Information-Theoretic Results on Communication Problems with Feed-forward and Feedback

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Thesis Outline

Source Coding with Feed-forward

- Point to Point Source Coding (Chap. 2)
 - Problem Statement
 - Rate-Distortion Function
 - Error Exponents
- Computation (Chap. 4)
- Multiple Descriptions (Chap. 5)

Ochannel Coding with Feedback (Chap. 3,4)

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Ochannel Coding with Feedback (Chap. 3,4)



- Source X, reconstruction \hat{X}
- Distortion measure $d(X, \hat{X})$

Rate-distortion function

Minimum rate R for distortion level D (Shannon '59)



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- Example: Temperature at nearby cities
- Presence of $Y \Rightarrow$ lower rate for given distortion D

Rate-distortion function $R_{WZ}(D)$ [Wyner,Ziv '76]



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The source field itself available with delay at decoder.



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Source Coding with Feed-Forward

 Feed-forward⇒ Decoder knows some of the past source samples.



FF with delay k, block length N.

To produce X̂_n, the decoder knows index W and (X₁,..., X_{n-k}).

Prediction and FF

Stock Market Example

- Behavior of a particular stock over an *N*-day period.
- Stock price- modeled as a k + 1-state Markov chain.
- Value of stock: Markov source $\{X_n\}$



Insider and Investor



- Insider- a priori knowledge about behavior of stock
- Investor has stock for *N* days, needs to know when value drops.
- Insider: gives information to investor at rate R.

Insider and Investor



Reconstruction

- Decision of investor on day $n: \hat{X}_n$
- $\hat{X}_n = 1 \Rightarrow$ price drop from day n 1 to n
- Otherwise $\hat{X}_n = 0$
- Distortion = 1 if \hat{X}_n is wrong

Insider and Investor



Reconstruction

- Decision of investor on day $n: \hat{X}_n$
- $\hat{X}_n = 1 \Rightarrow$ price drop from day n-1 to n
- Otherwise $\hat{X}_n = 0$
- Distortion = 1 if \hat{X}_n is wrong



Feed-forward!

- Before day n, investor knows previous stock values Xⁿ⁻¹, makes decision X̂_n
- Minimum info (in bits/sample) needed to predict drops with distortion *D*?

Feed-Forward: A Formal Definition

Previous: [Weissman et al 03], [Pradhan 04], [Martinian et al 04]



- Source X: Alphabet \mathcal{X} , reconstruction alphabet $\widehat{\mathcal{X}}$
- Encoder: Rate R , $e: \mathcal{X}^N \to \{1, \dots, 2^{NR}\}$
- **Decoder**: knows all the past (n k) source samples to reconstruct *n*th sample.

$$g_n: \{1,\ldots,2^{NR}\} imes \mathcal{X}^{n-k} o \widehat{\mathcal{X}}, \quad n=1,\ldots,N.$$

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A Formal Definition (contd.)

• Distortion measure $d_N(X^N, \hat{X}^N)$.

GOAL

Given any source X, find the least R such that

 $E[d_N(x^N, \hat{x}^N)] \leq D.$

• Rate-Distortion function with Feed-forward!

Directed Information

- Inspired by [Marko '73]: work on bidirectional communication
- [Massey '90] The directed information flowing from A^N to B^N

$$I(A^N \to B^N) = \sum_{n=1}^N I(A^n; B_n | B^{n-1}).$$

• Interestingly :

$$I(A^N \to B^N) = I(A^N; B^N) - \sum_{n=2}^N I(B^{n-1}; A_n | A^{n-1})$$

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Causality in Information Flow

$$I(A^N; B^N) = I(A^N \to B^N) + I(0B^{N-1} \to A^N)$$

- $I(A^N \rightarrow B^N)$: How causal knowledge of A_n 's reduces the uncertainty in B_n
- Example: GNP vs Money Supply [Geweke '82]



Interpretation of Directed Information



Without FF, need $I(\hat{X}^N; X^N)$ bits to represent X^N with $\hat{X^N}$.

• With feed-forward, to produce X_n , the decoder knows X^{n-1}

Number of bits required is reduced by $I(\hat{X}_n; X^{n-1} | \hat{X}^{n-1})$.

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• With delay 1 feed-forward, we need

$$I(\hat{X}^{N}; X^{N}) - \sum_{n=2}^{N} I(\hat{X}_{n}; X^{n-1} | \hat{X}^{n-1})$$
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• Directed information from \hat{X}^N to X^N !

Delay k feed-forward



- With delay k feed-forward, to produce \hat{X}_n , the decoder knows X^{n-k} .
- No. of bits: $I(\hat{X}^{N}; X^{N}) \sum_{n=k+1}^{N} I(\hat{X}_{n}; X^{n-k} | \hat{X}^{n-1})$
- Not Directed Information- will denote it $I_k(\hat{X}^N \to X^N)$ - 'k-directed information'.

Rate-distortion function

- Optimize I_k over $P_{\hat{X}^N|X^N}$ that satisfy distortion
- No savings for discrete memoryless source w/ single-letter distortion measure

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- with feed-forward, optimal joint distribution may not be.

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a_1, a_2, \ldots : random sequence

• $\limsup_{n \text{ prob}} a_n = \overline{a}$: Smallest number α such that

$$\lim_{n\to\infty}\Pr(a_n>\alpha)=0.$$

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$$i_k(\hat{x}^n \to x^n) = \frac{1}{n} \log \frac{P(x^n, \hat{x}^n)}{P(x^n) \cdot \prod_{i=1}^n P(\hat{x}_i | \hat{x}^{i-1}, x^{i-k})}$$

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Rate-Distortion Theorem for a general source

•
$$\mathbf{P}_{\mathbf{X}} = \{P_{X_1}, P_{X^2}, \dots, P_{X^N}, \dots\}$$

•
$$\mathbf{P}_{\hat{\mathbf{X}}|\mathbf{X}} = \{P_{\hat{X}_1|X_1}, P_{\hat{X}^2|X^2}, \dots, P_{\hat{X}^N|X^N}, \dots\}$$

Theorem

For arbitrary source X with distribution P_X , the rate-distortion function with feed-forward delay k, the infimum of all achievable rates at probability-1 distortion D, is given by

$$\mathsf{R}_{\mathrm{ff}}(D) = \inf_{\mathsf{P}_{\hat{\mathbf{X}}|\mathbf{X}}: \rho(\mathsf{P}_{\hat{\mathbf{X}}|\mathbf{X}}) \leq D} \overline{I}_k(\hat{X} \to X),$$

where

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- Directed information and why it occurs
- Feed-forward rate-distortion result for general sources, distortions

Next . .

How to compute the rate-distortion function?

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How to compute the rate-distortion function?

Source Coding Optimization



- Source X with distribution **P**_X.
- Find $\mathbf{P}_{\hat{\mathbf{X}}|\mathbf{X}}$ to minimize $\overline{I}_k(\hat{X} \to X)$ s.t

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- Pick a conditional distribution $\mathbf{P}_{\hat{\mathbf{X}}|\mathbf{X}} = \{P_{\hat{\mathbf{X}}^n|\mathbf{X}^n}\}$
- For what sequence of distortion measures d_n does $P_{\hat{X}|X}$ achieve the infimum in the rate-distortion formula ?
- Approach- similar in spirit to [Csiszar and Korner]

Structure of Distortion Function

Theorem

A stationary, ergodic source X characterized by $\mathbf{P}_{\mathbf{X}} = \{P_{X^n}\}_{n=1}^{\infty}$ with feed-forward delay k. $\mathbf{P}_{\hat{\mathbf{X}}|\mathbf{X}} = \{P_{X^n|X^n}\}_{n=1}^{\infty}$ is a conditional distribution such that the joint distribution is stationary and ergodic. Then $\mathbf{P}_{\hat{\mathbf{X}}|\mathbf{X}}$ achieves the rate-distortion function if for all sufficiently large n, the distortion measure satisfies

$$d_n(x^n, \hat{x}^n) = -c \cdot rac{1}{n} \log rac{P_{X^n, \hat{X}^n}(x^n, \hat{x}^n)}{ec{P}^k_{\hat{X}^n|X^n}(\hat{x}^n|x^n)} + d_0(x^n),$$

where

$$\vec{P}^{k}_{\hat{X}^{n}|X^{n}}(\hat{x}^{n}|x^{n}) = \prod_{i=1}^{n} P_{\hat{X}_{i}|X^{i-k},\hat{X}^{i-1}}(\hat{x}_{i}|x^{i-k},\hat{x}^{i-1})$$

and c is any positive number and $d_0(.)$ is an arbitrary function.

Stock example revisited



• Value of the stock: Markov source $\{X_n\}$

• Decision of investor on day n: \hat{X}_n (0 or 1)

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Stock example revisited



• Value of the stock: Markov source $\{X_n\}$



• Decision of investor on day $n: \hat{X}_n$ (0 or 1)



- $R_{ff}(D)$: Minimum rate the investor needs to predict drops with distortion D.
- Try first-order Markov joint distribution
- Distortion can be cast in the required form!

Proposition

The minimum rate in (bits/sample) is

$$\sum_{i=1}^{k-1} \pi_i \Big[h(p_i, q_i, 1-p_i-q_i) - h(\epsilon, 1-\epsilon) \Big]
onumber \ + \pi_k \left(h(q_k, 1-q_k) - h(\epsilon, 1-\epsilon)
ight)$$

where

Computing Rate-distortion function with FF

- Predict' a conditional distribution
- **2** Check if distortion function can be put into required form.

Next ... A multi-terminal problem

Computing Rate-distortion function with FF

- Predict' a conditional distribution
- **2** Check if distortion function can be put into required form.

Next ... A multi-terminal problem

Multiple Descriptions



- Source X: compressed into packets
- Packets may be dropped
- Compress X into two *different* packets

Multiple Descriptions



- Compress X into two *different* packets



Multiple Descriptions



- Rate R₁ yields reconstruction X̂₁ with distortion D₁
 Rate R₂ yields reconstruction X̂₂ with distortion D₂
 - Want better quality D_0 if both packets received
 - \hat{X}_1 and \hat{X}_2 need to *refine* each other!

Rate	Distortion
R_1 bits/sample	D_1
R_2 bits/sample	D_2
$R_1 + R_2$ bits/sample	D_0

GOAL

Given i.i.d source P_X :

Find all achievable $(R_1, R_2, D_1, D_2, D_0)$

- Still an open problem
- Studied by [Cover, El Gamal], [Ahlswede], [Zhang, Berger], . .
- Best known rate-region: [Zhang, Berger '87]

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Network example with feed-forward



- Source propagates to destination with delay
- To reconstruct \hat{X}_n , decoder has packet and X^{n-k}



- Alice has an i.i.d binary source \sim Bernoulli(1/2)
- Bob and Carol: distortion d using R_B , R_C bits/sample
- Dave gets Bob's bits and Carol's bits- needs to reconstruct perfectly!

Characterize

 $r_{sum}(d) \triangleq$ Smallest sum-rate $R_B + R_C$ for distortion (d, d, 0)



Same model as before, one extra feature...

- After Carol reconstructs each sample, Alice reveals the value to her. *Feed-forward*
- Before reconstructing each sample, Carol knows past source samples

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General feed-forward model



- Encoder mappings: $e_m : \mathcal{X}^N \to \{1, \dots, 2^{NR_m}\}, \quad m = 1, 2.$
- Mappings for decoders 1 and 0:

$$g_1 : \{1, \dots, 2^{NR_1}\} \to \widehat{\mathcal{X}}_1^N$$
$$g_0 : \{1, \dots, 2^{NR_1}\} \times \{1, \dots, 2^{NR_2}\} \to \widehat{\mathcal{X}}_0^N$$

- A sequence of mappings for decoder 2: $g_{2n}: \{1, \dots, 2^{NR_2}\} \times \mathcal{X}^{n-k} \to \widehat{\mathcal{X}}_2, \quad n = 1, \dots, N.$

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Rate region w/o FF



Zhang-Berger '87

 $P_{U,\hat{X}_1,\hat{X}_2,\hat{X}_0|X}$ such that

 $\begin{aligned} & \textit{Ed}_{m}(X;\hat{X}_{m}) \leq \textit{D}_{m}, \quad m=0,1,2 \\ & \textit{R}_{1} > \textit{I}(X;\hat{X}_{1}\textit{U}) \quad \textit{R}_{2} > \textit{I}(X;\hat{X}_{2}\textit{U}) \\ & \textit{R}_{1} + \textit{R}_{2} > \textit{I}(X;\hat{X}_{1}\textit{U}) + \textit{I}(X;\hat{X}_{2}\textit{U}) + \textit{I}(X;\hat{X}_{0}|\hat{X}_{1}\hat{X}_{2}\textit{U}) \\ & \quad + \textit{I}(\hat{X}_{1};\hat{X}_{2}|\textit{X}\textit{U}) \end{aligned}$

Correlation in MD

Stringent Decoder 0 distortion \Rightarrow Need correlation in \hat{X}_1, \hat{X}_2

'Cloud' Center

- U: Cloud center of X sent to all decoders
- Rate I(U; X) each for decoder 1 and 2

Penalty Term

- Can't have $\hat{X}_1 \sim P(\hat{X}_1|XU)$ and $\hat{X}_2 \sim P(\hat{X}_2|XU)$ indep'ly
- Need to be *jointly* distributed: $\sim P(\hat{X}_1, \hat{X}_2 | XU)$
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Coding Strategy

• Consider *B* long blocks of source samples

 $X_1 \ldots X_N \quad X_{N+1} \ldots X_{2N} \qquad \ldots \ldots X_{NB}$

• While encoding one block, give 'preview' of next block



Restricted encoding for user 1- within cell j



After reconstructing block b

- User 1 gets 'preview' of block b+1
- User 2 knows it too- due to FF!

Block-Markov, superposition; [Cover,Leung], [Willems] for MAC

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Rate region with FF



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$$\begin{split} & \textit{Ed}_m(X; \hat{X}_m) \leq D_m, \quad m = 0, 1, 2 \\ & R_1 > I(X; \hat{X}_1 U) \\ & R_2 > I(X; \hat{X}_2 | U) + \max\{0, R_1 - I(X \hat{X}_2; \hat{X}_1 | U)\} \\ & R_1 + R_2 > I(X; \hat{X}_1 U) + I(X; \hat{X}_2 | U) + I(X; \hat{X}_0 | \hat{X}_1 \hat{X}_2 U) \\ & + I(\hat{X}_1; \hat{X}_2 | X U) + \max\{0, R_1 - I(X \hat{X}_2; \hat{X}_1 | U)\} \end{split}$$

Improvement over [Zhang-Berger]



• Fix
$$R_1 = I(X; \hat{X}_1 U) + \epsilon$$

• If max
$$\{0, R_1 - I(X\hat{X}_2; \hat{X}_1|U)\} = 0$$
:

- Savings in $R_2 = I(U; X)$ bits/sample
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Example: No feed-forward



Without FF [Zhang Berger '87]

$$r_{sum}(d) \geq 2-h\left(\frac{4d+1-\sqrt{12d^2-4d+1}}{2}\right)$$

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• (a): Lower bound without feed-forward [Zhang, Berger '87]

(b): Achievable sum-rate with FF- better than optimal w/o FF

• (c): Rate-dist lower bound with FF- $R_B + R_C > 2(1 - h(d))$

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Summary

- Feed-forward: helps build correlation
- Single-letter achievable region for MD with FF
- FF helps even for an i.i.d source
- How to use feed-forward to all decoders?
- FF for Gaussian multiple descriptions [Pradhan IT '07]

Summary

- Source Coding with feed-forward
 - -Directed Information
 - -Rate-Distortion Function
 - -How to evaluate optimization
 - -Multiple Descriptions with FF
- Channel Coding with feedback

Some questions...

- Feedback/FF in multi-terminal setting
- Noisy feedback/FF
- Applications of directed-info like quantities

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