

EECS598: Prediction and Learning: It's Only a Game

Fall 2013

Lecture 13: Universal Portfolios Review and Online Convex Optimization

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Announcements

- Homework 2 due in one week.
- It is time to think about the projects. There will be discussion on early November.

13.1 Universal Portfolios Review

Consider the problem with n stocks and $\mathbf{b}^t \in (0, \infty)^n$ as the “Price relatives” where $b_i^t = \frac{\text{Price}_{t+1}(i)}{\text{Price}_t(i)}$. Our decision is the portfolio $\mathbf{w} \in \Delta_n$, which means we have w_i fraction of wealth invest in stock i . From time t to time $t + 1$, the wealth grows by $\mathbf{w}\mathbf{b}^t = \sum_i w_i b_i^t$. Without loss of generality, we can consider the starting wealth $c = 1$.

13.1.1 Constant Rebalanced Portfolio

We focus on a certain class of policies called the constant rebalanced portfolio(CRP), where the \mathbf{w} is constant for each day. Then the total wealth of a CRP policy after T days would be $V_T(\mathbf{w}) = \prod_{t=1}^T (\mathbf{w} \cdot \mathbf{b}^t)$.

13.1.2 Cover’s Algorithm - UCRP

To compete with the best CRP under all scenarios, we consider a Universal CRP(UCRP), where the money is invested in *all* CRP’s evenly. In another word, all CRP got an infinitesimal amount of weight. We try to develop a regret bound of the UCRP to the best CRP.

Proof ideas: Define the ball around w_0 as $\text{Ball}_\epsilon(w_0) = \{w : (1 - \epsilon)w_0 + \epsilon v, \text{ where } v \in \Delta_n\}$. By this definition, we can see the following two properties:

- $\text{Vol}[\text{Ball}_\epsilon(w_0)] = \epsilon^{n-1} \text{Vol}(\Delta_n)$
- $V_T(w) \geq V_T(w_0)(1 - \epsilon)^T$

With these, we are ready to show the regret bound of the UCRP to the best CRP using these two properties.

Proof:

$$\begin{aligned}
 \text{Wealth(UCRP)} &= \frac{1}{\text{Vol}(\Delta_n)} \int_{w \in \Delta_n} V_T(w) d\mu(w) \\
 &\geq \frac{1}{\text{Vol}(\Delta_n)} \int_{w \in \text{Ball}_\epsilon(w^*)} V_T(w) d\mu(w) \\
 &\geq \frac{1}{\text{Vol}(\Delta_n)} \int_{w \in \text{Ball}_\epsilon(w^*)} d\mu(w) V_T(w^*) (1 - \epsilon)^T \\
 &= \epsilon^{n-1} (1 - \epsilon)^T V_T(w^*)
 \end{aligned}$$

It is clear that we need to tune the parameter ϵ to achieve a good bound. The best bound could be a little bit complex to achieve and we just let $\epsilon = \frac{1}{T}$ and it can give a good enough result.

To derive the bound, we let $\epsilon = \frac{1}{T}$ and look at the log-wealth of UCRP and CRP:

$$\begin{aligned}
 \text{Logwealth(UCRP)} &\geq -(n-1) \log \frac{1}{\epsilon} + T \log(1 - \epsilon) + \log V_T(w^*) \\
 &= -(n-1) \log T + O(1) + \log V_T(w^*)
 \end{aligned}$$

This gives the bound of UCRP. \square

The properties of UCRP:

- The regret is $O(n \log T)$ in log space.
- UCRP only very mild assumption on price relatives (could be very large or small).
- UCRP requires no tuning (although the analysis part need some tuning).
- Here we ignore the transition cost, but it is easy to be added. ([BK99])
- UCRP is not efficient to implement. (See [KV03] for details)
- Lots of work within Online Convex Optimization (OCO) framework on efficient algorithm with the same regret bound.

13.2 Online Convex Optimization (OCO)

The general framework in OCO looks like this:

- We have a decision space $X \in \mathbb{R}^n$ for $t = 1, \dots, T$. Assume that X is convex, closed and bounded.
- Player chooses $x_t \in X$, then nature chooses $l_t : X \xrightarrow{\text{convex}} \mathbb{R}$,
- Player suffers $l_t(x_t)$ and observe l_t , (we will go to bandit setting where player only observe $l_t(x_t)$ soon!)

We want to minimize the regret $\sum_{t=1}^T l_t(x_t) - \min_{x^* \in X} \sum_{t=1}^T l_t(x^*)$.

There are a variety of problems fall into the category of OCO, here are some examples:

Setting	Action	Data	Loss
Prediction with expert advice	$\mathbf{w} \in \Delta_n$	$(\mathbf{f}, y) \in (0, 1)^{n+1}$	$l(\mathbf{w}, \mathbf{f}, y)$
Action setting	$\mathbf{P} \in \Delta_n$	$\mathbf{1}_t \in [0, 1]^n$	$\mathbf{p} \mathbf{1}^t$
Linear Prediction (pattern recognition)	$\mathbf{w} \in \mathbb{R}_n$	$(\mathbf{x}, y) \in \mathbb{R} \times \{-1, 1\}$	$\max(0, y \mathbf{w} \mathbf{x})$
Portfolio Selection	$\mathbf{w} \in \Delta_n$	$\mathbf{b}^t \in (0, \infty)^n$	$-\log \mathbf{w} \mathbf{b}$

13.2.1 Online Gradient Decent (OGD)[Zino03]

First we define the projection function $\text{Proj}_X(z) = \text{argmin}_{x \in X} \|x - z\|_2$. Then we know: 1) $\text{Proj}_X(z) = z$ iff $z \in X$. 2) $\|x - z\|_2 \geq \|x - \text{Proj}_X(z)\|_2$.

The main idea of OGD method is the following:

- Choose an initial x_0 arbitrarily in X .
- For each $t = 1, \dots, T$
 - $z_{t+1} = x_t - \eta \nabla l_t(x_t)$,
 - $x_{t+1} = \text{Proj}_X(z_{t+1})$.

It can be shown that if $\nabla_t := \nabla l_t(x_t)$ has norm $\leq G$, and $D := \max_{x, y \in X} \|x - y\|_2$. Then $\text{regret}(\text{ODG}) \leq \eta \frac{G^2 T}{2} + \frac{D^2}{2\eta}$. Apply tuning to η and let $\eta = \frac{D}{G\sqrt{T}}$, we can have $\text{regret}(\text{ODG}) \leq GD\sqrt{T}$.

References

- [BK99] Avrim Blum and Adam Kalai. Universal portfolios with and without transaction costs. *Mach. Learn.*, 35(3):193–205, June 1999.
- [KV03] Adam Kalai and Santosh Vempala. Efficient algorithms for universal portfolios. *J. Mach. Learn. Res.*, 3:423–440, March 2003.
- [Zino03] Martin Zinkevich. Online convex programming and generalized infinitesimal gradient ascent. 2003.