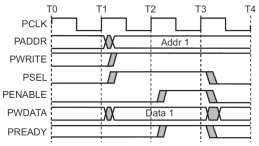
EECS 373
Design of Microprocessor-Based Systems

Branden Ghena
University of Michigan



Lecture 5: Memory and Peripheral Buses
January 22, 2015

Announcements

- I'm not Prabal
  - You probably noticed
  - He's in Washington DC doing this and this
  - He regrets that he cannot be here today and will not have OH
- We're working on additional GSI/IA office hours (OH)
  - Pat Pannuto 10-11am MW in EECS Learning Center
    - (Glass rooms between BBB and Dow)

Outline

- Announcements
- Accessing memory from Assembly/C
- Busses: The glue that connects the pieces
- ARM Advanced Peripheral Bus (APB)
- ARM Advanced High-performance Bus Light (AHB-Lite)

Accessing memory locations from C

- Memory has an address and value
- Can equate a pointer to desired address
- Can set/get de-referenced value to change memory

```
#define  SYSREG_SOFT_RST_CR  0xE0042030

uint32_t *reg = (uint32_t *)(SYSREG_SOFT_RST_CR);

main () {
  *reg |= 0x00004000; // Reset GPIO hardware
  *reg &= ~(0x00004000);
}
```


What happens when this "instruction" executes?

```
#include <stdio.h>
#include <inttypes.h>

#define REG_FOO 0x40000140

main () {
  uint32_t *reg = (uint32_t *)(REG_FOO);
  *reg += 3;

  printf("0x%x\n", *reg); // Prints out new value
}
```


"*reg += 3" is turned into a ld, add, str sequence

- Load instruction
  - A bus read operation commences
  - The CPU drives the address "reg" onto the address bus
  - The CPU indicated a read operation is in process (e.g. R/W#)
  - Some "handshaking" occurs
  - The target drives the contents of "reg" onto the data lines
  - The contents of "reg" is loaded into a CPU register (e.g. r0)
- Add instruction
  - An immediate add (e.g. add r0, #3) adds three to this value
- Store instruction
  - A bus write operation commences
  - The CPU drives the address "reg" onto the address bus
  - The CPU indicated a write operation is in process (e.g. R/W#)
  - Some "handshaking" occurs
  - The CPU drives the contents of "r0" onto the data lines
  - The target stores the data value into address "reg"

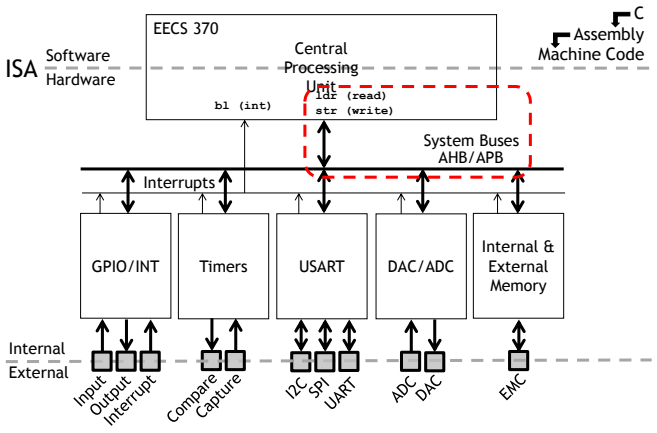
Some useful C keywords

- const
  - Makes variable value or pointer parameter unmodifiable
  - const foo = 32;
- register
  - Tells compiler to locate variables in a CPU register if possible
  - register int x;
- static
  - Preserve variable value after its scope ends
  - Does not go on the stack
  - static int x;
- volatile
  - Opposite of const
  - Can be changed in the background
  - volatile int I;

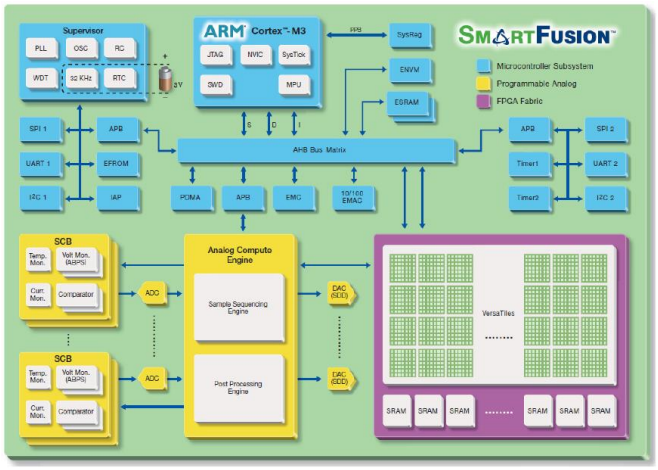
Outline

- Announcements
- Accessing memory from Assembly/C
- Busses: The glue that connects the pieces
- ARM Advanced Peripheral Bus (APB)
- ARM Advanced High-performance Bus Light (AHB-Lite)

Busses: the glue that connects the pieces



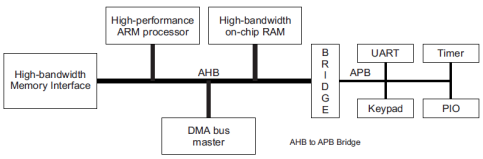
Actel SmartFusion system/bus architecture



Why have so many busses?

- Many designs considerations
  - Master vs Slave
  - Internal vs External
  - Bridged vs Flat
  - Memory vs Peripheral
  - Synchronous vs Asynchronous
  - High-speed vs low-speed
  - Serial vs Parallel
  - Single master vs multi master
  - Single layer vs multi layer
  - Multiplexed A/D vs demultiplexed A/D
- Discussion: what are some of the tradeoffs?

Advanced Microcontroller Bus Architecture (AMBA)
- Advanced High-performance Bus (AHB)
- Advanced Peripheral Bus (APB)



AHB

- High performance
- Pipelined operation
- Burst transfers
- Multiple bus masters
- Split transactions

APB

- Low power
- Latched address/control
- Simple interface
- Suitable of many peripherals

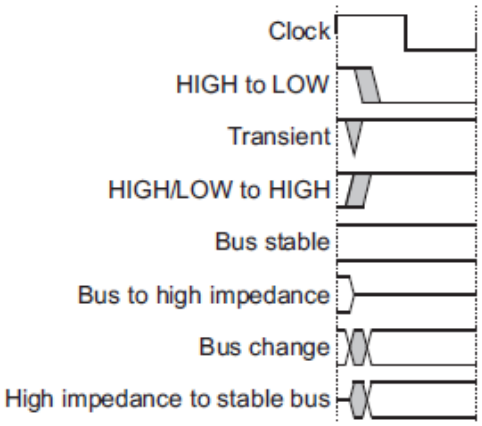
Outline

- Announcements
- Accessing memory from Assembly/C
- Busses: The glue that connects the pieces
- ARM Advanced Peripheral Bus (APB)
- ARM Advanced High-performance Bus Light (AHB-Lite)

APB: a simple bus that is easy to work with

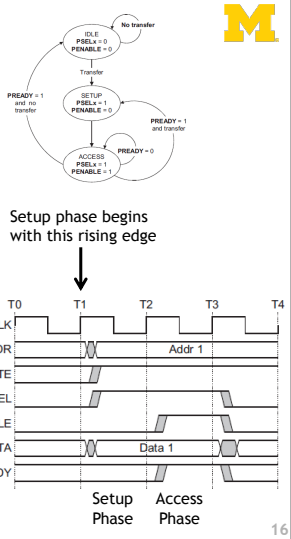
- Low-cost
- Low-power
- Low-complexity
- Low-bandwidth
- Non-pipelined
- Ideal for peripherals

Notation



APB bus state machine

- IDLE
  - Default APB state
- SETUP
  - When transfer required
  - PSELx is asserted
  - Only one cycle
- ACCESS
  - PENABLE is asserted
  - Addr, write, select, and write data remain stable
  - Stay if PREADY = L
  - Goto IDLE if PREADY = H and no more data
  - Goto SETUP is PREADY = H and more data pending

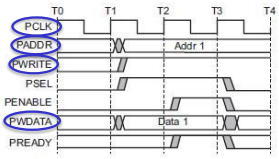


APB signal definitions

- PCLK: the bus clock source (rising-edge triggered)
- PRESETn: the bus (and typically system) reset signal (active low)
- PADDR: the APB address bus (can be up to 32-bits wide)
- PSELx: the select line for each slave device
- PENABLE: indicates the 2nd and subsequent cycles of an APB xfer
- PWRITE: indicates transfer direction (Write=H, Read=L)
- PWDATA: the write data bus (can be up to 32-bits wide)
- PREADY: used to extend a transfer
- PRDATA: the read data bus (can be up to 32-bits wide)
- PSLVERR: indicates a transfer error (OKAY=L, ERROR=H)

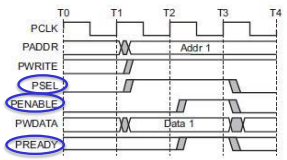
APB bus signals in action

- PCLK
  - Clock
- PADDR
  - Address on bus
- PWRITE
  - 1=Write, 0=Read
- PWDATA
  - Data written to the I/O device. Supplied by the bus master/processor.

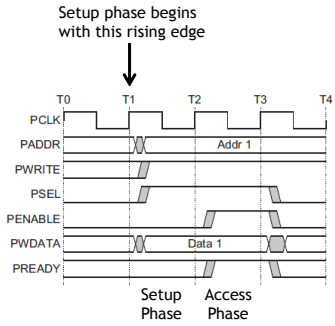


APB bus signals

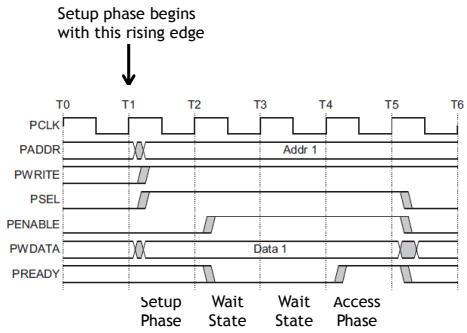
- PSEL
  - Asserted if the current bus transaction is targeted to ***this*** device
- PENABLE
  - High during entire transaction *other than* the first cycle.
- PREADY
  - Driven by target. Similar to our #ACK. Indicates if the target is *ready* to do transaction.
    Each target has it's own PREADY



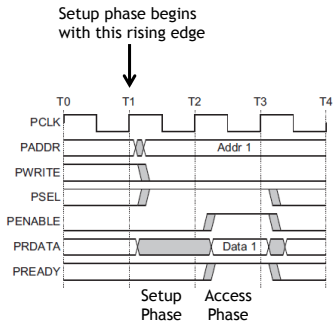
A write transfer with no wait states



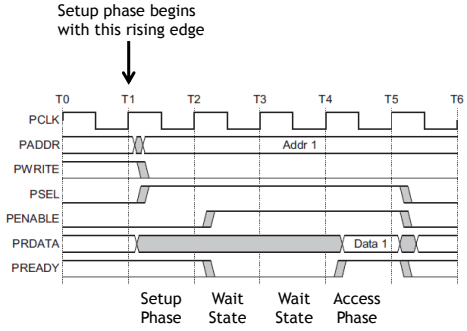
A write transfer with wait states



A read transfer with no wait states



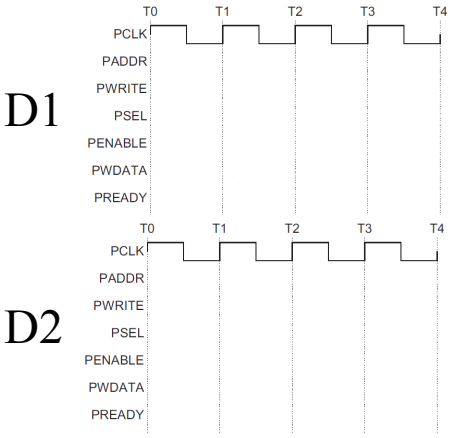
A read transfer with wait states



Example setup

- We will assume we have one bus master "CPU" and two slave devices (D1 and D2)
  - D1 is mapped to 0x00001000-0x0000100F
  - D2 is mapped to 0x00001010-0x0000101F

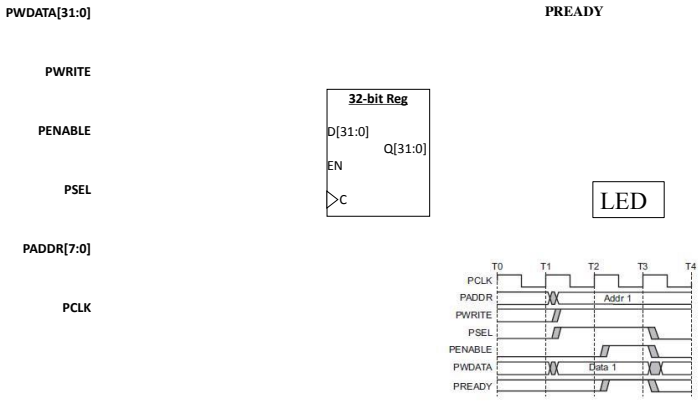
CPU stores to location 0x00001004 with no stalls



Writes

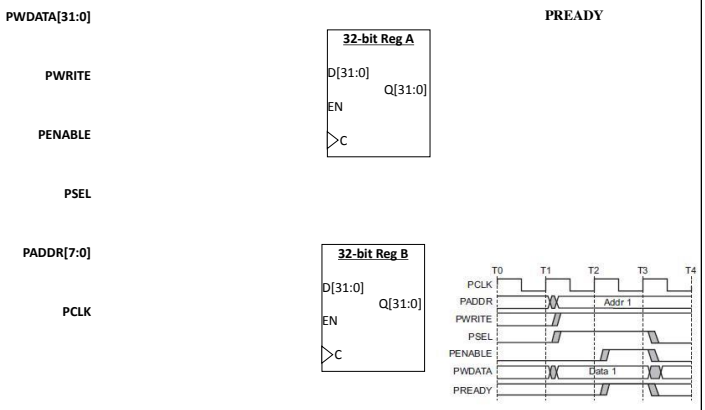
Let's do some hardware examples!

What if we want to have the LSB of this register control an LED?



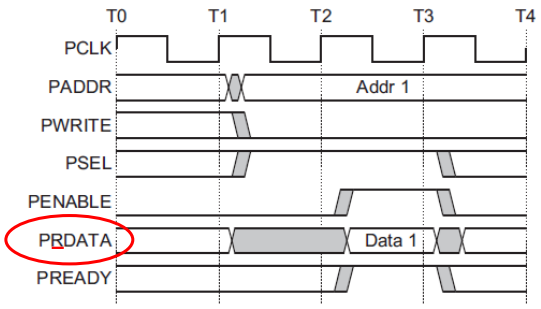
We are assuming APB only gets lowest 8 bits of address here…

Reg A should be written at address 0x00001000 Reg B should be written at address 0x00001004



We are assuming APB only gets lowest 8 bits of address here…

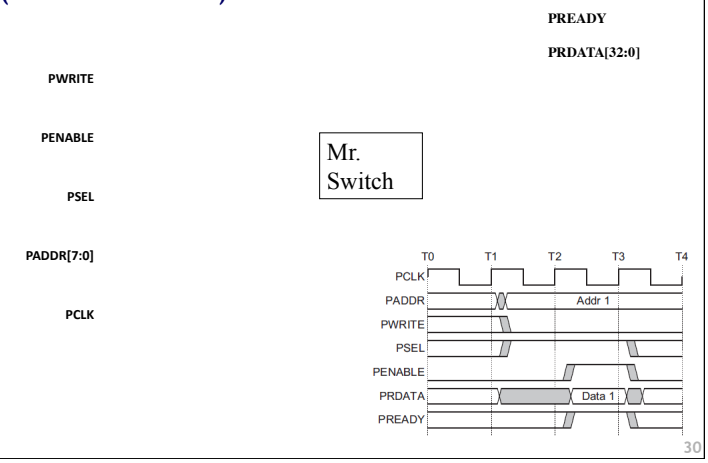
Reads…



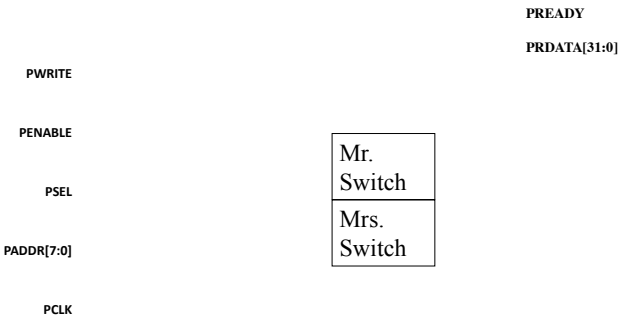
The key thing here is that each slave device has its own read data (PRDATA) bus!

Recall that "R" is from the initiator's viewpoint—the device drives data when read.

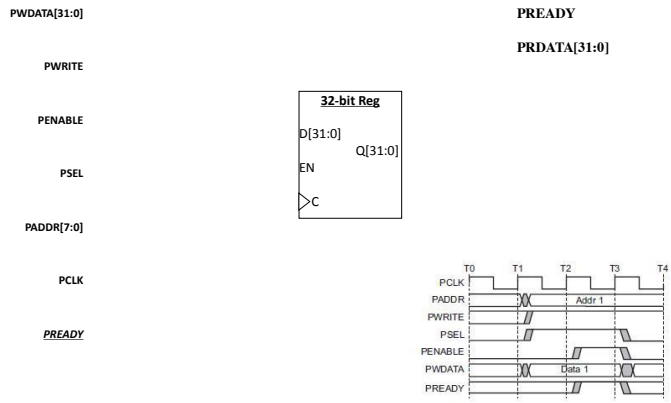
Let's say we want a device that provides data from a switch on a read to any address it is assigned. (so returns a 0 or 1)



Device provides data from switch A if address 0x00001000 is read from. B if address 0x00001004 is read from



All reads read from register, all writes write…



We are assuming APB only gets lowest 8 bits of address here…

Things left out…

- There is another signal, PSLVERR (APB Slave Error) which we can drive high if things go bad.
  - We'll just tie that to 0.
- PRESETn
  - Active low system reset signal
  - (needed for stateful peripherals)
- Note that we are assuming that our device need not stall.
  - We could stall if needed.
    - I can't find a limit on how long, but I suspect at some point the processor would generate an error.

Verilog!

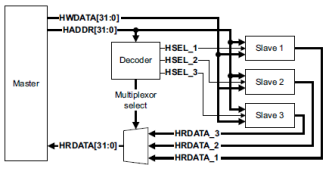
```
/*** APB3 BUS INTERFACE ***/
input PCLK,                     // clock
input PRESERN,                  // system reset
input PSEL,                     // peripheral select
input PENABLE,                  // distinguishes access phase
output wire PREADY,             // peripheral ready signal
output wire PSLVERR,            // error signal
input PWRITE,                   // distinguishes read and write cycles
input [31:0] PADDR,             // I/O address
input wire [31:0] PWDATA,       // data from processor to I/O device (32 bits)
output reg [31:0] PRDATA,       // data to processor from I/O device (32-bits)

/*** I/O PORTS DECLARATION ***/
output reg LEDOUT,              // port to LED
input SW                        // port to switch
);

assign PSLVERR = 0;
assign PREADY = 1;
```


Outline

- Announcements
- Accessing memory from Assembly/C
- Busses: The glue that connects the pieces
- ARM Advanced Peripheral Bus (APB)
- ARM Advanced High-performance Bus Light (AHB-Lite)

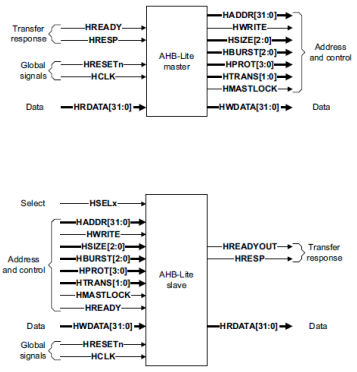
AHB-Lite supports single bus master and provides high-bandwidth operation

- Burst transfers
- Single clock-edge operation
- Non-tri-state implementation
- Configurable bus width



AHB-Lite bus master/slave interface

- Global signals
  - HCLK
  - HRESETn
- Master out/slave in
  - HADDR (address)
  - HWDATA (write data)
  - Control
    - HWRITE
    - HSIZE
    - HBURST
    - HPROT
    - HTRANS
    - HMASTLOCK
- Slave out/master in
  - HRDATA (read data)
  - HREADY
  - HRESP

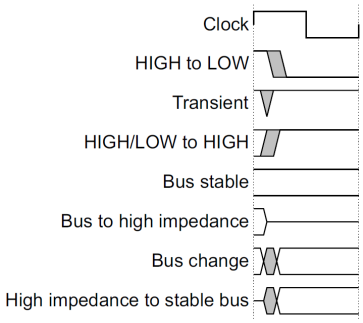


AHB-Lite signal definitions

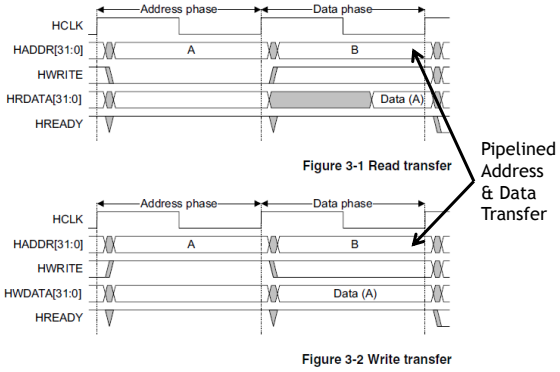
- Global signals
  - HCLK: the bus clock source (rising-edge triggered)
  - HRESETn: the bus (and system) reset signal (active low)
- Master out/slave in
  - HADDR[31:0]: the 32-bit system address bus
  - HWDATA[31:0]: the system write data bus
  - Control
    - HWRITE: indicates transfer direction (Write=1, Read=0)
    - HSIZE[2:0]: indicates size of transfer (byte, halfword, or word)
    - HBURST[2:0]: indicates single or burst transfer (1, 4, 8, 16 beats)
    - HPROT[3:0]: provides protection information (e.g. I or D; user or handler)
    - HTRANS: indicates current transfer type (e.g. idle, busy, nonseq, seq)
    - HMASTLOCK: indicates a locked (atomic) transfer sequence
- Slave out/master in
  - HRDATA[31:0]: the slave read data bus
  - HREADY: indicates previous transfer is complete
  - HRESP: the transfer response (OKAY=0, ERROR=1)

Key to timing diagram conventions

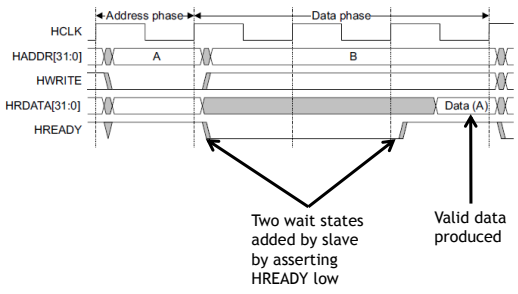
- Timing diagrams
  - Clock
  - Stable values
  - Transitions
  - High-impedance
- Signal conventions
  - Lower case 'n' denote active low (e.g. RESETn)
  - Prefix 'H' denotes AHB
  - Prefix 'P' denotes APB



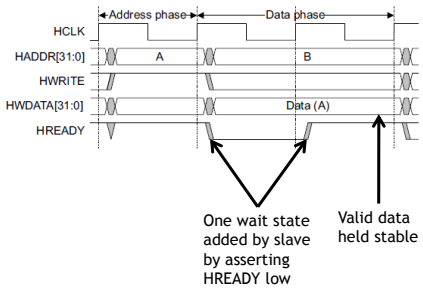
Basic read and write transfers with no wait states



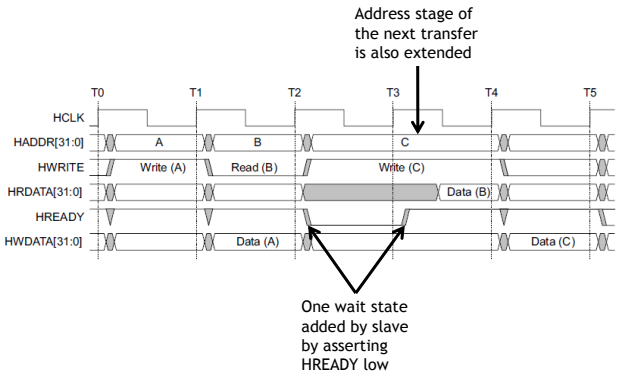
Read transfer with two wait states



Write transfer with one wait state



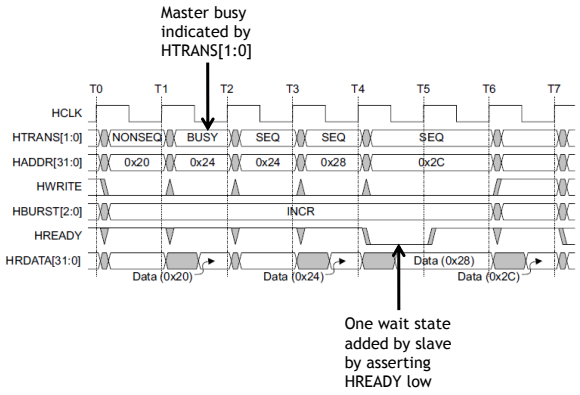
Wait states extend the address phase of next transfer



Transfers can be of four types (HTRANS[1:0])

- IDLE (b00)
  - No data transfer is required
  - Slave must OKAY w/o waiting
  - Slave must ignore IDLE
- BUSY (b01)
  - Insert idle cycles in a burst
  - Burst will continue afterward
  - Address/control reflects next transfer in burst
  - Slave must OKAY w/o waiting
  - Slave must ignore BUSY
- NONSEQ (b10)
  - Indicates single transfer or first transfer of a burst
  - Address/control unrelated to prior transfers
- SEQ (b11)
  - Remaining transfers in a burst
  - Addr = prior addr + transfer size

A four beat burst with master busy and slave wait



Controlling the size (width) of a transfer

- HSIZE[2:0] encodes the size
- The cannot exceed the data bus width (e.g. 32-bits)
- HSIZE + HBURST is determines wrapping boundary for wrapping bursts
- HSIZE must remain constant throughout a burst transfer

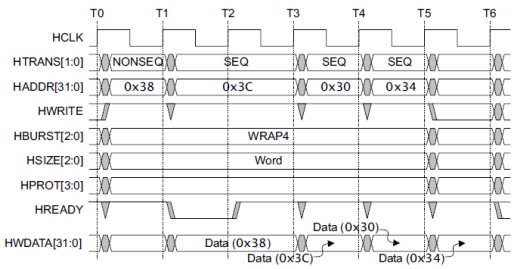
HSIZE[2]	HSIZE[1]	HSIZE[0]	Size (bits)	Description
0	0	0	8	Byte
0	0	1	16	Halfword
0	1	0	32	Word
0	1	1	64	Doubleword
1	0	0	128	4-word line
1	0	1	256	8-word line
1	1	0	512	-
1	1	1	1024	-

Controlling the burst beats (length) of a transfer

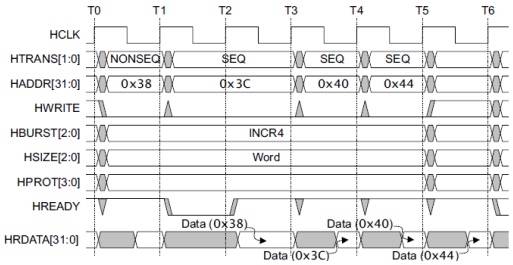
- Burst of 1, 4, 8, 16, and undef
- HBURST[2:0] encodes the type
- Incremental burst
- Wrapping bursts
  - 4 beats x 4-byte words wrapping
  - Wraps at 16 byte boundary
  - E.g. 0x34, 0x38, 0x3c, 0x30,…
- Bursts must not cross 1KB address boundaries

HBURST[2:0]	Type	Description
b000	SINGLE	Single burst
b001	INCR	Incrementing burst of undefined length
b010	WRAP4	4-beat wrapping burst
b011	INCR4	4-beat incrementing burst
b100	WRAP8	8-beat wrapping burst
b101	INCR8	8-beat incrementing burst
b110	WRAP16	16-beat wrapping burst
b111	INCR16	16-beat incrementing burst

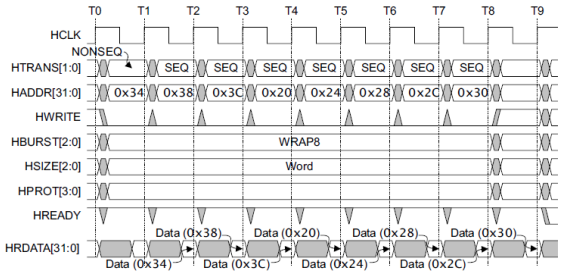
A four beat wrapping burst (WRAP4)



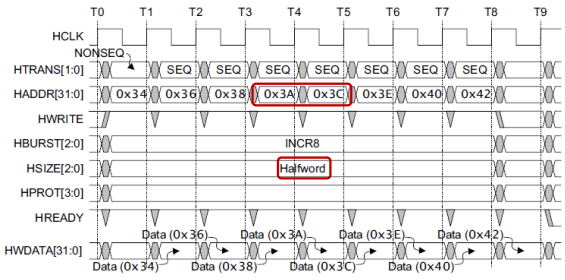
A four beat incrementing burst (INCR4)



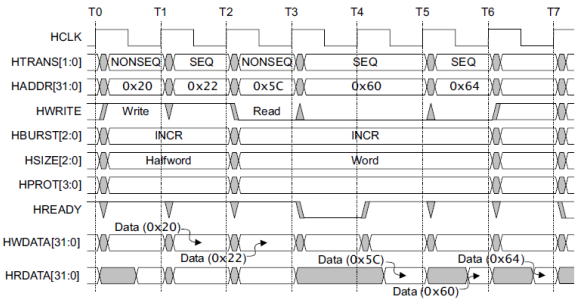
An eight beat wrapping burst (WRAP8)



An eight beat incrementing burst (INCR8) using half-word transfers

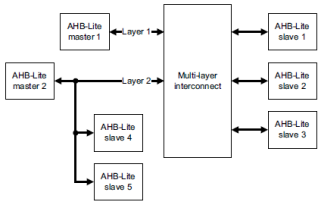


An undefined length incrementing burst (INCR)



Multi-master AHB-Lite requires a multi-layer interconnect

- AHB-Lite is single-master
- Multi-master operation
  - Must isolate masters
  - Each master assigned to layer
  - Interconnect arbitrates slave accesses
- Full crossbar switch often unneeded
  - Slaves 1, 2, 3 are shared
  - Slaves 4, 5 are local to Master 1



Questions?

Comments?

Discussion?