

# HW2 – Frequent Path Loop Invariant Code Motion

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# Loop Invariant Code Motion (LICM)

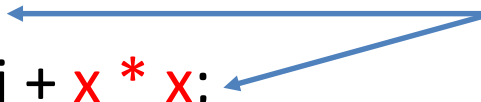
```
for (int i = 0; i < n; i++) {
```

```
    x = y + z;
```

```
    a[i] = 6 * i + x * x;
```

```
}
```

Their values don't  
change within the loop



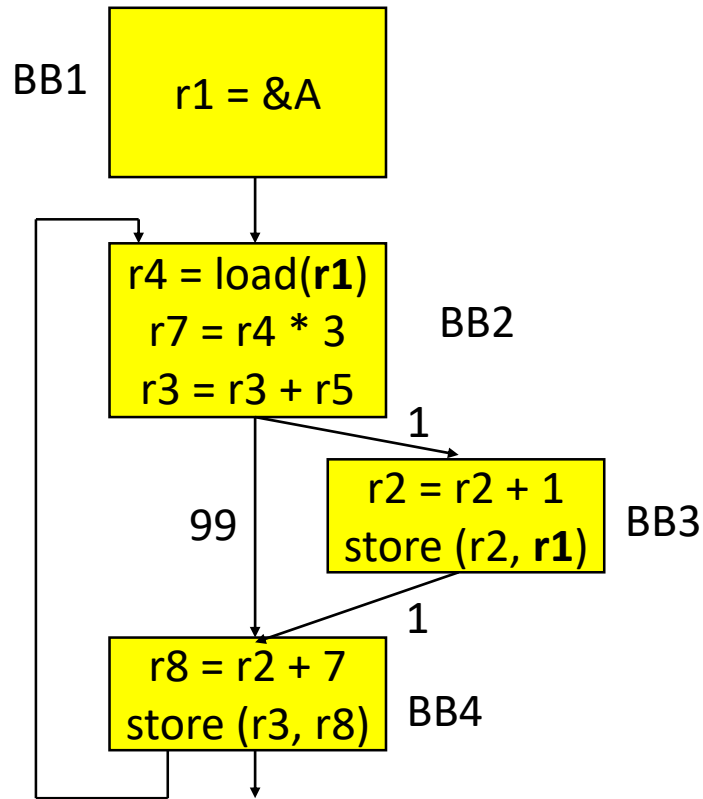
# Loop Invariant Code Motion (LICM)

```
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    x = y + z;  
    a[i] = 6 * i + x * x;  
}
```

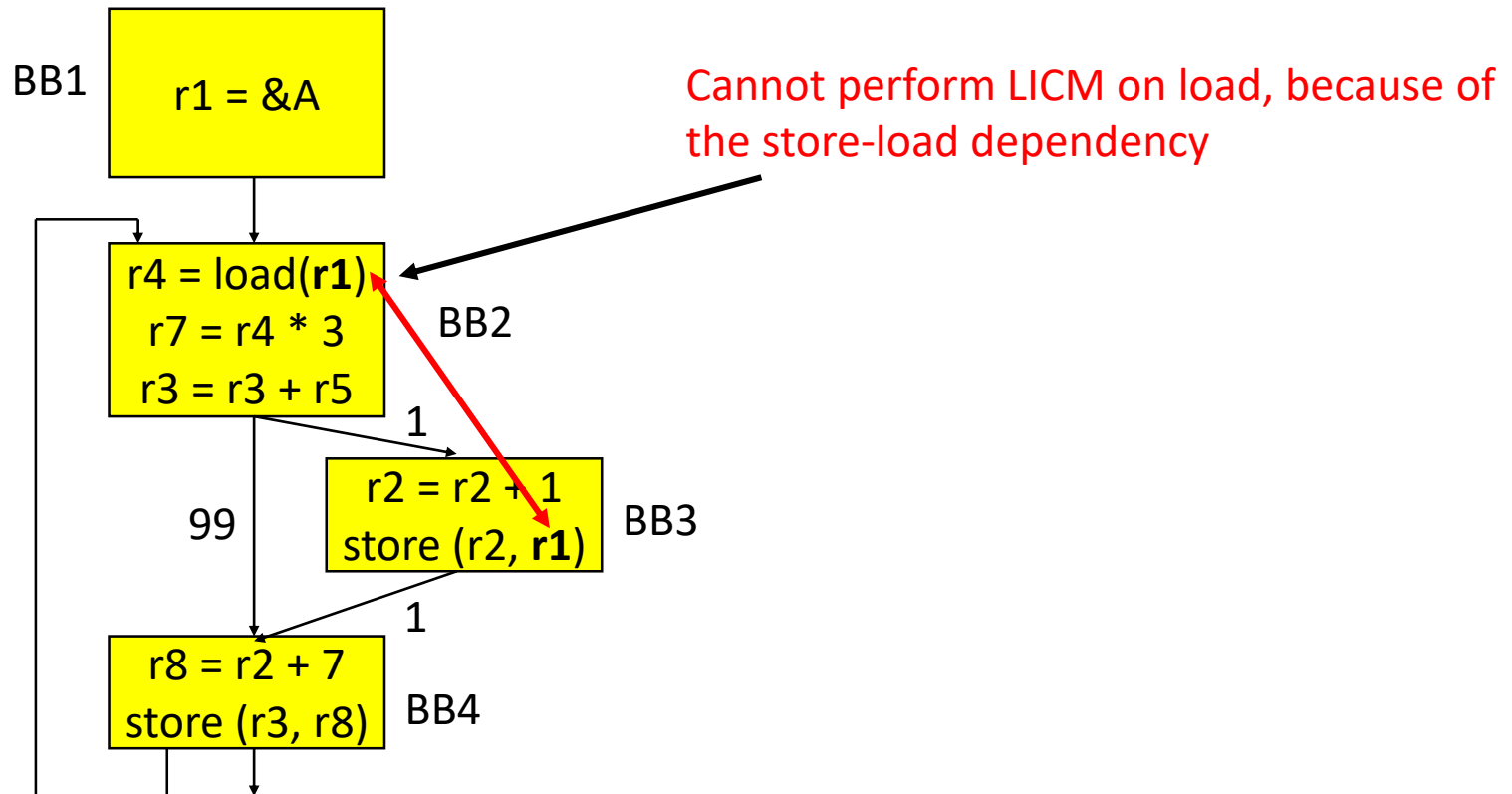
```
x = y + z;  
t1 = x * x;  
for (int i = 0; i < n; i++) {  
    a[i] = 6 * i + t1;  
}
```

- Move operations whose source operands do not change within the loop to the loop preheader
  - Execute them only 1x per invocation of the loop
  - Be careful with memory operations!
  - Be careful with ops not executed every iteration
- LICM code exists in LLVM!
  - /lib/Transforms/Scalar/LICM.cpp

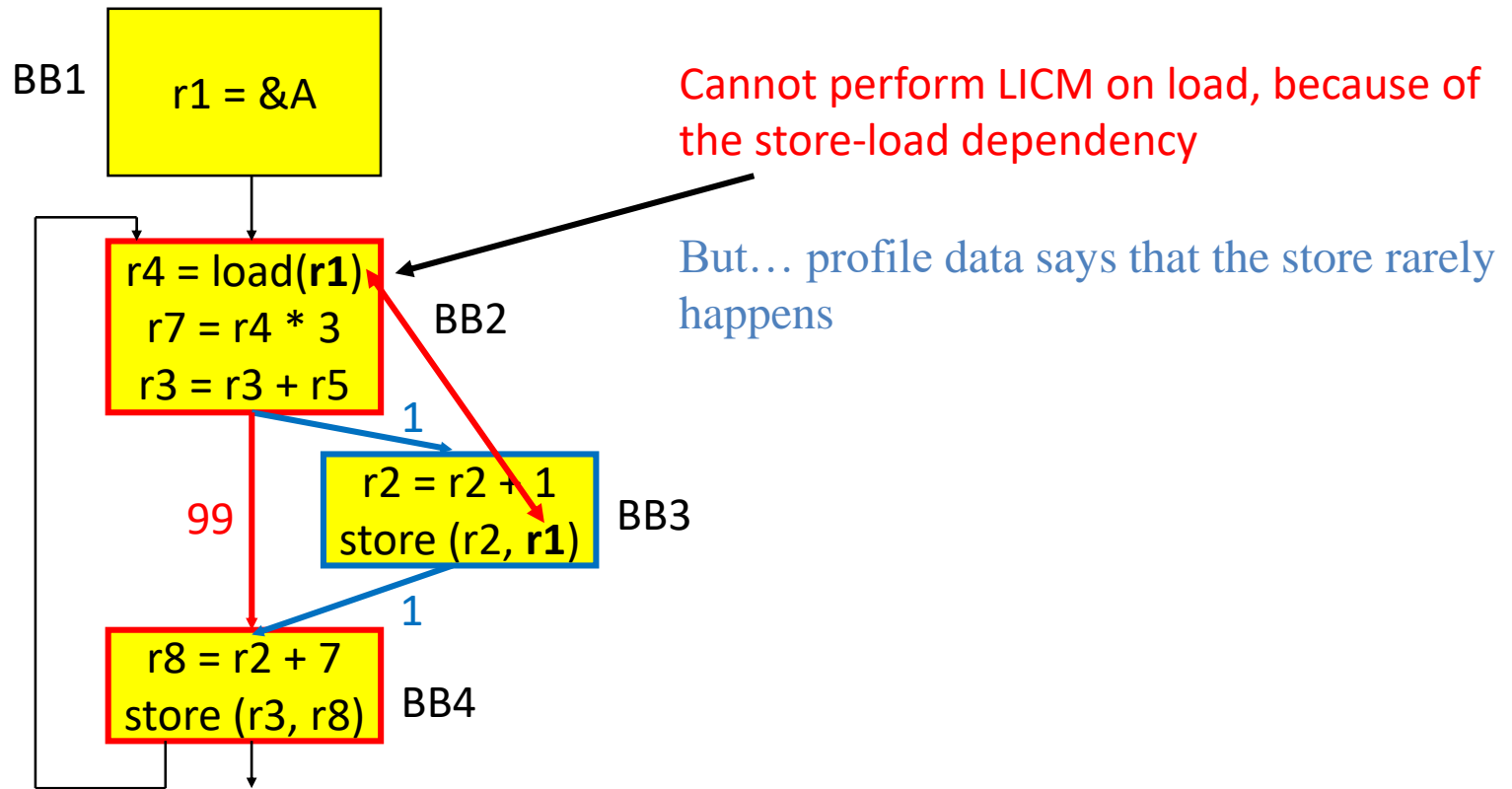
# Your Assignment: Frequent Path LICM



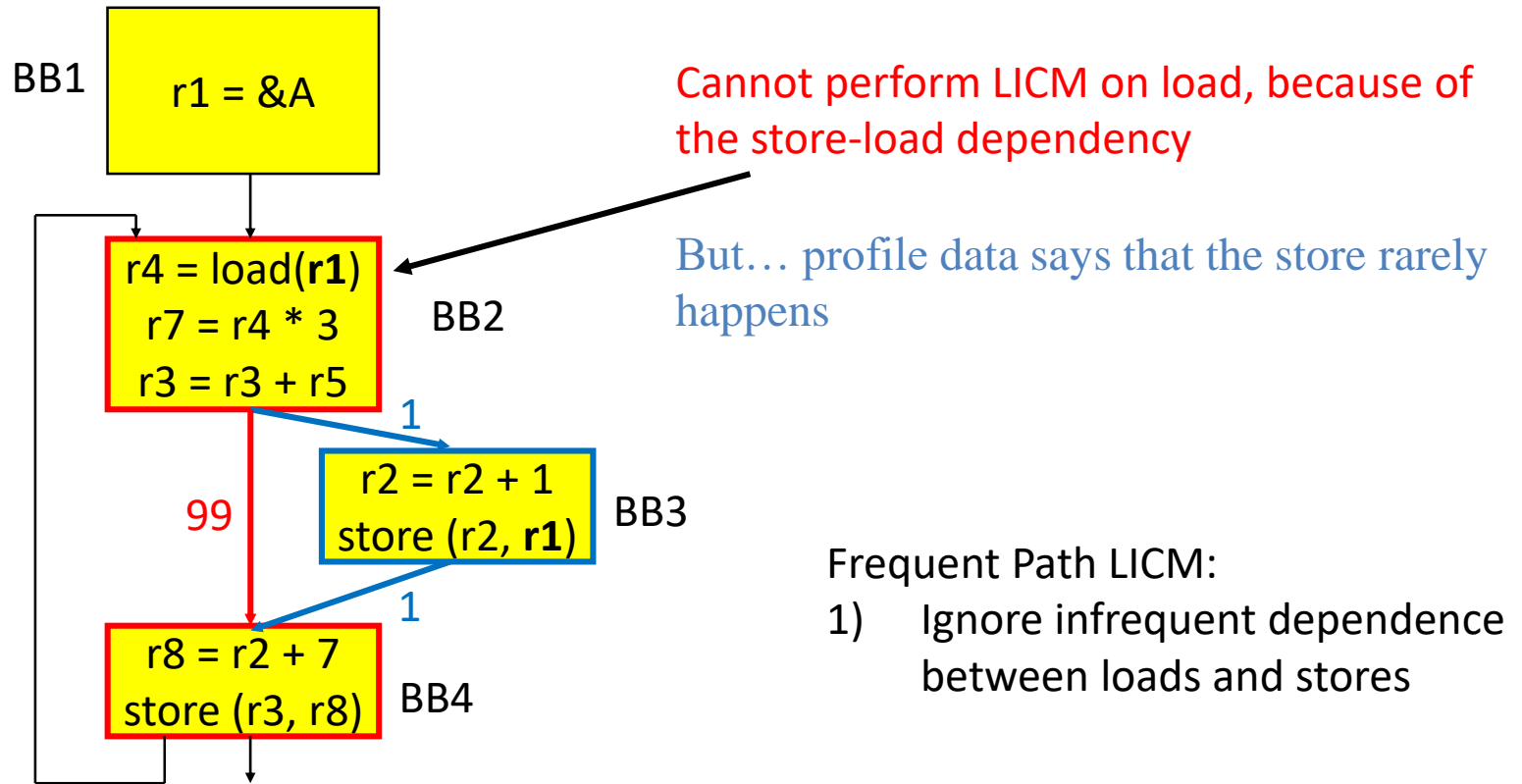
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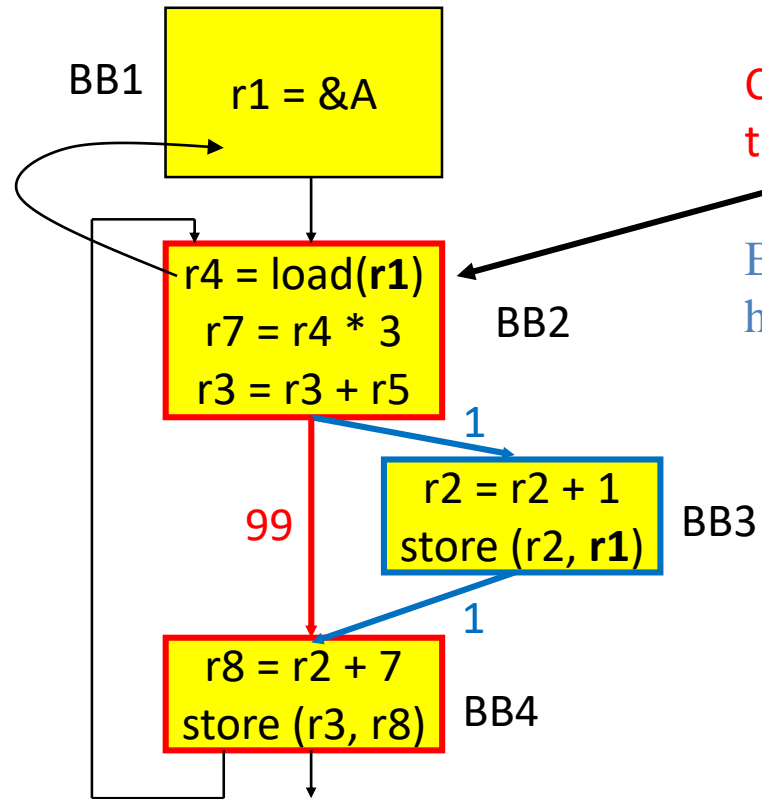
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Cannot perform LICM on load, because of the store-load dependency

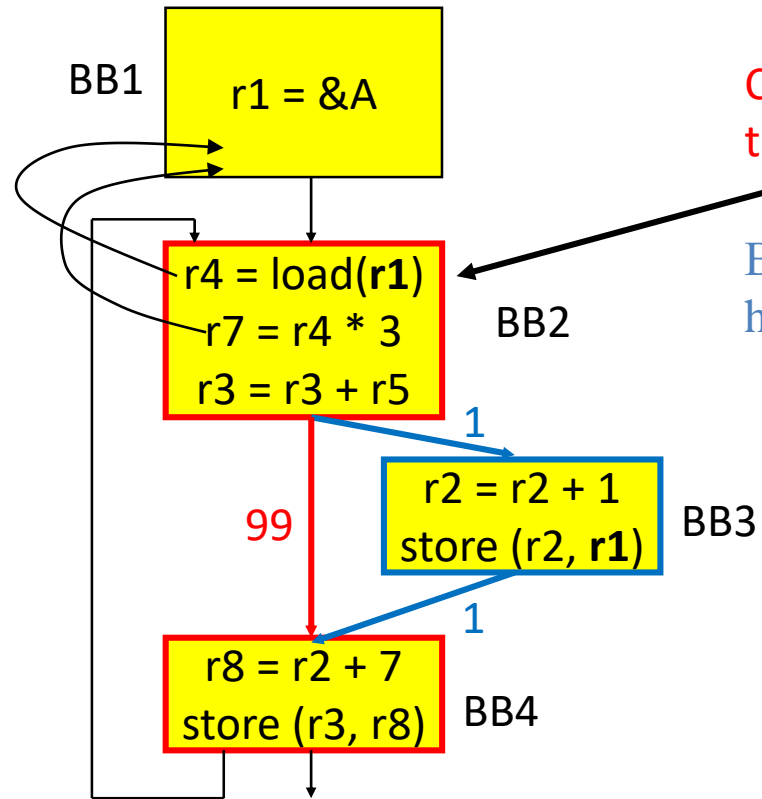
But... profile data says that the store rarely happens

Frequent Path LICM:

- 1) Ignore infrequent dependence between loads and stores
- 2) Perform LICM on load



# Your Assignment: Frequent Path LICM



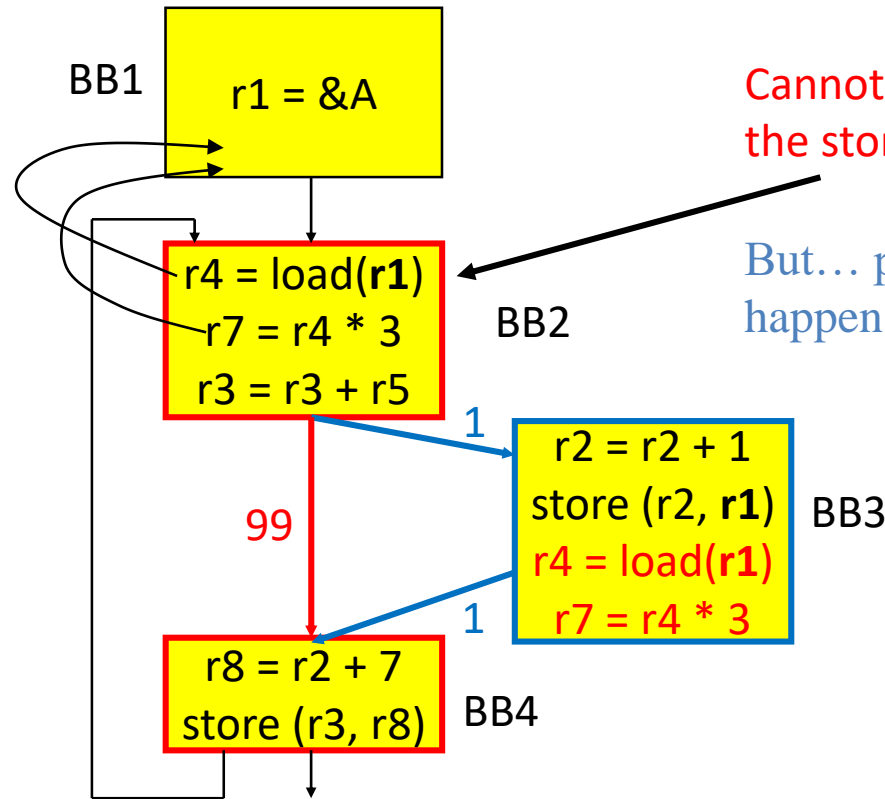
Cannot perform LICM on load, because of the store-load dependency

But... profile data says that the store rarely happens

Frequent Path LICM:

- 1) Ignore infrequent dependence between loads and stores
- 2) Perform LICM on load
- 3) Perform LICM on any consumers of the load that become invariant

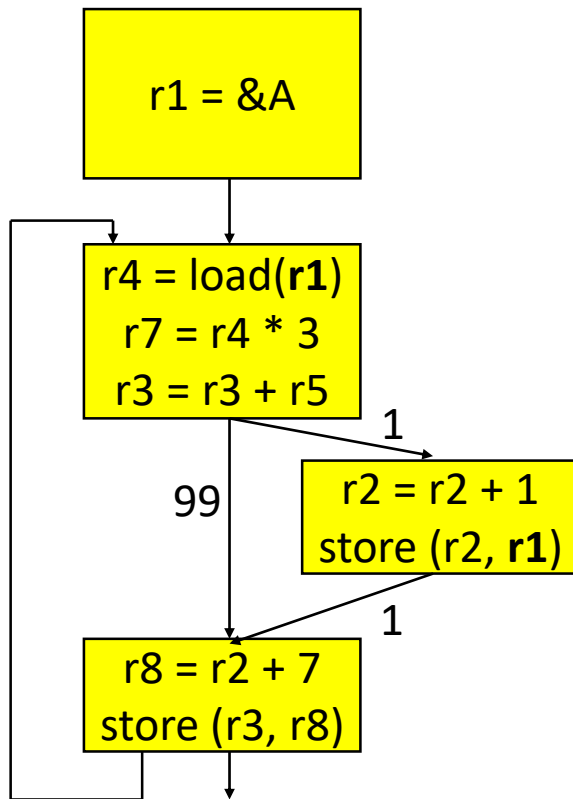
# Your Assignment: Frequent Path LICM



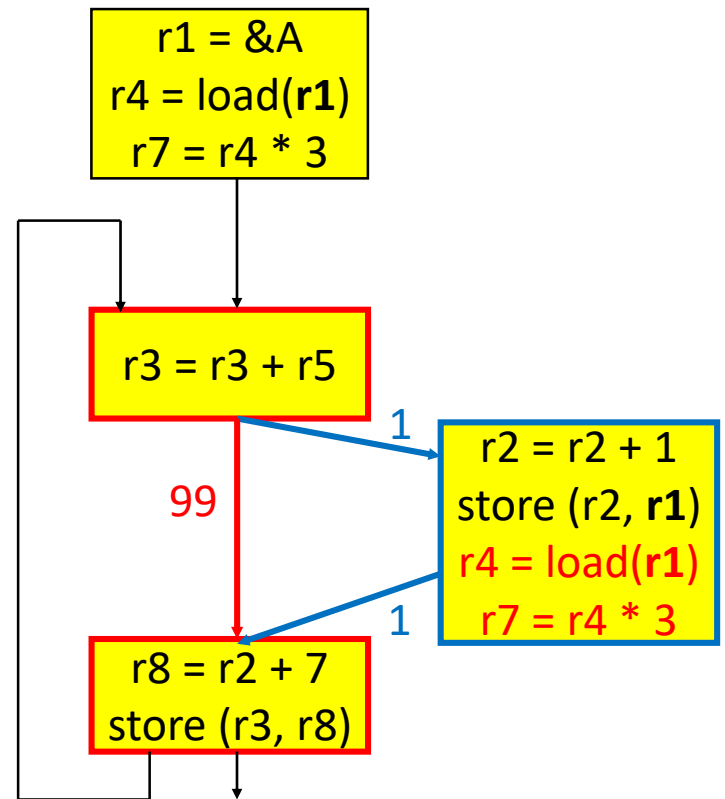
Frequent Path LICM:

- 1) Ignore infrequent dependence between loads and stores
- 2) Perform LICM on load
- 3) Perform LICM on any consumers of the load that become invariant
- 4) Insert fix-up code to restore correct execution

# Your Assignment: Frequent Path LICM



Before FPLICM



After FPLICM

# HW2: Frequent Path (FP) LICM

## Correctness:

- Identify the Frequent Path (edge probability  $\geq 80\%$ )
- Find store instructions among all infrequent BBs and their dependent load instructions in frequent BBs

*destination operand of infrequent store = source operand of frequent load*

- Hoist the FP invariants: Load instruction
- Replicate all hoisted instructions in the infrequent path

## Performance:

- Create a heuristic that determines to perform FP LICM or not.
  - Smart heuristic should apply optimization when it's profitable.
- Hoist the profitable FP invariants.
  - Load instruction
  - Consumers of the load that become invariant\* (For bonus points)

# FPLICM: What constitutes to FP

## Correctness:

- this can be accomplished by starting at the loop header and repeatedly following the  $\geq 80\%$  branch until a  $\geq 80\%$  loop backedge is taken.
- \*Note: This means that the cumulative probability of a BB might be lower than 80%
- **Anything not on the frequent path is on the infrequent path.**

## Performance:

- tune the parameter to achieve the highest performance gains

# HW2: Useful Resources

- run.sh
  - List of commands used in HW2
  - Check correctness of your pass!
- Project Template
  - HW2PASS.cpp
  - *runOnLoop(...)*      *inSubLoop(...)*
- Visualization Script – will be on piazza later
- Benchmarks
  - 6 correctness tests + README (Required)
    - Only need to hoist the dependent load instructions
    - Must generate the correct output after applying your FPLICM pass
    - Only submit the file created after your pass could run. hw2correct1.fplicm.bc  
=> hw2correct1\_base.bc. You do NOT have to test your pass on the performance benchmarks
  - 4 performance tests + README (Optional)
    - Hoist as many instructions as possible
    - Correctness first, then the performance
    - Same thing. except rename to hw2perf1\_bonus.bc

# LLVM Code of Interest

- The following slides present code from the LLVM codebase that may help you with HW2.
- Disclaimers:
  - Use of following API is your choice. There are many ways to do this assignment.
  - You are free to use any other code that exists in LLVM 12.0.1 or that you develop.
  - **Read the documentation/source before asking for help!**  
<http://llvm.org/docs/ProgrammersManual.html#helpful-hints-for-common-operations>

# Code: Manipulating Basic Blocks

- `SplitBlock(...)` splits a BB at a specified instr, returns ptr to new BB that starts with the instr, connects the BBs with an unconditional branch

```
// I is an Instruction*
BasicBlock *BB1 = I->getParent();
BasicBlock *BB3 =
    SplitBlock(BB1, I);
BasicBlock *BB2 =
    SplitEdge(BB1, BB3);
```

- `SplitEdge(...)` will insert a BB between two specified BBs

- Code found in:

- `<llvm-src-root>/include/llvm/Transforms/Utils/BasicBlockUtils.h`
- `<llvm-src-root>/lib/Transforms/Utils/BasicBlockUtils.cpp`



# Code: Creating and Inserting Instructions

- Various ways to create & insert instructions

```
// 1) create load, insert at end of
//     specified basic block
LoadInst *LD =
    new LoadInst(Val,
                  "loadflag",
                  BB1);
```
  - Hint: Instructions have a **clone()** member function
  - See specific instruction constructors/member functions in:
    - `<llvm-src-root>/include/llvm/IR/Instruction.s.h`
  - See general instruction functions available to all instructions in:
    - `<llvm-src-root>/include/llvm/IR/Instruction.h`
- ```
// 2) create branch using Create
//     method, insert before BB1's
//     terminating instruction
Branch::Create(BB1, BB2, flag,
                BB1->getTerminator());

// 3) create a store inst that stores
//     result of LD to some variable
//     (related to next slide)
StoreInst *ST =
    new StoreInst(LD, var);
//     inserting store into code
ST->insertAfter(LD);
```

# Code: Creating Variables

- Use `AllocInst` to allocate memory space on the stack.

```
// 1) Create a variable in the
//      function Entry block
AllocInst *Val = new AllocInst(
    I-&gtgetType(),
    0,
    nullptr,
    Entry->getTerminator()
);
```

```
// 2) store to the variable
StoreInst *ST = new StoreInst(
    Result,
    Val,
    Entry->getTerminator()
);
```

# Important: Maintaining SSA Form

- Static Single Assignment form requires unique destination registers for each instruction
  - Replicated instructions in your infrequent BB will write to different regs compared to the instructions in the preheader!
  - Store results of hoisted instrs to stack variables (see prev. slide)
  - Make sure AllocInst's are in function's entry BB!

# General Notes Regarding HW2

- **Start early!**
- Will be released on 9/20 (Mon)
- Make sure your optimization doesn't break a program!
- Start with script/template.
- Try the bonus part
- Check the piazza
- Running/Debugging
- Performance Competition: Generate correct **AND** fast bitcode