

How to Use a Short Basis: **Trapdoors for Hard Lattices** **and New Cryptographic Constructions**

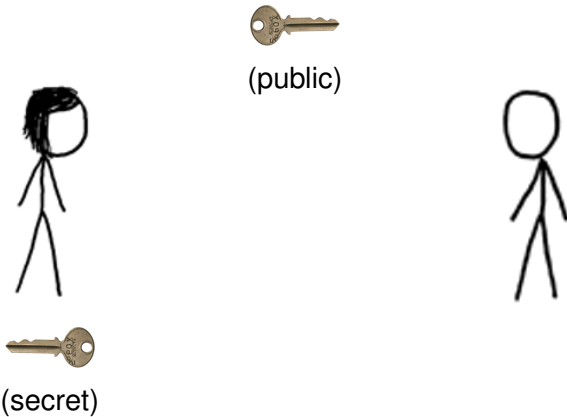
Chris Peikert
SRI

Work with Craig Gentry and Vinod Vaikuntanathan

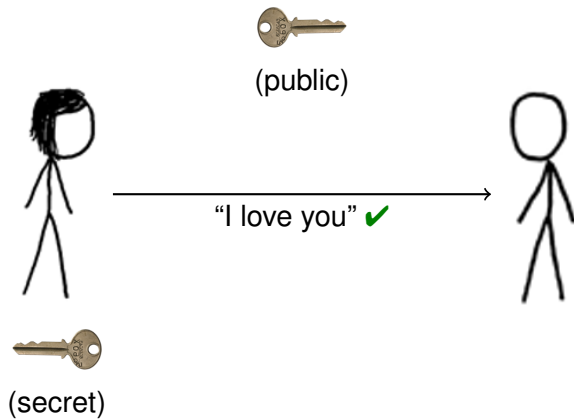
Digital Signatures



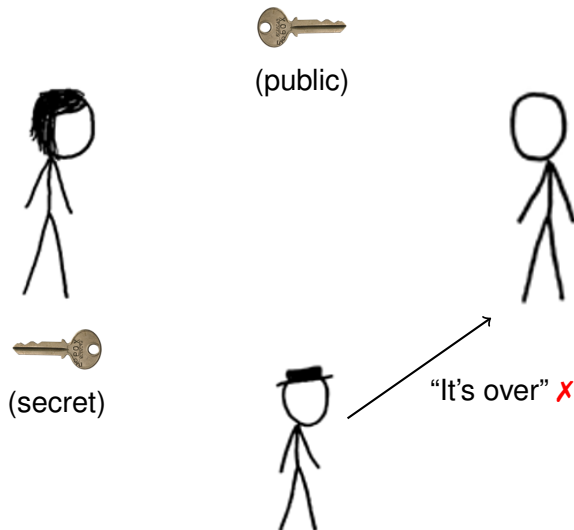
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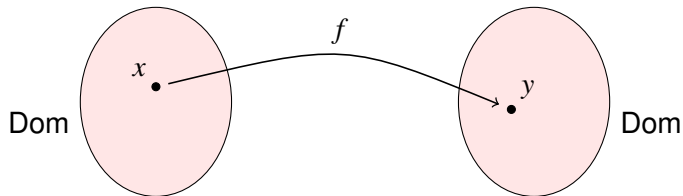


Trapdoor Permutations [DiffieHellman76]

- ▶ Public function f , secret “trapdoor” f^{-1}

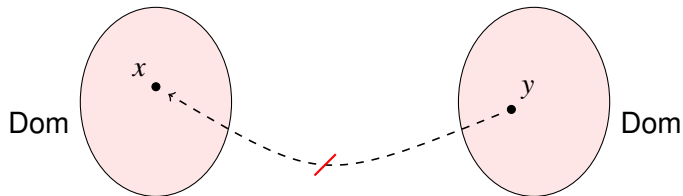
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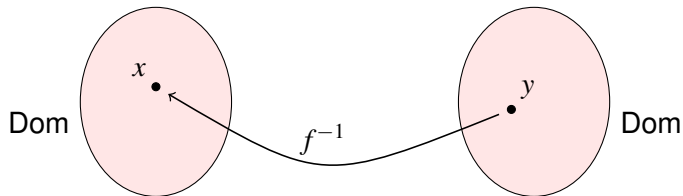
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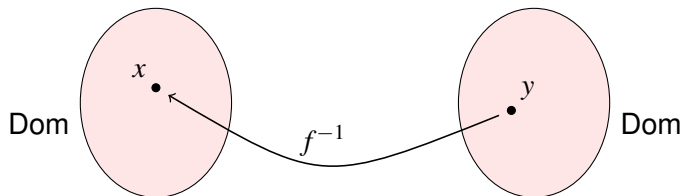
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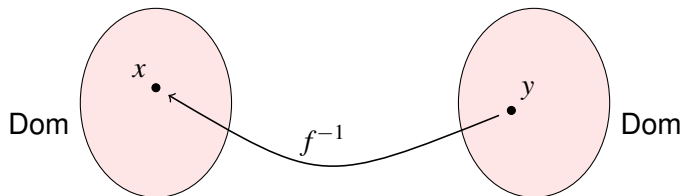
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 - ✓ “General assumption”
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- ▶ Candidates: [RSA78,Rabin79,Paillier99]
 - ✓ “General assumption”
 - ✓ Applications: digital signatures, OT, NIZK, ...
- ▶ All rely on hardness of **factoring**
 - ✗ Complex: 2048-bit exponentiation
 - ✗ Lack of diversity
 - ✗ Broken by quantum algorithms [Shor]

Lattice-Based Cryptography

What's To Like

- ▶ **Simple & efficient**: linear ops, small integers
- ▶ Resist **subexp & quantum** attacks (so far)
- ▶ Security from **worst-case** hardness [Ajtai, ...]

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- 1 One-way & collision-resistant functions [Ajtai,...,MicciancioRegev]
- 2 Public-key encryption [AjtaiDwork,Regev]
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What's Missing

- ▶ **Everything else!**

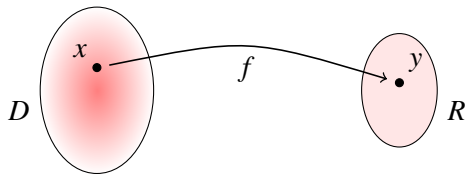
Practical signatures, protocols, “advanced” crypto, ...

Results: New Lattice-Based Crypto

- 1 Preimage sampleable trapdoor functions

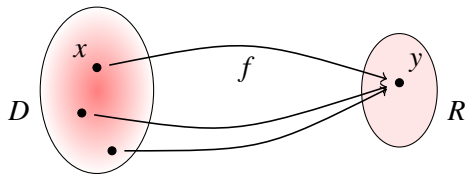
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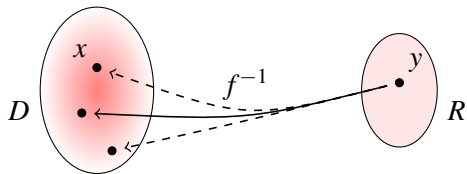
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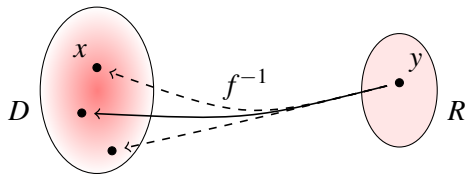
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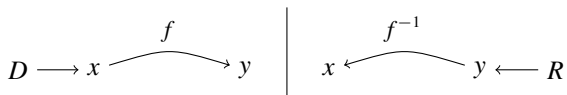


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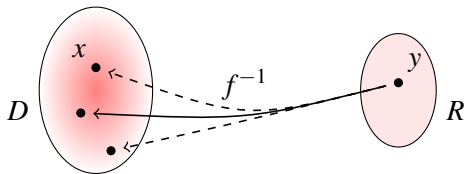


- Generate (x, y) in **two equivalent ways**:

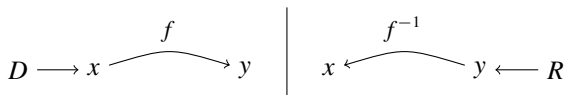


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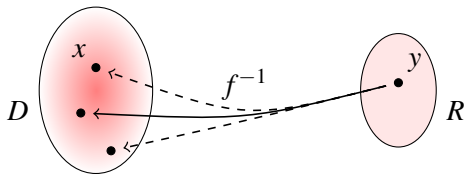
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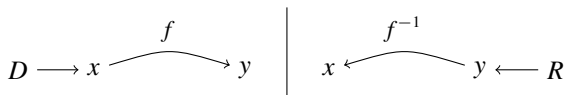
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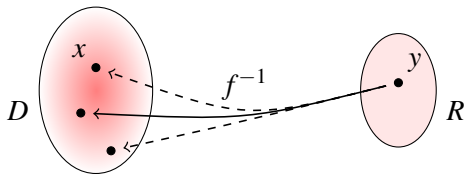


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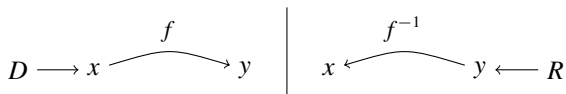
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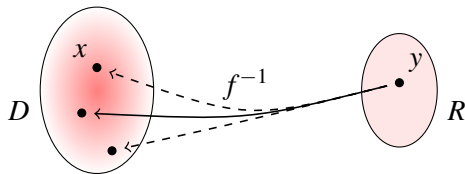
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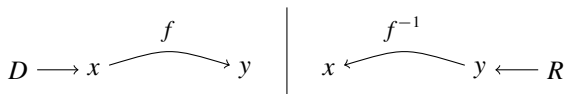
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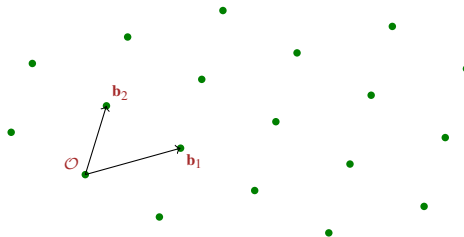
New Algorithmic Tool

- ▶ “Oblivious decoder” on lattices

Lattices

A lattice $\mathcal{L} \subset \mathbb{R}^n$ having basis $\mathbf{B} = \{\mathbf{b}_1, \dots, \mathbf{b}_n\}$ is:

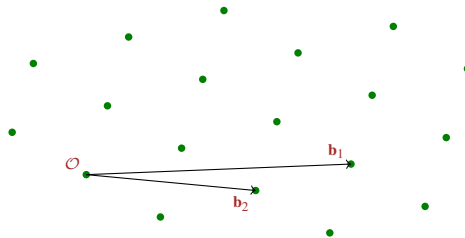
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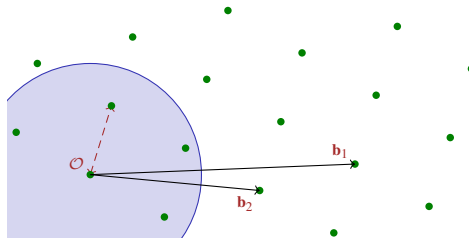
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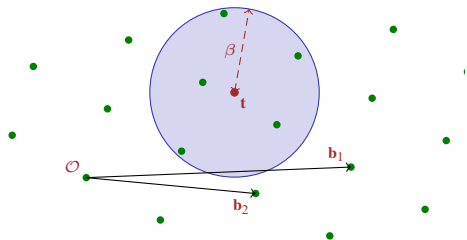
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- ▶ Given \mathbf{B} , find (nonzero) $\mathbf{v} \in \mathcal{L}$ within γ factor of shortest.

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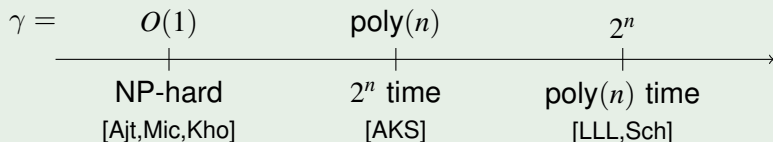
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Absolute Distance Decoding (ADD $_{\beta}$)

- ▶ Given \mathbf{B} and target $\mathbf{t} \in \mathbb{R}^n$, find some $\mathbf{v} \in \mathcal{L}$ within distance β .

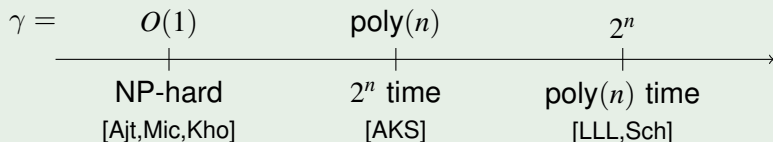
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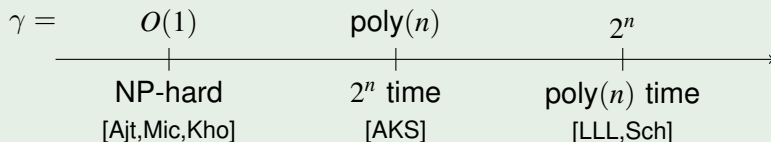
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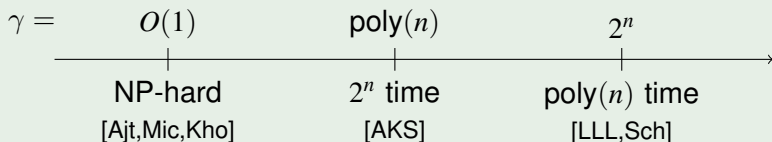
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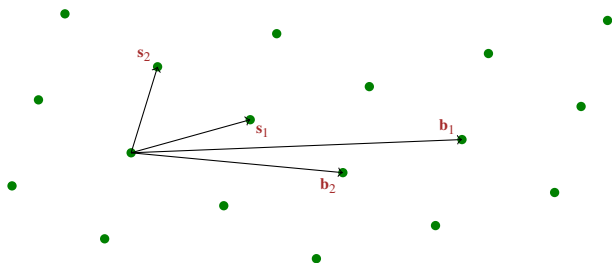
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Bottom Line

- ▶ On **random** lattices, SVP_{γ} and ADD_{β} seem **exponentially hard**

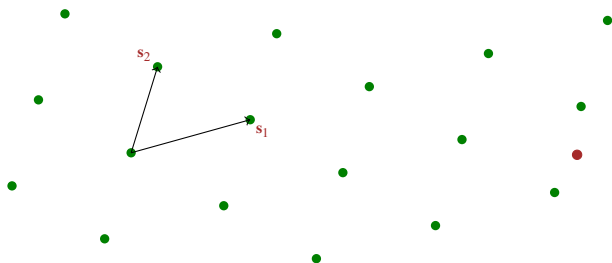
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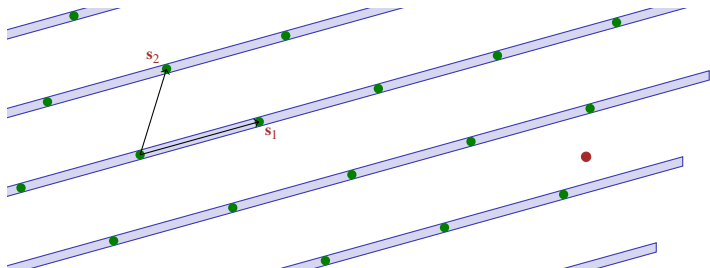
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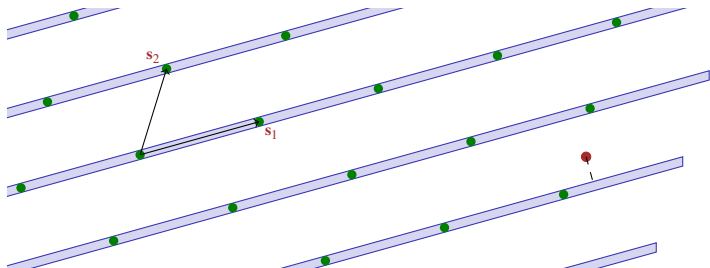
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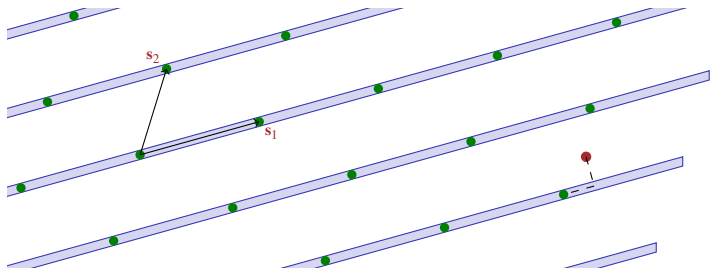
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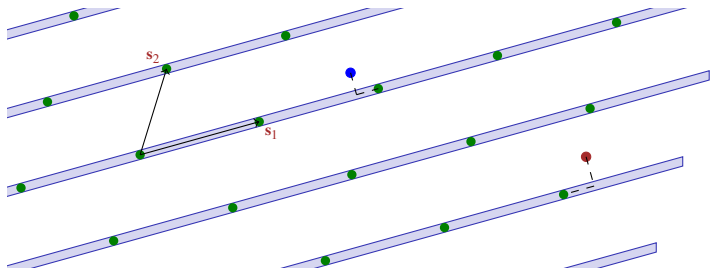
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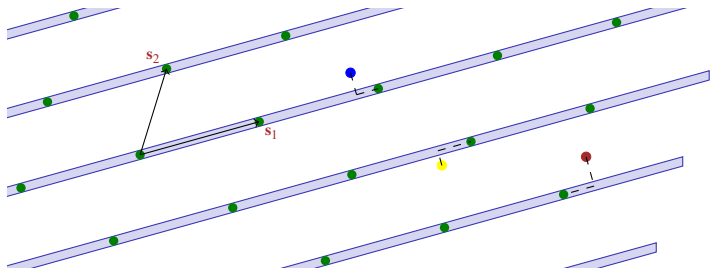
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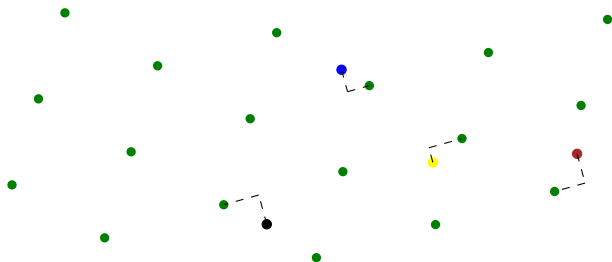
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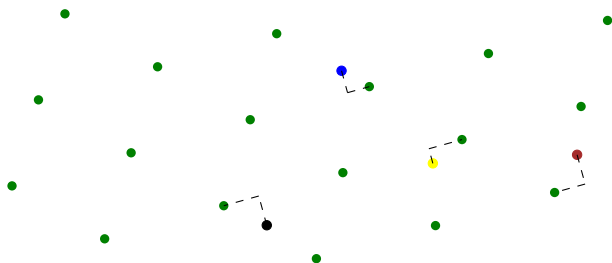
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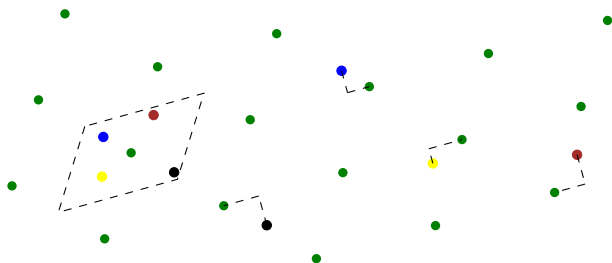


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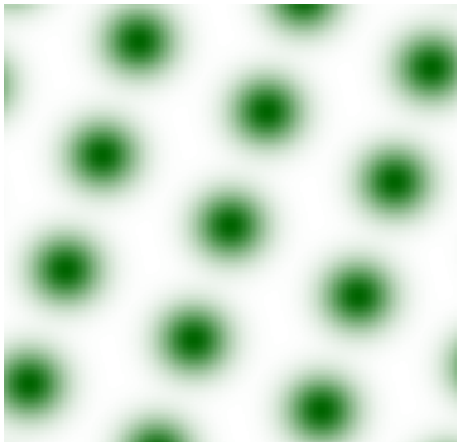
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- 2 Secret key leakage
 - **Total break** after several signatures [NguyenRegev]

Gaussians and Lattices



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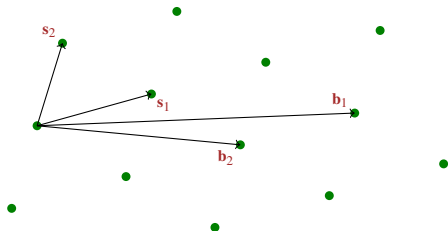


“Uniform” in \mathbb{R}^n when $\text{std dev} \geq \text{shortest basis}$

[Regev, MicciancioRegev]

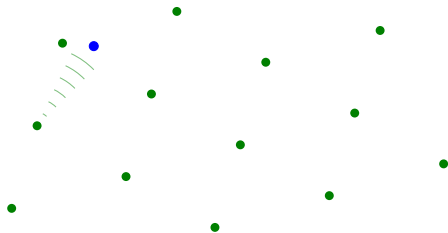
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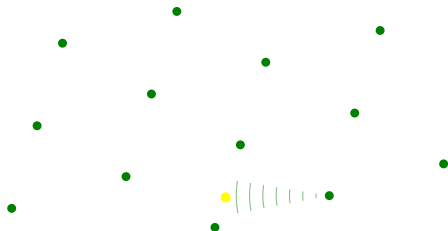
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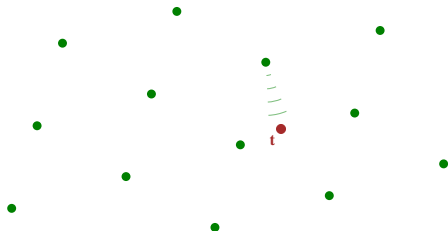
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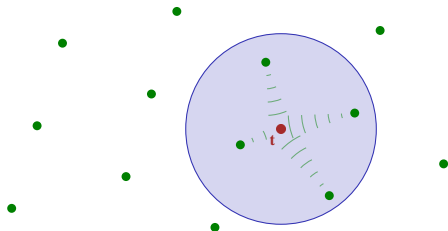
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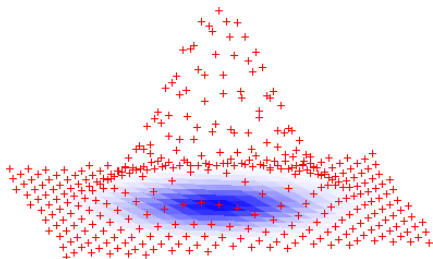
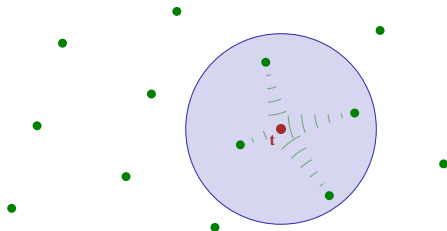
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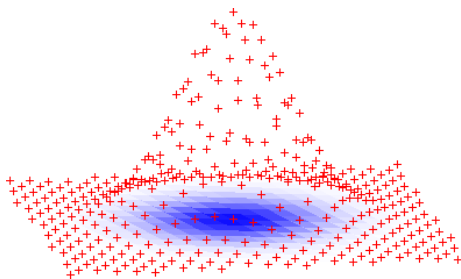
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- ▶ Conditional distribution is “discrete Gaussian” $D_{\mathcal{L}, \mathbf{t}}$



Analysis tool in
[Ban, AR, Reg, MR, Pei, . . .]

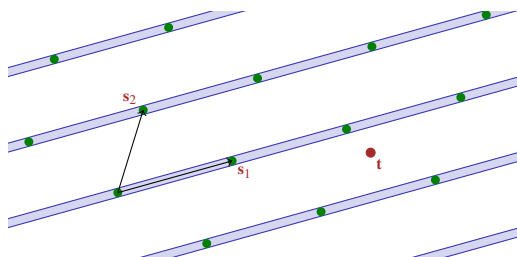
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- ▶ Given basis S , **samples** $D_{\mathcal{L},t}$ for any std dev $\geq \max\|s_i\|$
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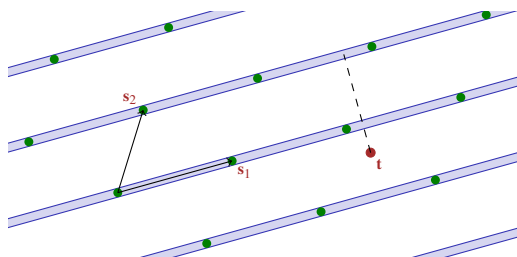
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- ▶ **Randomized** nearest-plane [Babai,Klein]



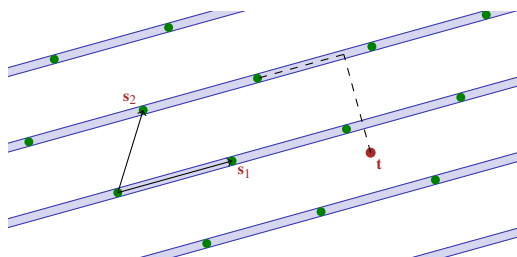
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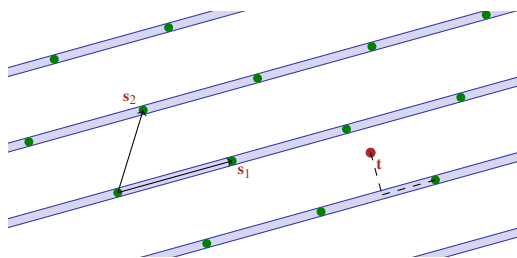
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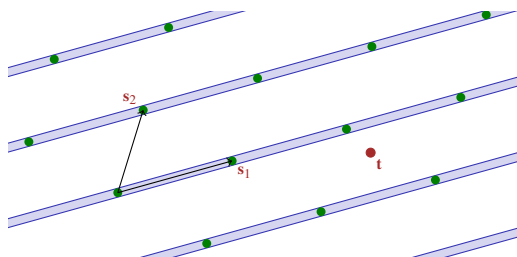
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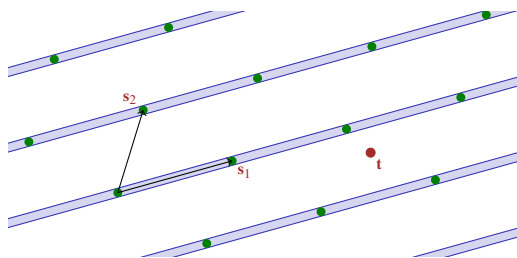
- ▶ Given basis S , samples $D_{\mathcal{L},t}$ for any std dev $\geq \max \|s_i\|$
 - Leaks nothing about S !
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[This work]: std dev $\geq \max \|\tilde{s}_i\| \Rightarrow$ samples $D_{\mathcal{L},t}$ exactly*

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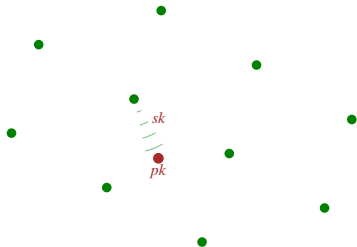
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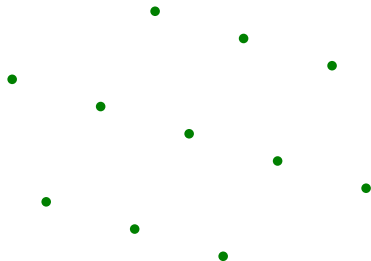
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Cryptosystem with Master Trapdoor

Primal \mathcal{L}

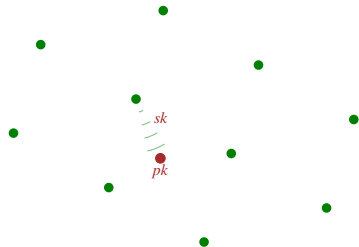


Dual \mathcal{L}^*

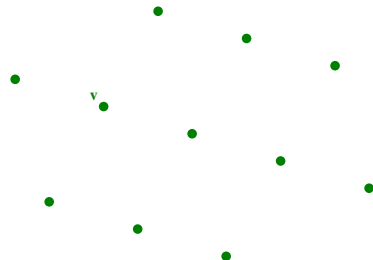


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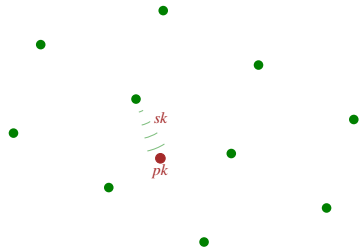
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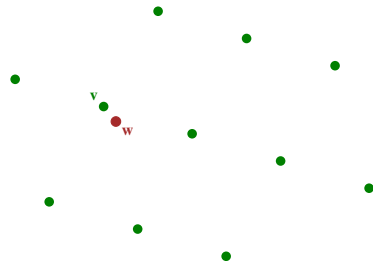
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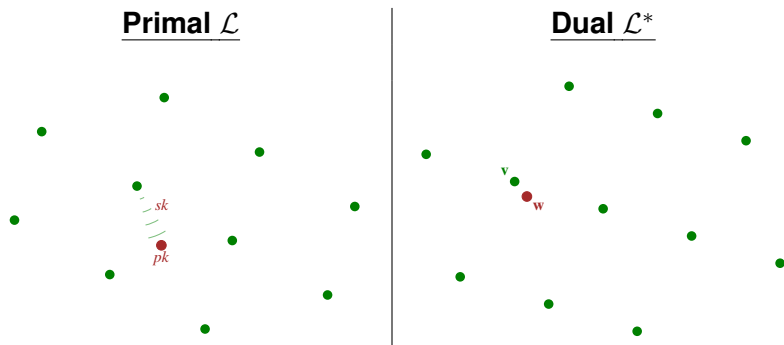


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“quasi”-agreement

Cryptosystem with Master Trapdoor



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- ▶ For $\mathbf{w} \approx \mathbf{v}$: $\langle \mathbf{v}, pk \rangle \approx \langle \mathbf{w}, sk \rangle \pmod{1}$ “quasi”-agreement
- ▶ Security: decoding \mathbf{w} , a.k.a. “learning with errors”
 - Quantum worst-case connection [Regev]
 - Now: classical worst-case hardness [P]

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Thanks!



(Artwork courtesy of xkcd.org)