

Removing Impediments to Loop Fusion Through Code Transformations

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What is loop fusion?

- Optimization technique that takes several loops and combines them into a single large loop
- Decreases the number of loop branches executed
- Creates opportunities for data reuse
- Generally we:
 - Scan the code to find pairs of normalized loops that can be fused
 - Greedily fuse them



Requirements for loop fusion

- Are control equivalent
- Have no dependencies
- Have conforming bounds *
- No intervening code between the loops *

*: these conditions are loosened in the proposed algorithm



Negative Effects of loop fusion

- Increased code size
- Increased register pressure within a loop
- Potential over-committing of hardware resources
- Formation of loops with more complex control flow



Loop Fusion Basic Example



Problem: Loop fusion takes place under very specific conditions



Overview

- Code often has only small sets of loops that meet general code fusion requirements
- Lays out new algorithms to make code more fuseable
 - Eliminating conditions that prevent loop fusion
 - Generating maximal fusion -> increases scope of later transformations
- Focuses on IBM XL compiler suite
- Significant improvements to number of fused loops as a result of new algorithms



Conventional Loop Optimizations



- Optimization transformations in conventional XL compilers
- Paper's innovations focuses on loop fusion
 - Increases scope for optimization in the loop distributor



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LOOPFUSION				
1.	foreach	NestLevel N_i from outermost to innermost		
2.		Gather identically control dependent loops in N_i		
		into LoopSets		
3.		foreach LoopSet S_i		
4.		Remove loops non-eligible for fusion from		
		S_i		
5.		$FusedLoops \leftarrow True$		
6.		Direction \leftarrow Forward		
7.		while $FusedLoops = True$		
8.		$\mathbf{if} S_i < 2$		
9.		break		
10		endif		
11		Build Control Flow Graph		
12		Compute Dominance Relation		
13	i.	FusedLoops =		
		$LoopFusionPass(S_i, Direction)$		
14		if Direction $=$ Forward		
15		Direction = Reverse		
16		else		
17		Direction = Forward		
18		endif		
19).	endfor		
20		end while		
21	. endfor			

Loop Fusion Algorithm

- New loop fusion algorithm is the main innovation of the paper
- Creates opportunities for loop fusion:
 - Can fuse loops that don't conform
 - Can fuse loops with intervening code
- Note: loops are normalized

Loop Fusion Pass Example



(b) Fortran 77 Code

(a) After Fusing i1 and i2 into i5

(b) After moving intervening code up

Fig. 6. Completing first forward pass in running example



Loc	PFUSION	
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8.	$\mathbf{if} \; S_i < 2$	
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18.	endif	
19.	endfor	
20.	end while	rec
21.	endfor	

Loop Fusion Algorithm

- Loop through each nest level:
- Partitions all loops into sets of loops that are control equivalent
- Loop through each loop set:
- Remove loops non-eligible for fusion from the set
- Alternate forward and reverse passes over the set to fuse loops until no more loops can be fused
- CFG and dominance relations recomputed in each iteration

$LoopFusionPass(S_i, Direction)$			
1.	FusedLoops = False		
2.	foreach pair of loops L_j and L_k in S_i , such that L_j		
	dominates L_k , in Direction		
3.	if INTERVENINGCODE $(L_j, L_k) = True$ and		
	ISINTERVENINGCODEMOVABLE $(L_j, L_k) = False$		
4.	continue		
5.	endif		
6.	$\sigma \leftarrow \kappa(L_j) - \kappa(L_k) $		
7.	if L_j and L_k are non-conforming and		
	σ cannot be determined at compile time		
8.	continue		
9.	endif		
10.	if $DependenceDistance(L_j, L_k) < 0$		
11.	continue		
12.	endif		
13.	MOVEINTERVENINGCODE $(L_j, L_k, Direction)$		
14.	if INTERVENINGCODE $(L_j, L_k) = False$		
15.	if L_j and L_k are non-conforming		
16.	$L_m \leftarrow \text{FuseWithGuard}(L_j, L_k)$		
17.	else		
18.	$L_m \leftarrow \operatorname{Fuse}(L_j, L_k)$		
19.	endif		
20.	$S_i \leftarrow S_i \cup L_m - \{L_j, L_k\}$		
21.	FusedLoops = True		
22.	else		
23.	continue		
24.	endif		
25.	endfor		
26.	return FusedLoops		

Loop Fusion Algorithm

- Iterate through loop pairs in dominance order if forward, etc
- If the loops aren't adjacent and intervening code can't be moved; or the loops don't conform and the difference btwn. their upper bounds can't be determined; or there's neg. dependence distance between the loops; can't fuse
- Move intervening code
- Fuse conforming loops, fuse non-conforming loops with guard





Number of loops fused with each version of the compiler

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Execution times for selected benchmarks

Conclusion

- Strengths
 - Increased number of loops fused
 - Enables technology for optimization to take place later in this compiler framework, ie, loop distribution
 - Increase granularity of parallelism and minimize loop synchronization
- Weaknesses
 - Inhibit software pipelining
- Future Works
 - Removing restrictions that all loops that are fused must be control equivalent
 - Variation of index set splitting to remove control flow splits
 - Loop Dependence Graph to solve data cache performance





