

# Homework 3, due Thurs, Feb 9, 2006

February 2, 2006

To be done in groups.

## 1 Unequal Divide and Conquer (revised from last time)

### 1.1 Guidelines

In these problems, you are not required to find  $c$  and  $n_0$  and prove the result by induction. It is sufficient to give a proof that appeals to results in the chapter. For example, you can say, without further comment, “ $5n^2 + 7n \leq O(n^2)$ ” or “ $\sum_{j=1}^n j^3 = \Theta(\int_0^n x^3 dx) = \Theta(n^4)$ ” or “a  $k$ -ary tree with height  $h$  has  $\Theta(k^h)$  leaves and nodes” ( $k > 1$ ). In fact, I’d prefer that you get some practice in this “bigger picture” way of thinking, and not always get bogged down with induction details.

### 1.2 The Problem

Exercises CLRS 4.2-4 and CLRS 4.2-5 illustrate that, if we use a divide-and-conquer approach to a problem, it is typically better to divide the problem into nearly equal-sized subproblems. The above and the last example of CLRS Section 4.2 also illustrate that it is not necessary that the subproblems be *exactly* the same size.

We now switch from *analyzing* to *designing* algorithms. Given a problem of size  $n$ , we will break it into two problems, of size  $a(n)$  and  $n - a(n)$ , where  $a(n)$  is a “common” function satisfying  $1 \leq a(n) \leq n/2$ . We assume the Divide and Recombine cost is linear, so we have recurrence

$$T(n) = T(a(n)) + T(n - a(n)) + cn.$$

All other things being equal, we’d want  $a(n) = n/2$  to minimize  $T(n)$ . But sometimes there is a separate (direct or indirect) cost involved in making  $a(n)$  exactly  $n/2$  and it’s easier to choose other  $a()$ ’s. Below we investigate which  $a()$ ’s are acceptable while still meeting certain overall cost requirements. Problem:

- How slowly-growing can  $a(n)$  be and still make  $T(n) \leq O(n \log^2(n))$ ?

Restrict attention to functions of the form  $n^r(\log(n))^s$ , where  $r$  and  $s$  are constant real numbers, and you only need to find  $s$  up to 1, additively. That is, your answer should be numbers  $r$  and  $s$  with

$$n^r(\log(n))^s \leq a(n) \leq n^r(\log(n))^{s+1}.$$

(If you’ve already found exactly the right  $s$ , please turn it in, but don’t attempt this if you haven’t started.)

## 2 New Problems

- CLRS 5.1-3, on biased coins. You may assume that different tosses of the coin are independent.
- CLRS 5.2-5, expected number of inversions