

## Programming Languages Topic of Ultimate Mastery

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CS 615 - TR 5:00-6:15 @ OLS 005

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(note: CS 615 == CS 655)

## Reasonable Initial Skepticism



## Today's Class

- Vague Historical Context
- Goals For This Course
- Requirements and Grading
- Course Summary
- Convince you that PL is useful

## Meta-Level Information

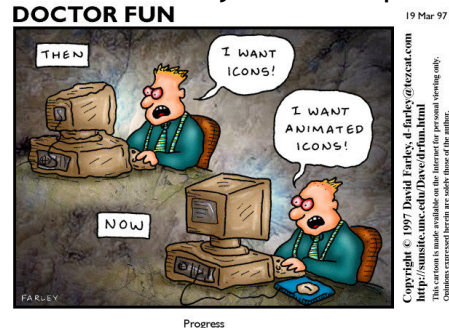
- Please interrupt at any time!
- Completely cromulent queries:
  - I don't understand: please say it another way.
  - Slow down, you talk too fast!
  - Wait, I want to read that!
  - I didn't get joke X, please explain.

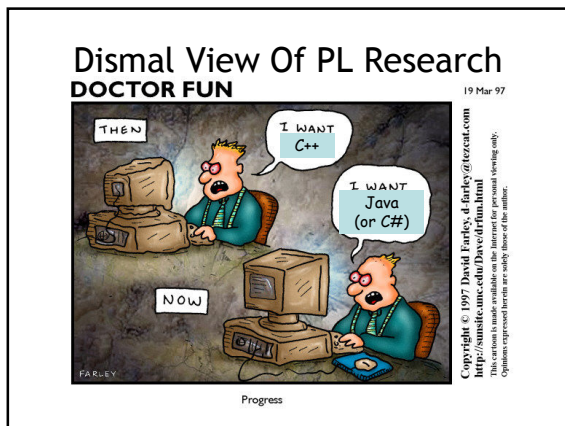


## What Have You Done For Us Lately?

- Isn't PL a solved problem?
    - PL is an old field within Computer Science
    - 1920's: "computer" = "person"
    - 1936: Church's Lambda Calculus (= PL!)
    - 1937: Shannon's digital circuit design
    - 1940's: first digital computers
    - 1950's: FORTRAN (= PL!)
    - 1958: LISP (= PL!)
    - 1960's: Unix
    - 1972: C Programming Language
    - 1981: TCP/IP
    - 1985: Microsoft Windows
    - 1992: Ultima Underworld / Wolfenstein 3D
- “... a prestigious line of work with a long and glorious tradition.” - Vizzini

## Don't We Already Have Compilers?





### Parts of Computer Science

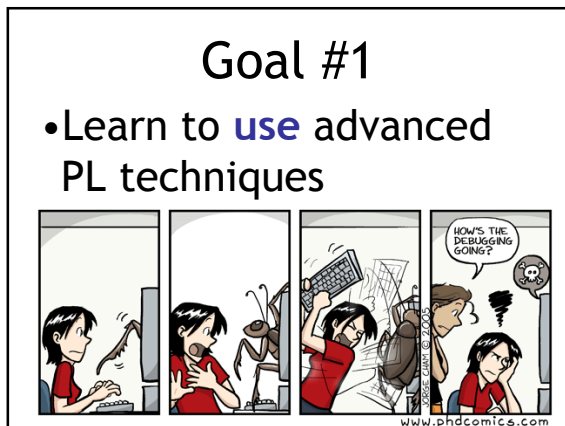
- CS = (Math × Logic) + Engineering
  - Science (from Latin *scientia* - knowledge) refers to a system of acquiring knowledge - based on empiricism, experimentation, and methodological naturalism - aimed at finding out the truth.
- We rarely actually do this in CS
  - “CS theory” = Math (logic)
  - “Systems” = Engineering (bridge building)

### Programming Languages

- Best of both worlds: **Theory** and **Practice!**
  - Only pure CS theory is more primal
- Touches most other CS areas
  - **Theory**: DFAs, PDAs, TMs, language theory (e.g., LALR)
  - **Systems**: system calls, assembler, memory management
  - **Arch**: compiler targets, optimizations, stack frames
  - **Numerics**: FORTRAN, IEEE FP, Matlab
  - **AI**: theorem proving, ML, search
  - **DB**: SQL, persistent objects, modern linkers
  - **Networking**: packet filters, protocols, even Ruby on Rails
  - **Graphics**: OpenGL, LaTeX, PostScript, even Logo (= LISP)
  - **Security**: buffer overruns, .net, bytecode, PCC, ...
  - **Software Engineering**: obvious

### Overarching Theme

- I assert (**and shall convince you**) that
- PL is one of the most **vibrant** and **active** areas of CS research today
  - It has theoretical and practical meatiness
  - It intersects most other CS areas
- You will be able to use PL techniques in **your own projects**



### Useful Complex Knowledge

- A proof of the fundamental theorem of calculus
- A proof of the max-flow min-cut theorem
- Nifty Tree node insertion (e.g., B-Trees, AVL, Red-Black)
- The code for the Fast Fourier Transform
- And so on ...

## No Useless Memorization

- I will not waste your time with useless memorization
- This course will cover complex subjects
- I will teach their details to help you understand them the first time
- But you will never have to memorize anything low-level
- Rather, learn to apply broad concepts

## Goal #2

- When (not if) you design a language, it will avoid the mistakes of the past and you'll be able to describe it formally

## Story: The Clash of Two Features

- Real story about bad programming language design
- Cast includes famous scientists
- ML ('82) is a functional language with polymorphism and monomorphic references (i.e. pointers)
- Standard ML ('85) innovates by adding polymorphic reference
- It took 10 years to fix the "innovation"

## Polymorphism (Informal)

- Code that works uniformly on various types of data
- Examples:
  - `length :  $\alpha$  list  $\rightarrow$  int` (takes an argument of type "list of  $\alpha$ ", returns an integer, for any type  $\alpha$ )
  - `head :  $\alpha$  list  $\rightarrow$   $\alpha$`
- Type inference:
  - generalize all elements of the input type that are not used by the computation

## References in Standard ML

- Like "updatable pointers" in C
- Type constructor: `ptr  $\tau$`
- Expressions:
  - `alloc :  $\tau \rightarrow$  ptr  $\tau$`  (allocate a cell to store a  $\tau$ )
  - `*e :  $\tau$`  when `e : ptr  $\tau$`  (read through a pointer)
  - `*e := e'` with `e : ptr  $\tau$`  and `e' :  $\tau$`   
(write through a pointer)
- Works just as you might expect

## Polymorphic References: A Major Pain

Consider the following program fragment:

Code	Type inference
<code>fun id(x) = x</code>	<code>id : <math>\alpha \rightarrow \alpha</math></code> (for any $\alpha$ )
<code>val c = alloc id</code>	<code>c : ptr (<math>\alpha \rightarrow \alpha</math>)</code> (for any $\alpha$ )
<code>fun inc(x) = x + 1</code>	<code>inc : int <math>\rightarrow</math> int</code>
<code>*c := inc</code>	Ok, since <code>c : ptr (int <math>\rightarrow</math> int)</code>
<code>(*c) ("hi")</code>	Ok, <code>c : ptr (string <math>\rightarrow</math> string)</code>

## Reconciling Polymorphism and References

- Type system **fails to prevent a type error!**
- Common solution:
  - value restriction: generalize only the type of values!
    - easy to use, simple proof of soundness
- **X Features  $\Rightarrow$  X<sup>2</sup> Complication**
- To see what went wrong we needed to understand semantics, type systems, polymorphism and references

## Story: Java Bytecode Subroutines

- **Java bytecode** programs contain **subroutines** (jsr) that run in the caller's stack frame (*why?*)
- jsr complicates the formal semantics of bytecodes
  - Several verifier bugs were in code implementing jsr
  - 30% of typing rules, 50% of soundness proof due to jsr
- It is **not worth it:**
  - In 650K lines of Java code, 230 subroutines, saving 2427 bytes, or 0.02%
  - 13 times more space could be saved by renaming the language back to Oak
    - [In 1994], the language was renamed "Java" after a trademark search revealed that the name "Oak" was used by a manufacturer of video adapter cards.

## Recall Goal #2

- When (not if) you **design** a language, it will avoid the mistakes of the past and you'll be able to describe it formally

## Goal #3

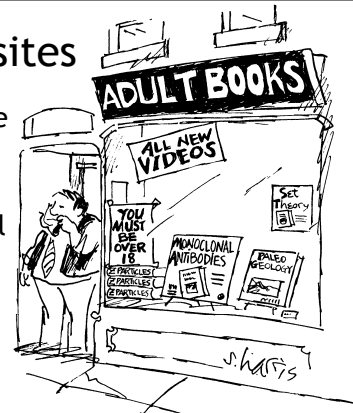
- Understand **current PL research** (PLDI, POPL, OOPSLA, TOPLAS, ...)

## Final Goal: Fun



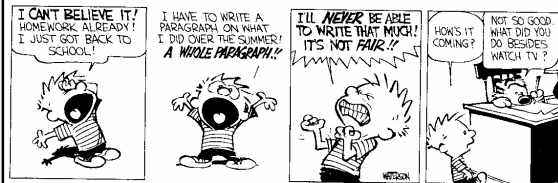
## Prerequisites

- Undergraduate compilers course
- "Mathematical maturity"



## Assignments

- Short Homework Assignments (4)
- Long Homework Assignment (1)
- Daily Reading (~2 papers per class)
- **Final Project**



## Homework Problem Sets

- Some material can be “mathy”
- Much like Calculus, practice is handy
- **Short:** ~3 theory + 1 coding per HW
- You have **one week** to do each one
- **Long:** analysis of real C programs
- NB: I will offer suggestions and comments on your **English prose**.

## Final Project

- Literature survey, implementation project, or research project
- Write a 10-page paper (a la PLDI)
- Give a 10-15 minute presentation
- On the topic of your choice
  - I will help you find a topic (many examples)
  - Best: **integrate PL with your current research**

## How Hard Is This Class?



## This Shall Be Avoided



In 1930, the Republican-controlled House of Representatives, in an effort to alleviate the effects of the... Anyone? Anyone? ... the Great Depression, passed the ... Anyone? Anyone? The tariff bill? The Hawley-Smoot Tariff Act? Which, anyone? Raised or lowered? ... raised tariffs, in an effort to collect more revenue for the federal government. Did it work? Anyone? Anyone know the effects?

## Key Features of PL



## Programs and Languages

- Programs
  - What are they trying to do?
  - Are they doing it?
  - Are they making some other mistake?
  - Were they hard to write?
  - Could we make it easier?
  - Should you run them?
  - How should you run them?
  - How can I run them faster?

## Programs and Languages

- Languages
  - Why are they annoying?
  - How could we make them better?
  - What tasks can they make easier?
  - What cool features might we add?
  - Can we stop mistakes before they happen?
  - Do we need new paradigms?
  - How can we help out My Favorite Domain?

## Common PL Research Tasks

- Design a new language feature
- Design a new type system / checker
- Design a new program analysis
- Find bugs in programs
- (Help people to) Fix bugs in programs
- Transform programs (source or assembly)
- Interpret and execute programs
- Prove things about programs
- Optimize programs

## Grand Unified Theory

- Design a new type system
- Your type-checker becomes a bug-finder
- No type errors  $\Rightarrow$  proof program is safe
- Design a new language feature
- To prevent the sort of mistakes you found
- Write a source-to-source transform
- Your new feature works on existing code

## CS 615 - Core Topics

- Operational semantics
- **Type theory**
- Verification conditions
- Abstract interpretation
- Lambda Calculus
- Type systems



"SO, BY A VOTE OF 8 TO 2 WE HAVE DECIDED TO SKIP THE INDUSTRIAL REVOLUTION COMPLETELY, AND GO RIGHT INTO THE ELECTRONIC AGE."

## Special Topics

- Object-Oriented Languages
- Software Model Checking
- Type Systems for Resource Management
- Automated Deduction / Theorem Proving
- **What do you want to hear about?**



## In Our Next Exciting Episode

$$\frac{\langle e, \sigma \rangle \Downarrow n}{\langle x := e, \sigma \rangle \Downarrow \sigma[x := n]}$$

Def:  $\sigma[x := n](x) = n$   
 $\sigma[x := n](y) = \sigma(y)$

$$\frac{\langle b, \sigma \rangle \Downarrow \text{false}}{\langle \text{while } b \text{ do } c, \sigma \rangle \Downarrow \sigma}$$

$$\frac{\langle b, \sigma \rangle \Downarrow \text{true} \quad \langle c; \text{while } b \text{ do } c, \sigma \rangle \Downarrow \sigma'}{\langle \text{while } b \text{ do } c, \sigma \rangle \Downarrow \sigma'}$$

## Meat Today? Judgments

hypotheses therefore result

$$\langle e, \sigma \rangle \Downarrow n$$

$$\{Z\} \models Z \wedge ((P \Rightarrow Q) \Rightarrow P) \Rightarrow P$$

$$\Theta; \Delta; \Sigma; \Gamma \vdash x : \Gamma(x)$$

These are examples of judgments. A judgment can really show anything. A **judgment** is something that is true (or can be proved) about or within a domain of discourse.

## Rules of Inference

Hypothesis<sub>1</sub> ... Hypothesis<sub>N</sub>

Conclusion

$$\frac{\Gamma \vdash b : \text{bool} \quad \Gamma \vdash e1 : \tau \quad \Gamma \vdash e2 : \tau}{\Gamma \vdash \text{if } b \text{ then } e1 \text{ else } e2 : \tau}$$

- For any given proof system, a finite number of rules of inference (or schema) are listed somewhere
- Rule instances should be **easily checked**
- **What is the definition of “NP”?**

## Derivation

$$\frac{\frac{\Gamma(x) = \text{int}}{\Gamma \vdash x : \text{int}} \text{ var} \quad \frac{\Gamma(x) = \text{int}}{\Gamma \vdash 3 : \text{int}} \text{ int}}{\Gamma \vdash x > 3 : \text{bool}} \text{ gt} \quad \frac{\frac{\Gamma(x) = \text{int}}{\Gamma \vdash x : \text{int}} \text{ var} \quad \frac{\Gamma(x) = \text{int}}{\Gamma \vdash x - 1 : \text{int}} \text{ sub}}{\Gamma \vdash x := x - 1} \text{ assign}}{\Gamma \vdash \text{while } x > 3 \text{ do } x := x - 1 \text{ done}} \text{ while}$$

- **Tree-structured** (conclusion at bottom)
- May include multiple sorts of rules-of-inference
- Could be constructed, typically are not (we just prove you could construct them)

## For Our Next Exciting Episode

- See webpage under “Lectures”
- [Read Winskel Chapter 2](#)
- [Read Hoare article](#)
- Peruse the optional readings

