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## ADVANCED COMPUTER NETWORKS

Roy, A., et al., "Inside the Social Network's (Datacenter) Network" *Proc. of ACM SIGCOMM '15*, 45(4):123-137, Oct. 2015

## Microsoft Datacenter Traffic

Previous Microsoft studies found datacenter traffic to be:

- 50-80% rack local
- frequently concentrated and bursty
- bimodal in packet sizes (ACK/MTU)
- on/off
- mostly in small flows, <5 concurrent large flows

## Facebook Datacenter Traffic

Characteristics of Facebook datacenter traffic:

- neither rack local nor all pairs
- demand is wide-spread, uniform, and stable [due to load balancing](#)
- small packets, continuous arrivals, [not on/off](#)
- many concurrent flows [due to connection pooling](#)
- rapidly changing, internally bursty heavy hitters, [reducing the efficacy of traffic engineering](#)
- only Hadoop's MapReduce-style traffic agrees with Microsoft's characterization

## Implications

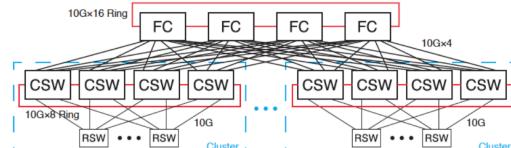
Datacenter network designs assume:

- worst-case, all-pair traffic matrix, with equal frequency and intensity ⇒ [maximize bisection bandwidth](#)
- hot-spots, due to oversubscription, to be alleviated with [bypass, secondary connectivities](#) (wireless, optical)
  - which requires traffic demand to be [predictable and stable](#) to be feasible
- stylized traffic allows for [specialized switch design](#) (buffer sizing, port count, etc.)

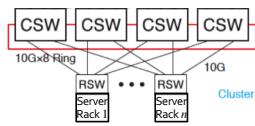
# Datacenter Topology

Similar to Google's first gen network:

- multiple **sites** connected by a backbone
- each site contains one or more buildings (**datacenters**)
- each datacenter contains multiple **clusters**



- each **cluster** employs a 3-tier, 4-post topology
- 10-Gbps **servers**



# Cluster

Unit of deployment

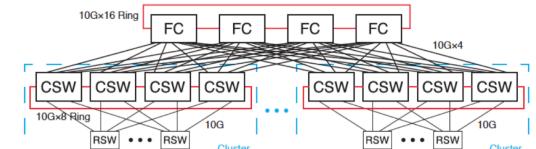
May be of a single function, e.g., cache cluster

Or multi-function: front-end cluster comprising web/front-end servers, load balancers, and cache servers

Inter-cluster, intra-datacenter connected by FC switches

Similar to Google,

- inter-datacenter, intra-site connected by aggregation switch
- inter-site connected by datacenter router



# Server

Each server has precisely one role:

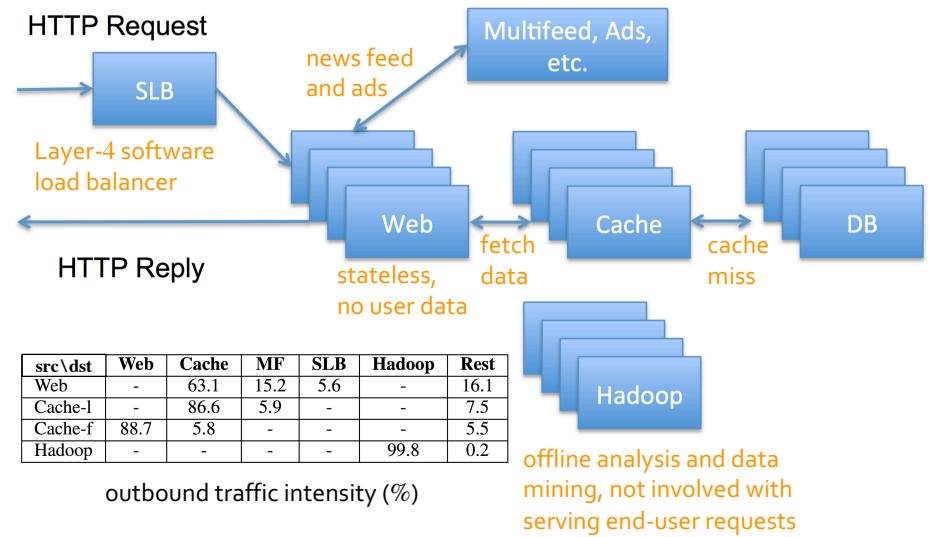
- web/front-end server
- mysql (db) server
- cache leader
- cache follower
- multifeed server to assemble news feed and serve ads
- Hadoop server for offline analysis and data mining

A small number of servers can be dynamically repurposed

No virtual machines (same as Microsoft)

Each rack contains only servers of the same role

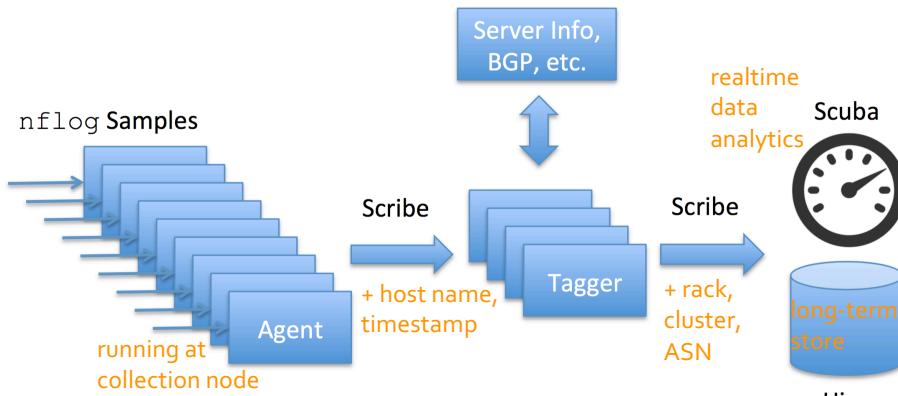
# Services



# Data Collection

Cannot collect every packet, instead use:

1. [Fbflow](#): sample packet headers (1:30K sampling rate) across entire global network



# Data Collection

Cannot collect every packet, instead use:

2. [Port mirroring](#): collect all packet headers of a [single machine](#) or [rack](#) for a few minutes
  - by [mirroring a ToR port](#) to a collection host on the same rack
    - [placement opportunistic](#), depending on space availability
    - a [kernel module](#) sitting atop the Ethernet driver extracts headers and spools it to [remote storage](#)
  - no loss
  - deployed at 5 different (type of) racks to monitor:
    - a [rack](#) of web servers
    - a Hadoop node
    - a cache leader node
    - a cache follower node
    - a multifeed node

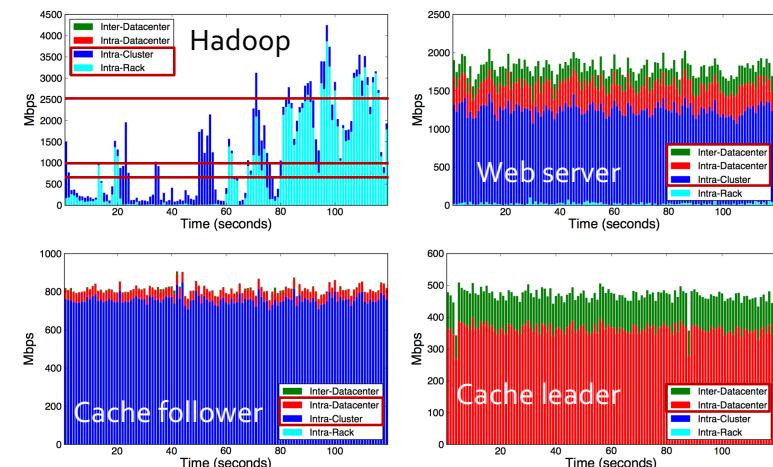
# Traffic Characterization

Characterize traffic across 3 different types of cluster:  
Hadoop, Web/front-end, and cache clusters

## Utilization:

- server to ToR links: < 1%,  
even in heaviest utilized Hadoop cluster, it's < 5%
- ToR to CSW links: median: 10-20%,  
with the busiest 5% reaching 23-46%
- CSW to FC links: higher

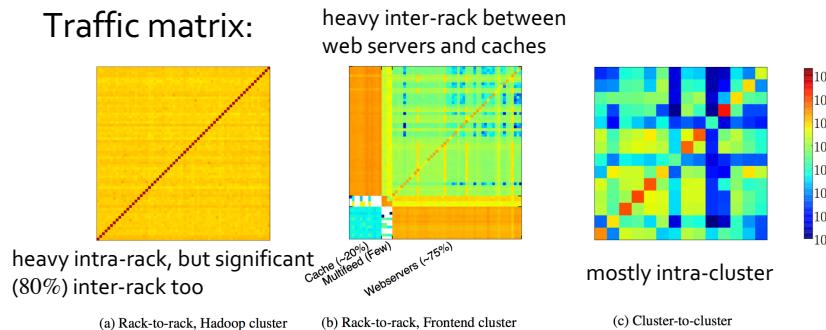
# Locality



Relative proportions of the locality are stable despite diurnal traffic pattern

# Implications of Locality

Traffic matrix:

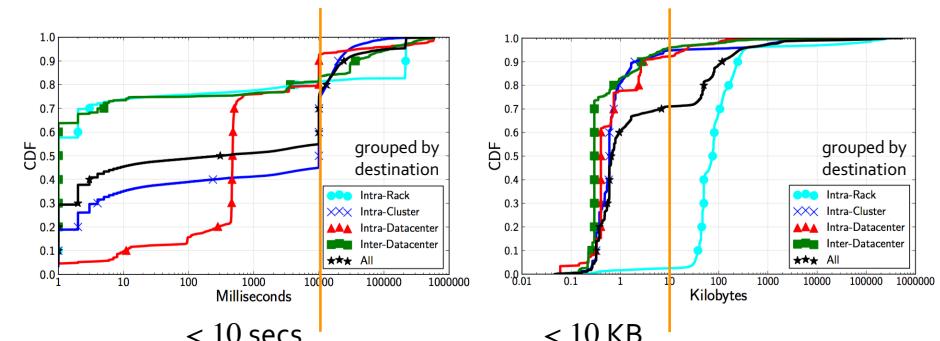


Homogenous topology will lead to over-/under-provisioning in different parts of the datacenter

Stability of traffic patterns means no need for rapid reconfigurability

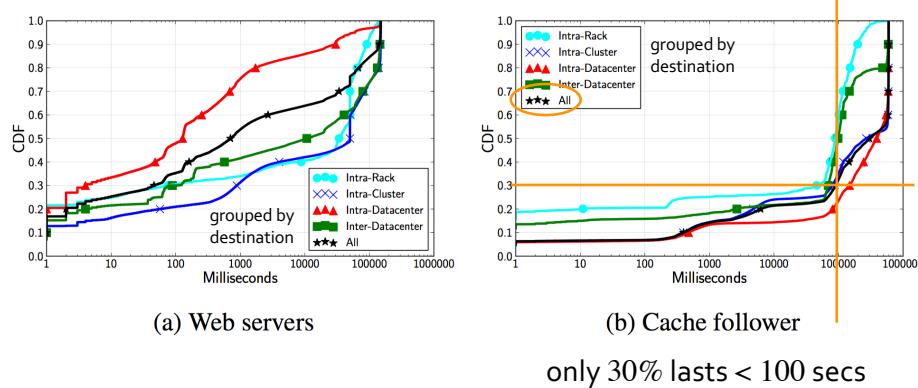
# Outbound Flow Characteristics

Most Hadoop flows are short and small, but varies across servers



# Outbound Flow Characteristics

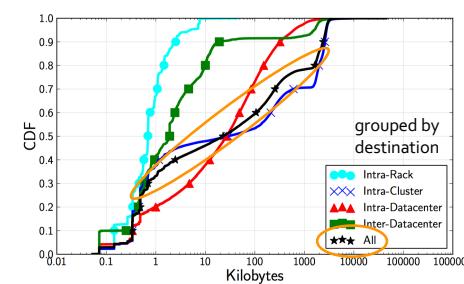
Non-Hadoop flows are more uniform across servers due to [load balancing](#) and last longer due to [connection pooling](#), but traffic per flow is bursty [surely on/off?]



# Outbound Flow Characteristics

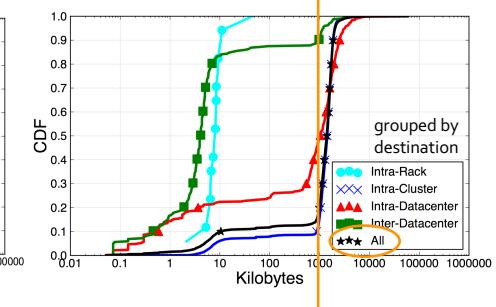
Cache flow sizes reflect [load balancing](#) over time

individual connection sizes



widely distributed

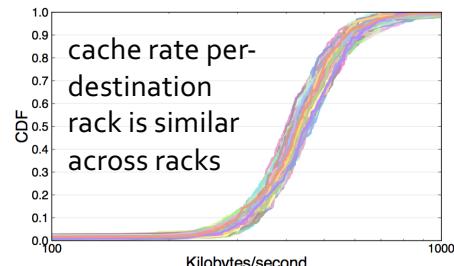
per-destination host flow sizes



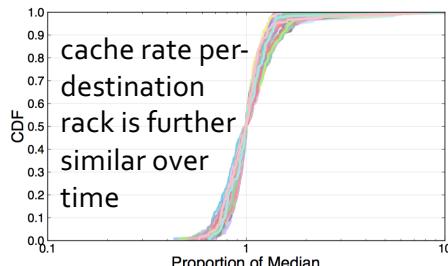
tightly distributed around 1 MB

# Impact of Load Balancing

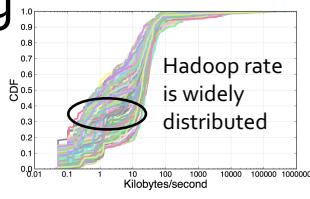
Load balancing smooths out traffic, reducing effectiveness of traffic engineering



(b) Cache (rate, each series is 1 second)



(c) Cache (stability, each series is a rack)



# Impact of Load Balancing

Load is monitored  $\Rightarrow$  large increases in load would be actively mitigated

Hot objects are temporarily cached at the web servers

Persistently popular objects are replicated across caches

Top-50 most popular objects are evenly spread across all caches

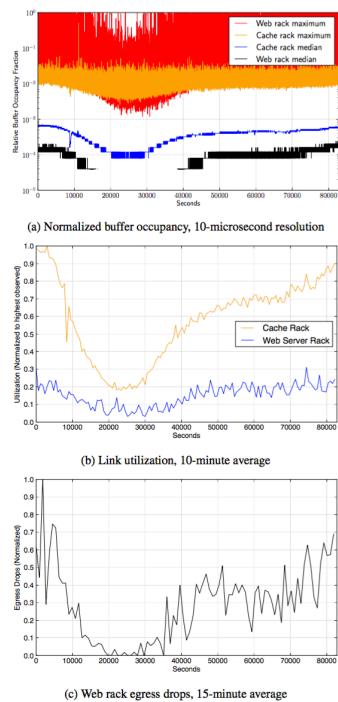
No heavy hitters (set of flows responsible for 50% of traffic volume) due to load balancing and caching

# Switch Design

Low traffic volume (in bytes), but high packet rate: even at 10% utilization, median packet size of 175 bytes means 85% of link packet forwarding capacity [no Nagle?]

Packet arrivals from a single source host are not ON/OFF, but arrivals for a single destination host are ON/OFF

Buffers overflow, especially for web servers [LRD traffic after all?]



# Discussion

Traffic observed reflects the design and implementation of a single service, is it the best design and implementation?

Traffic characteristics change as:

- service changes, e.g., more videos
- implementation or design changes, e.g., is having a cache cluster the best design?
  - or would it be better to spread cache servers across clusters?

Can all datacenter traffic be so regularized?

If so, are remaining datacenter hard problems (research issues) above the network layer?