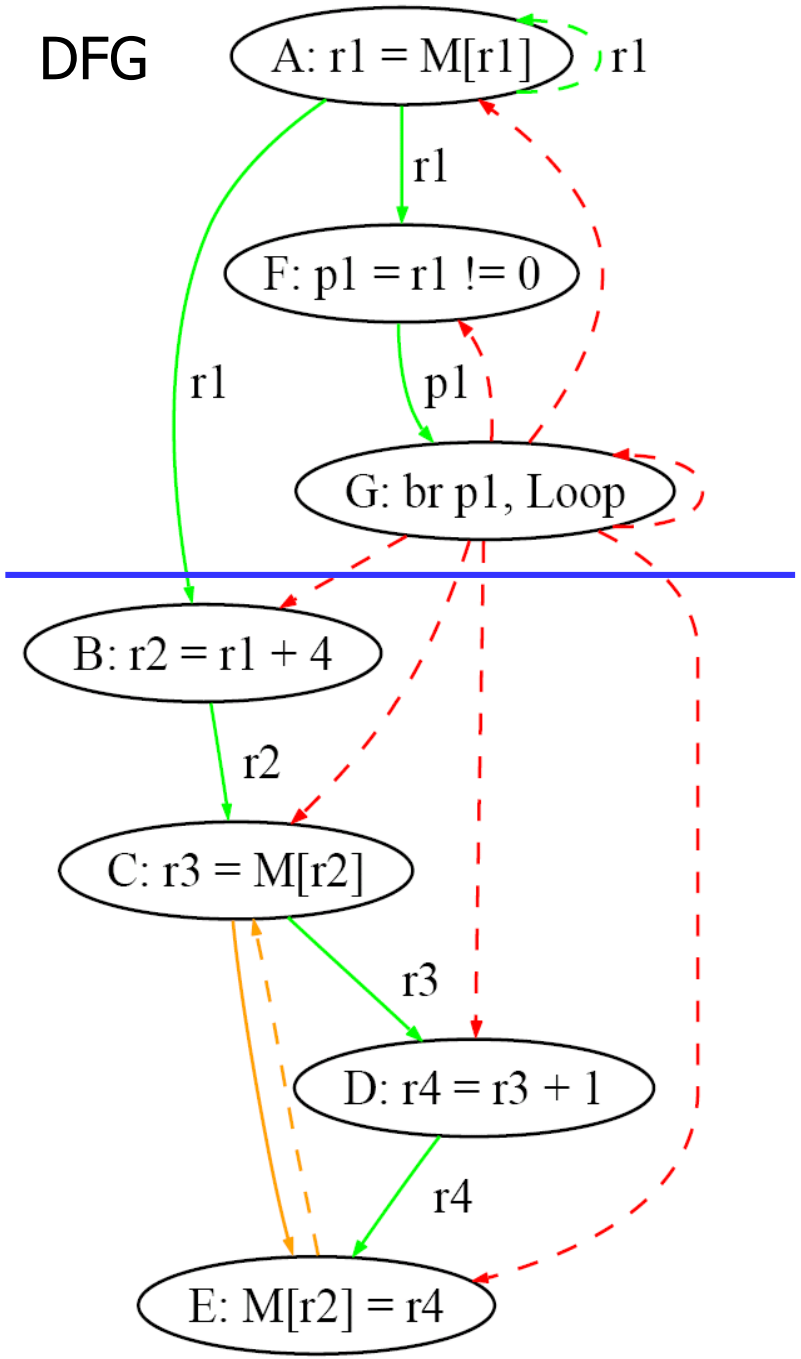
Decoupled Software Pipelining (DSWP)

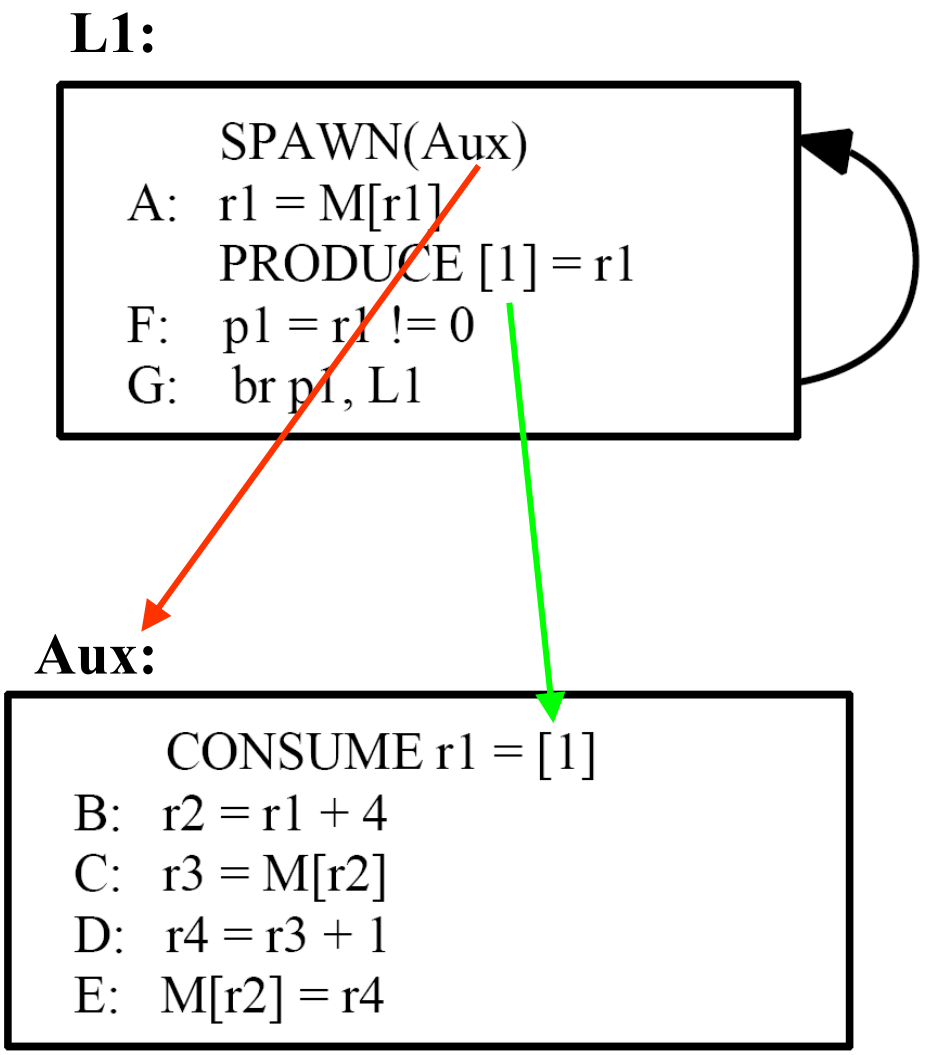
```
A: while(node)
B:   ncost = doit(node);
C:   cost += ncost;
D:   node = node->next;
```



Inter-thread communication latency is a one-time cost

Implementing DSWP





register

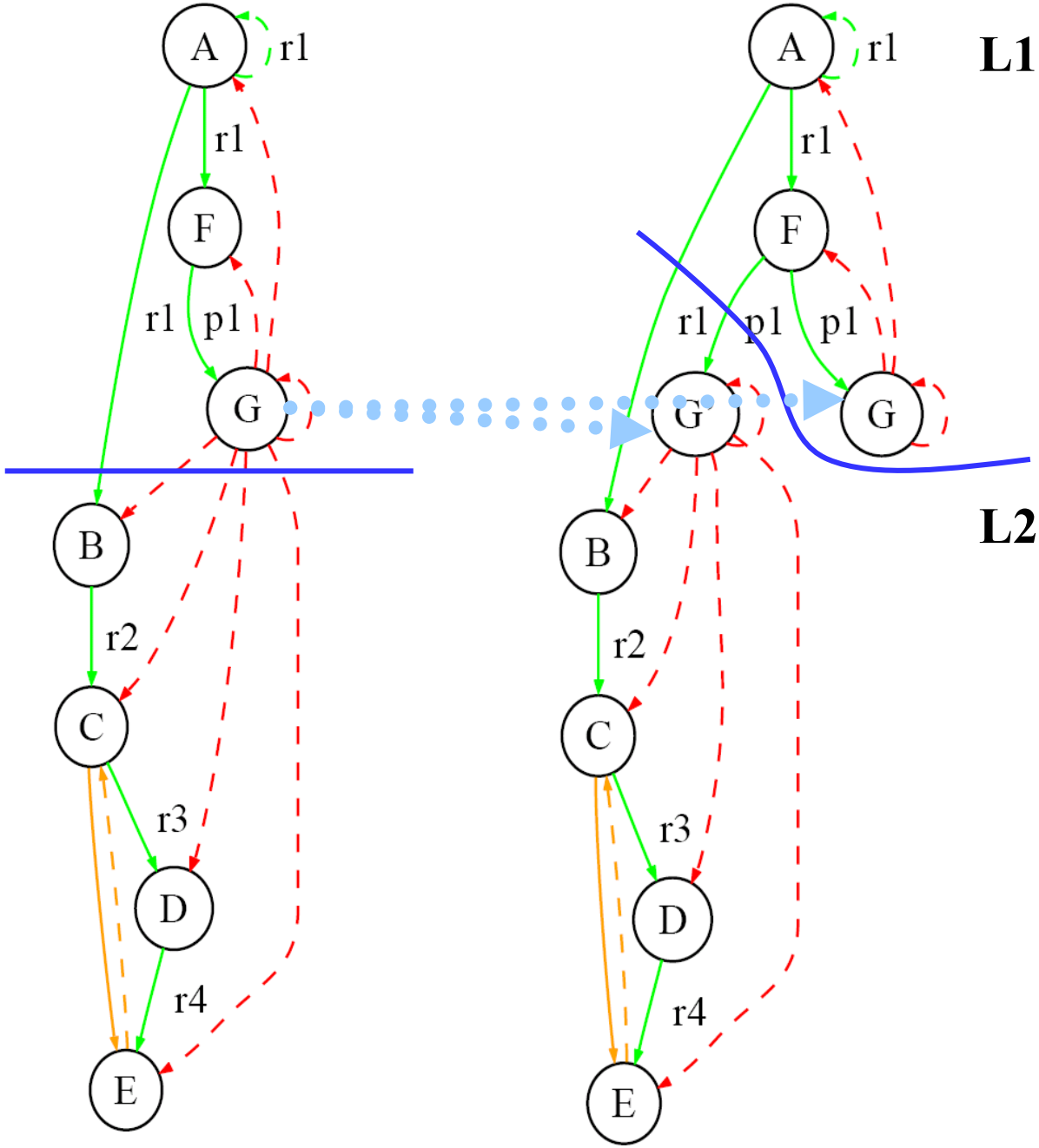
control

memory

intra-iteration

loop-carried

Optimization: Node Splitting To Eliminate Cross Thread Control



L1

```
A:  r1 = M[r1]
    PRODUCE [1] = r1
F:  p1 = r1 != 0
    PRODUCE [2] = p1
G:  br p1, L1
```

L2

```
    CONSUME r1 = [1]
B:  r2 = r1 + 4
C:  r3 = M[r2]
D:  r4 = r3 + 1
E:  M[r2] = r4
    CONSUME p1 = [2]
G': br p1, L2
```

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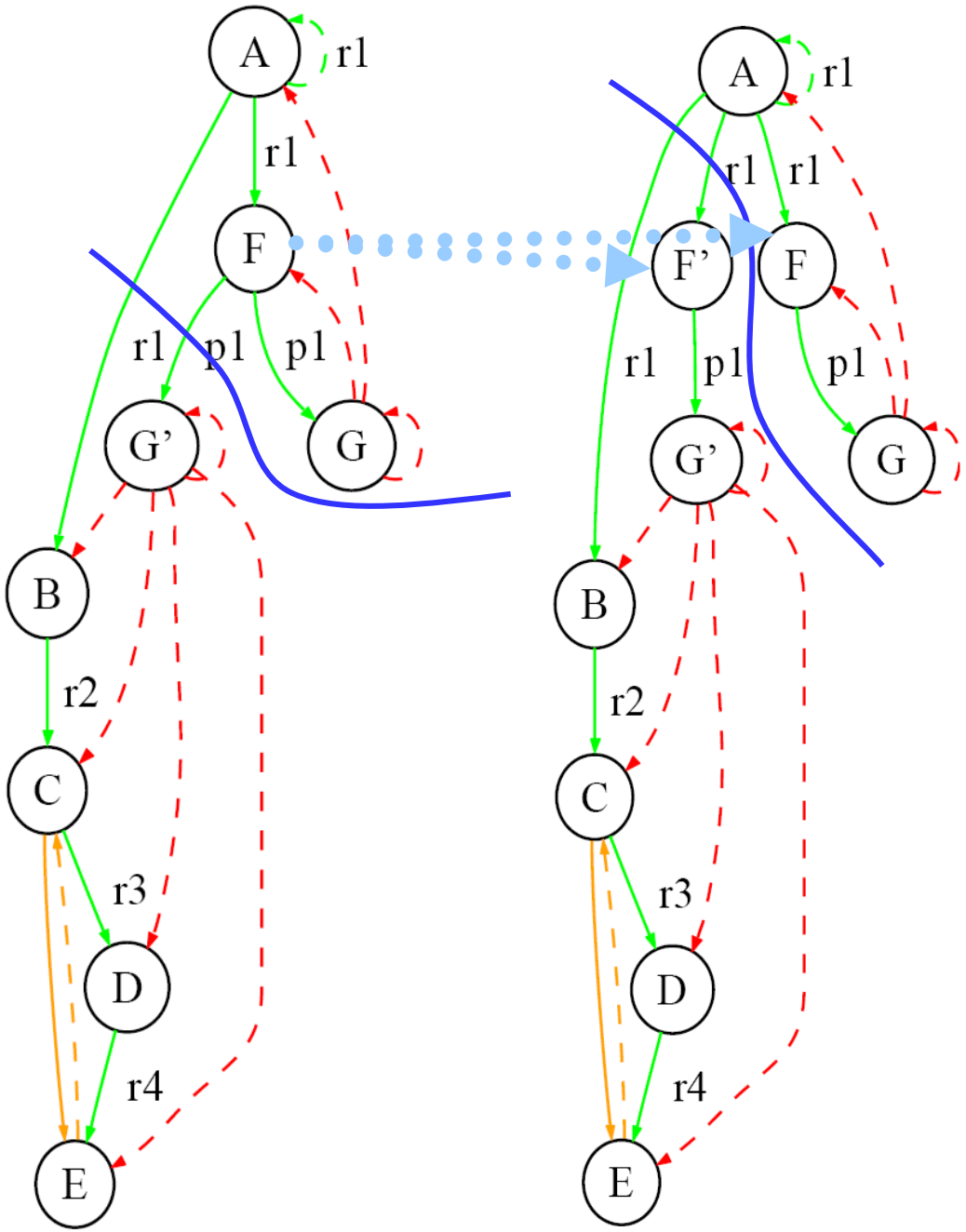
Optimization: Node Splitting To Reduce Communication



L1

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L2

```
     CONSUME r1 = [1]
B:   r2 = r1 + 4
C:   r3 = M[r2]
D:   r4 = r3 + 1
E:   M[r2] = r4
F':  p1 = r1 != 0
G':  br p1, L2
```

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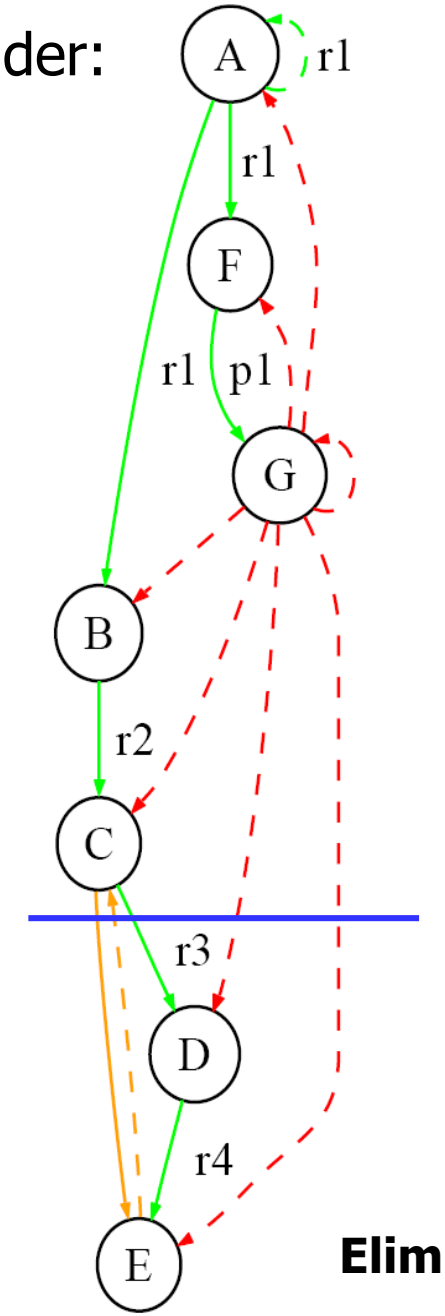
Constraint: Strongly Connected Components

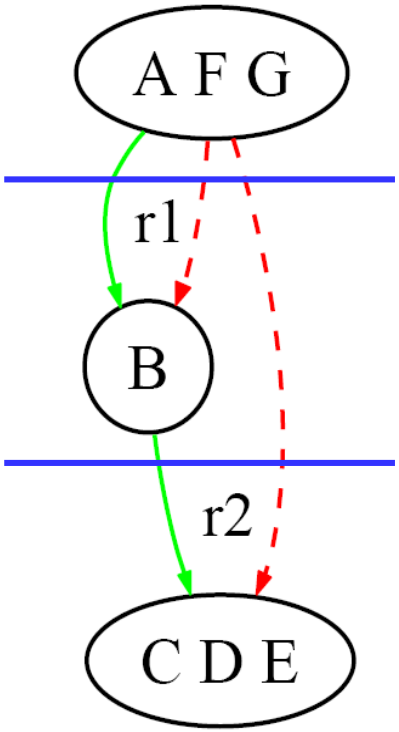
Consider:




```
      SPAWN(Aux)
A:    r1 = M[r1]
B:    r2 = r1 + 4
C:    r3 = M[r2]
      PRODUCE [1] = r3
      CONSUME r0 = [2]
F:    p1 = r1 != 0
G:    br p1, L1
```

```
      CONSUME r3 = [1]
D:    r4 = r3 + 1
E:    M[r2] = r4
      PRODUCE [2] = r0
```

Solution: DAG_{SCC}



Eliminates pipelined/decoupled property

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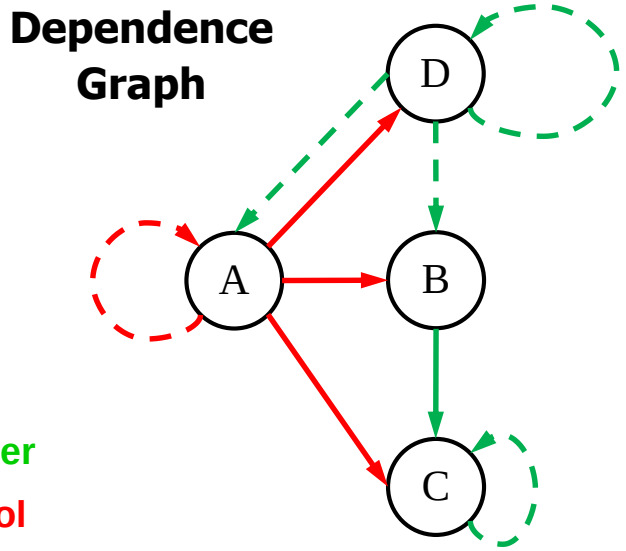
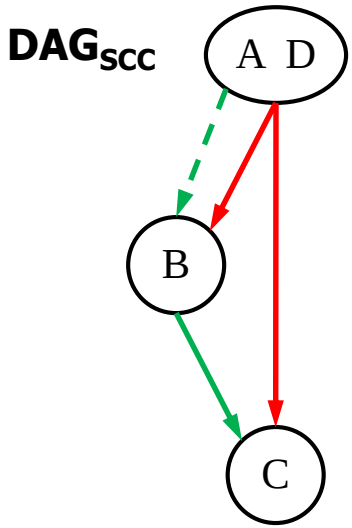
Speculation – Break Statistically Unlikely Dependences

```
A: while(node)
B:    ncost = doit(node);
C:    cost += ncost;
D:    node = node->next;
```





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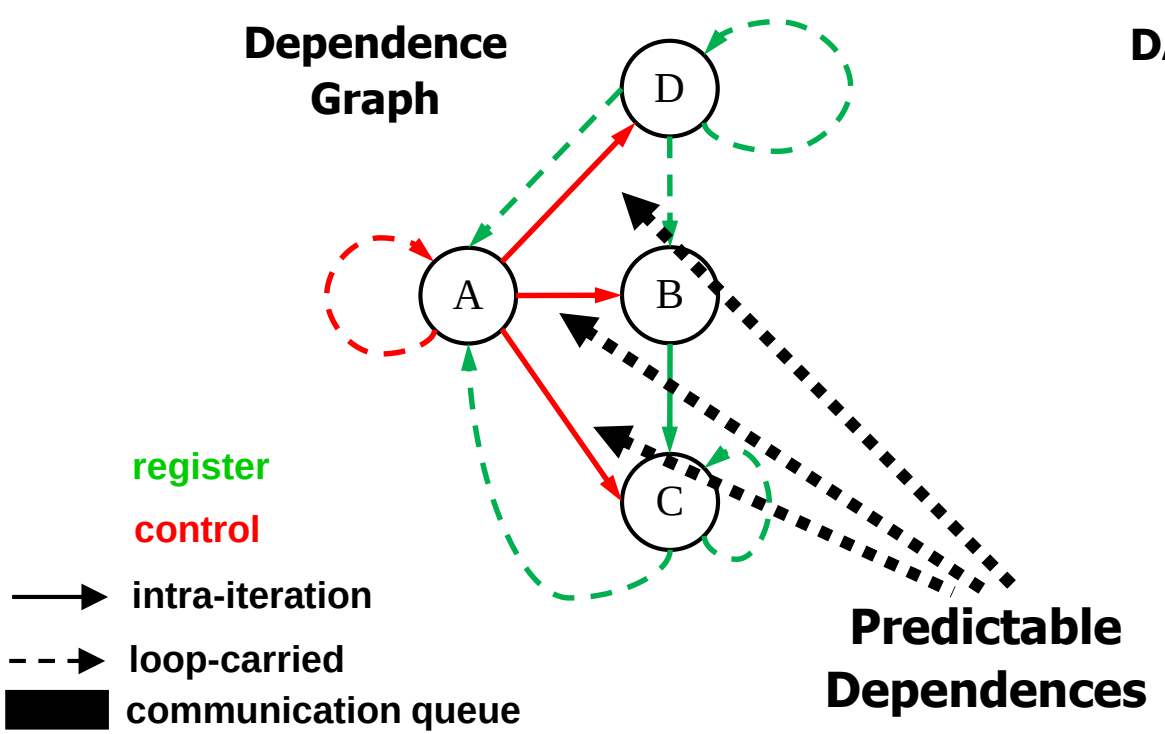
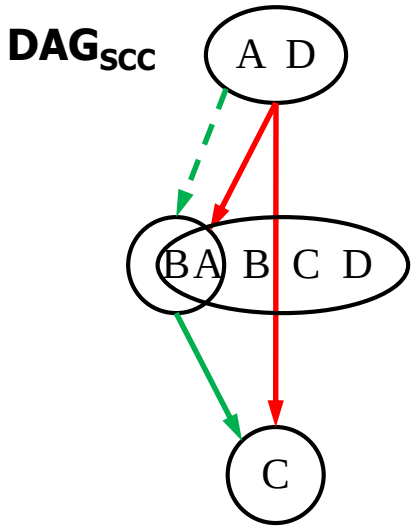
loop-carried

communication queue

Why Speculation?

```
A: while(cost < T && node)
B:   ncost = doit(node);
C:   cost += ncost;
D:   node = node->next;
```





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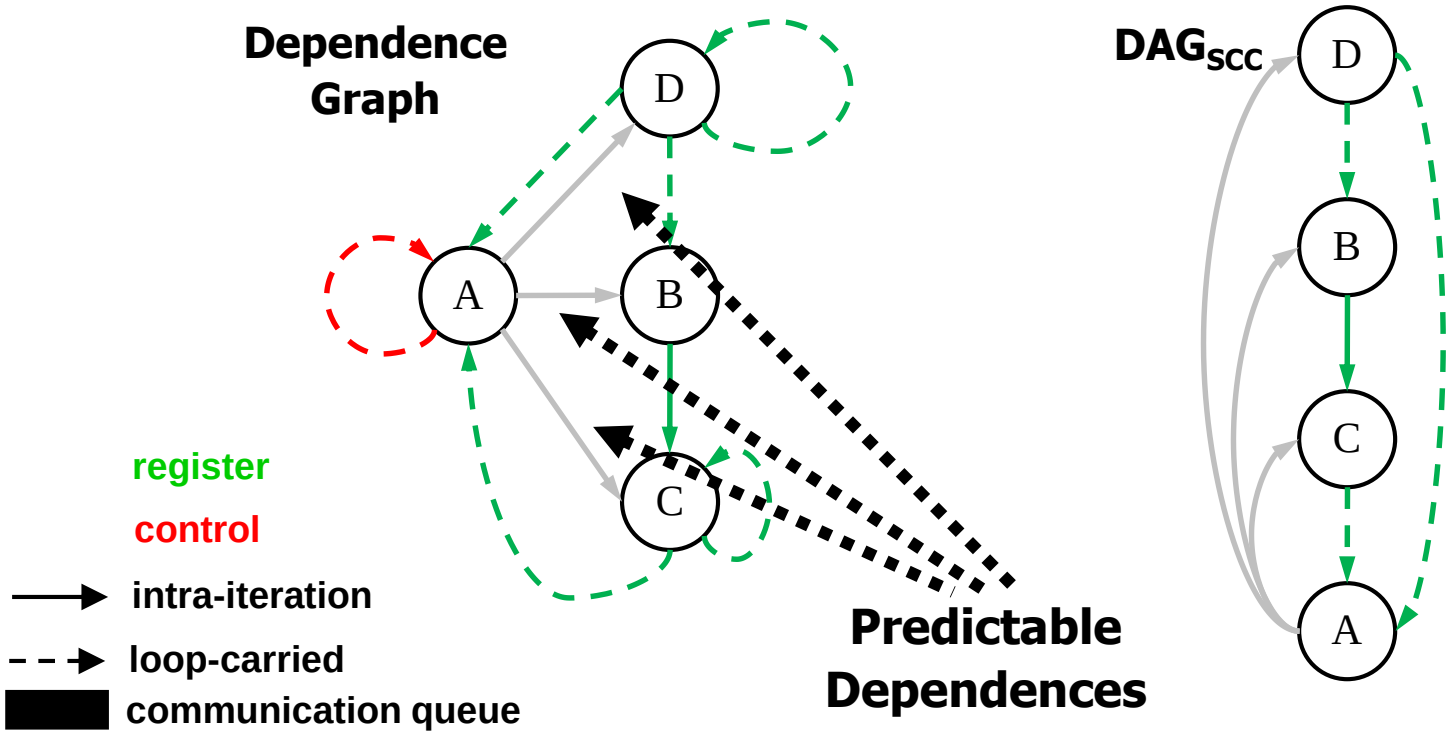
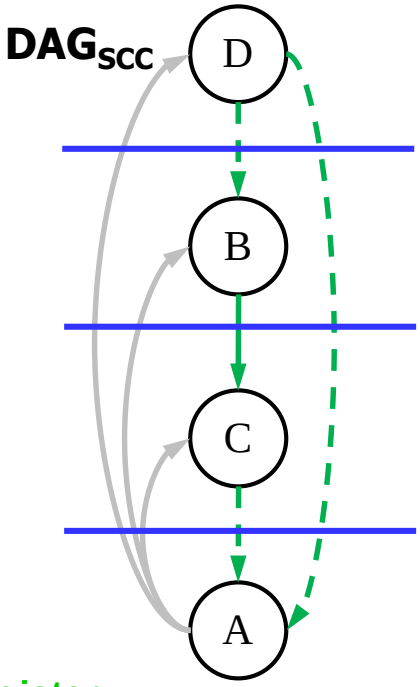


Execution Paradigm



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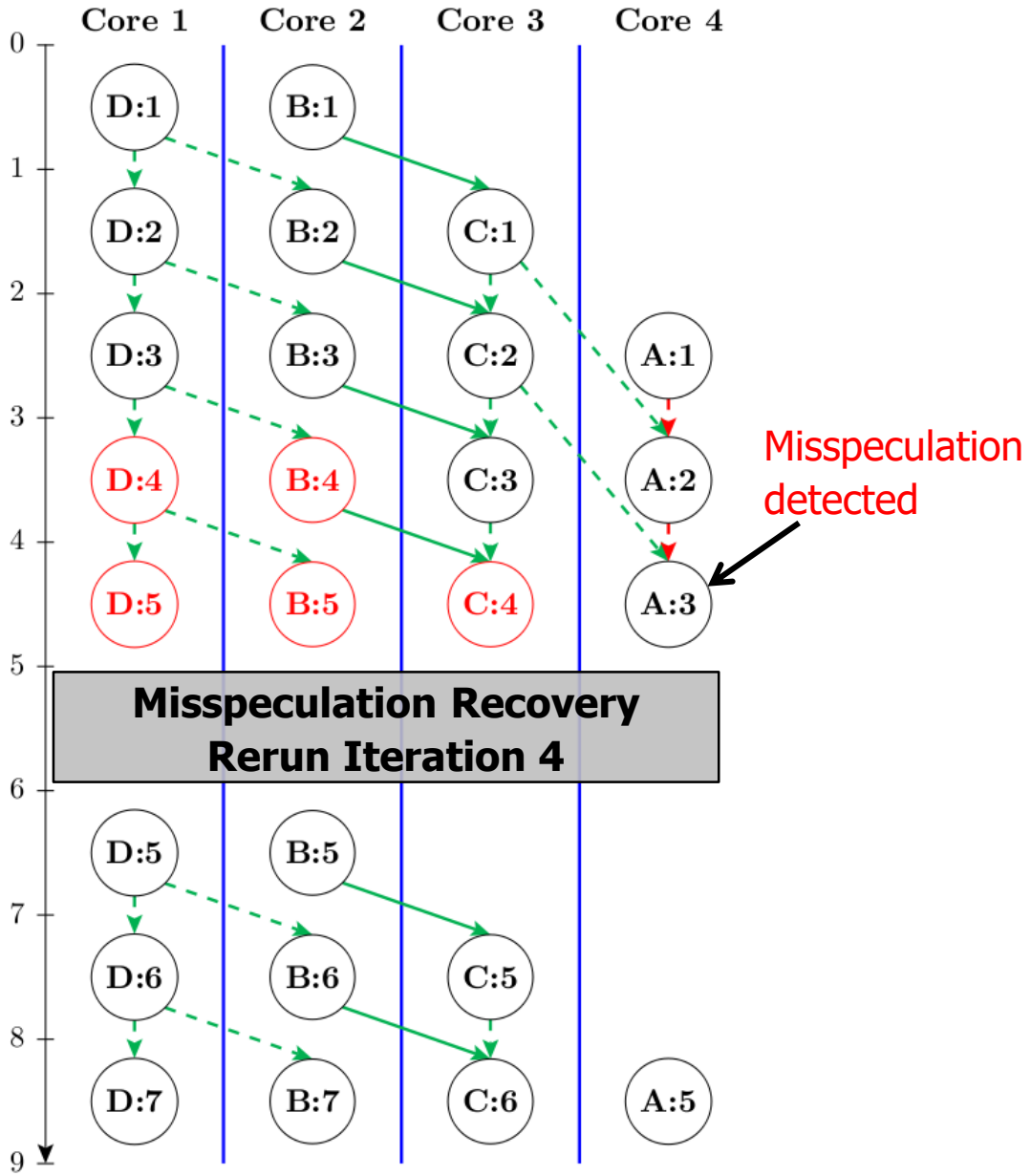
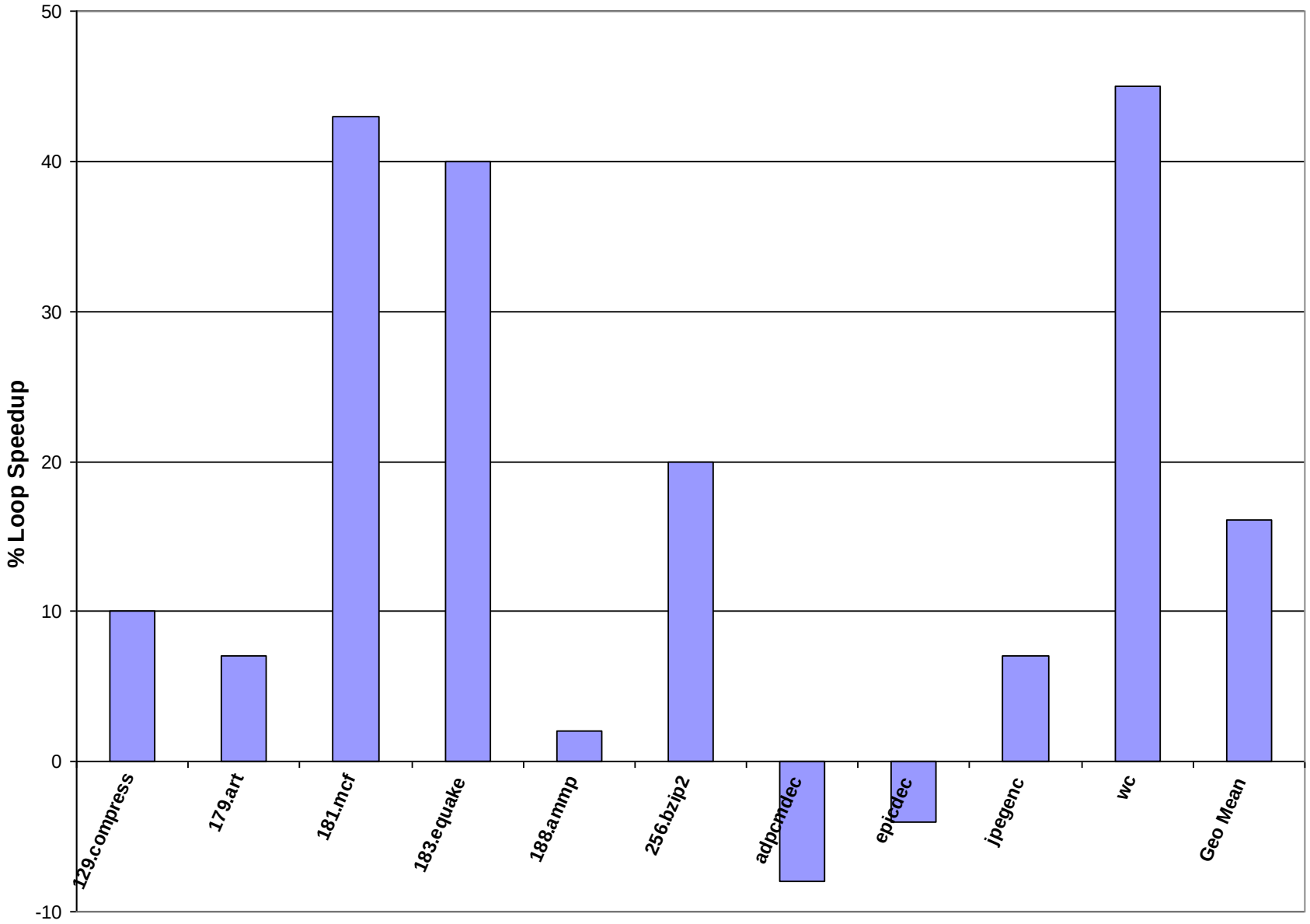


Evaluation: Dual Core vs Single Core



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References

- "Automatic Thread Extraction with Decoupled Software Pipelining," G. Ottoni, R. Rangan, A. Stoler, and D. I. August, *Proceedings of the 38th IEEE/ACM International Symposium on Microarchitecture*, Nov. 2005
- "Revisiting the Sequential Programming Model for Multi-Core," M. J. Bridges, N. Vachharajani, Y. Zhang, T. Jablin, and D. I. August, Proc 40th IEEE/ACM International Symposium on Microarchitecture, December 2007.