

# EECS 583 – Class 18

## Research Topic 1

### Breaking Dependencies, Dynamic Parallelization

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*University of Michigan*

*November 21, 2011*

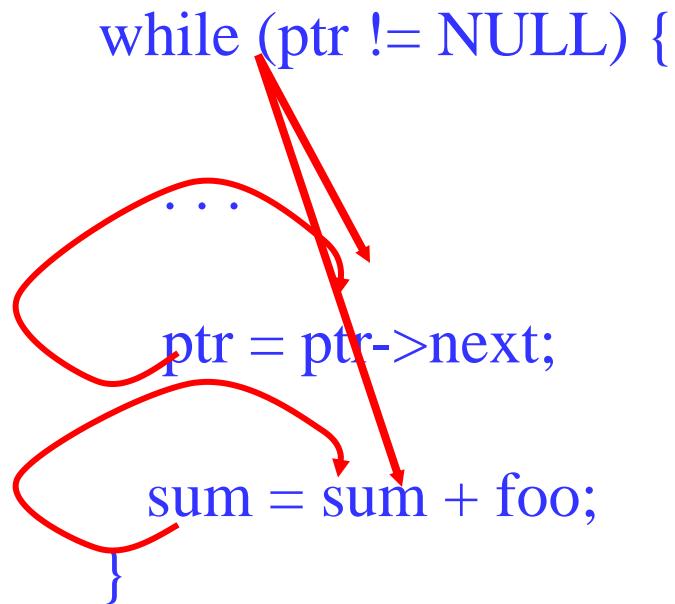
# Announcements & Reading Material

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- ❖ No class on Wednes (no paper summary either!)
- ❖ We are grading the exams
- ❖ Today's class reading
  - » “Spice: Speculative Parallel Iteration Chunk Execution,” E. Raman, N. Vachharajani, R. Rangan, and D. I. August, *Proc 2008 Intl. Symposium on Code Generation and Optimization*, April 2008.
- ❖ Next class reading (Monday, Nov 28)
  - » “Exploiting coarse-grained task, data, and pipeline parallelism in stream programs,” M. Gordon, W. Thies, and S. Amarasinghe, *Proc. of the 12th International Conference on Architectural Support for Programming Languages and Operating Systems*, Oct. 2006.

# The Data Dependence Problem

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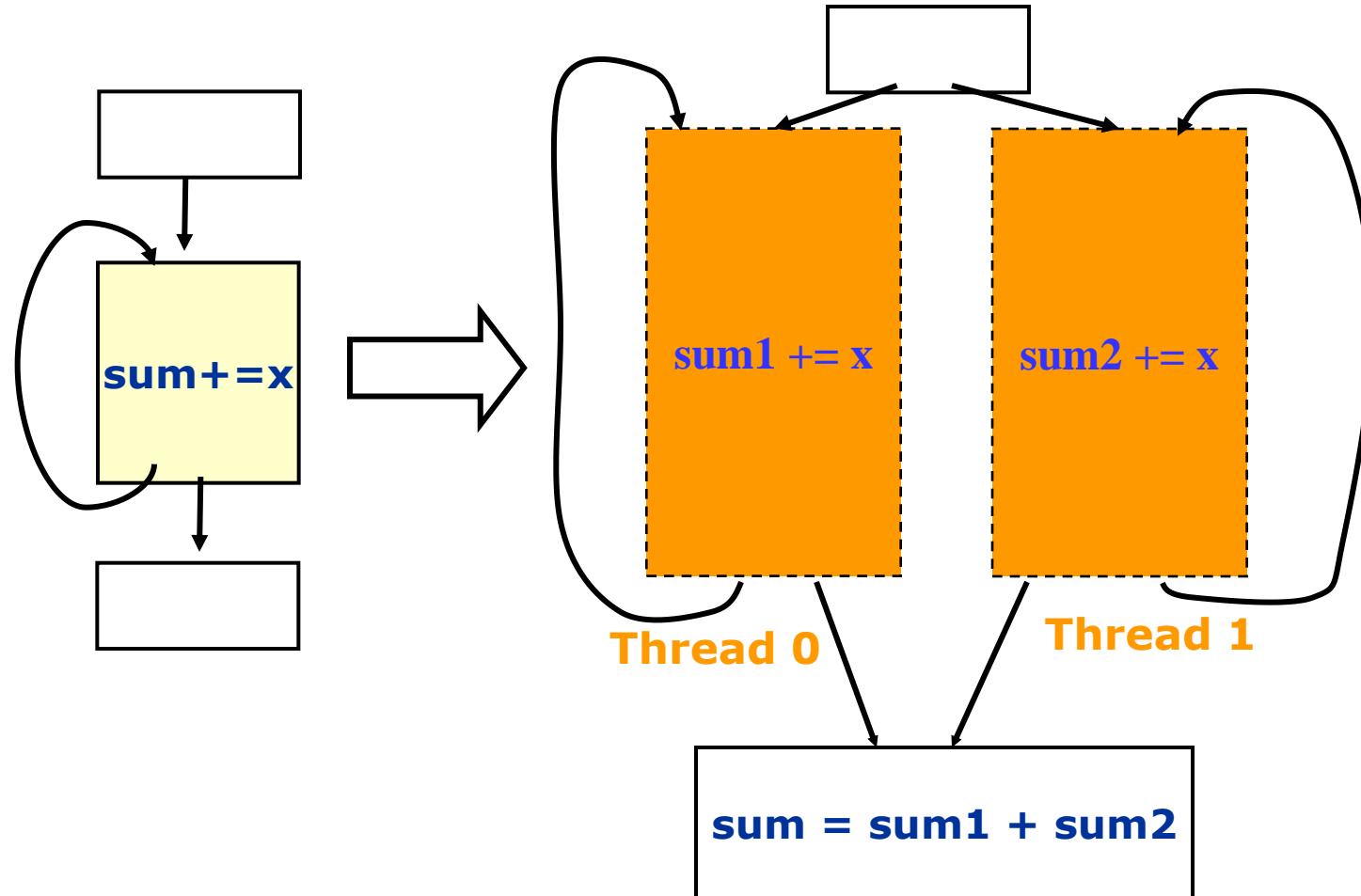


- How to deal with control dependences?
- How to deal with linked data structures?
- How to remove programmatic dependences?

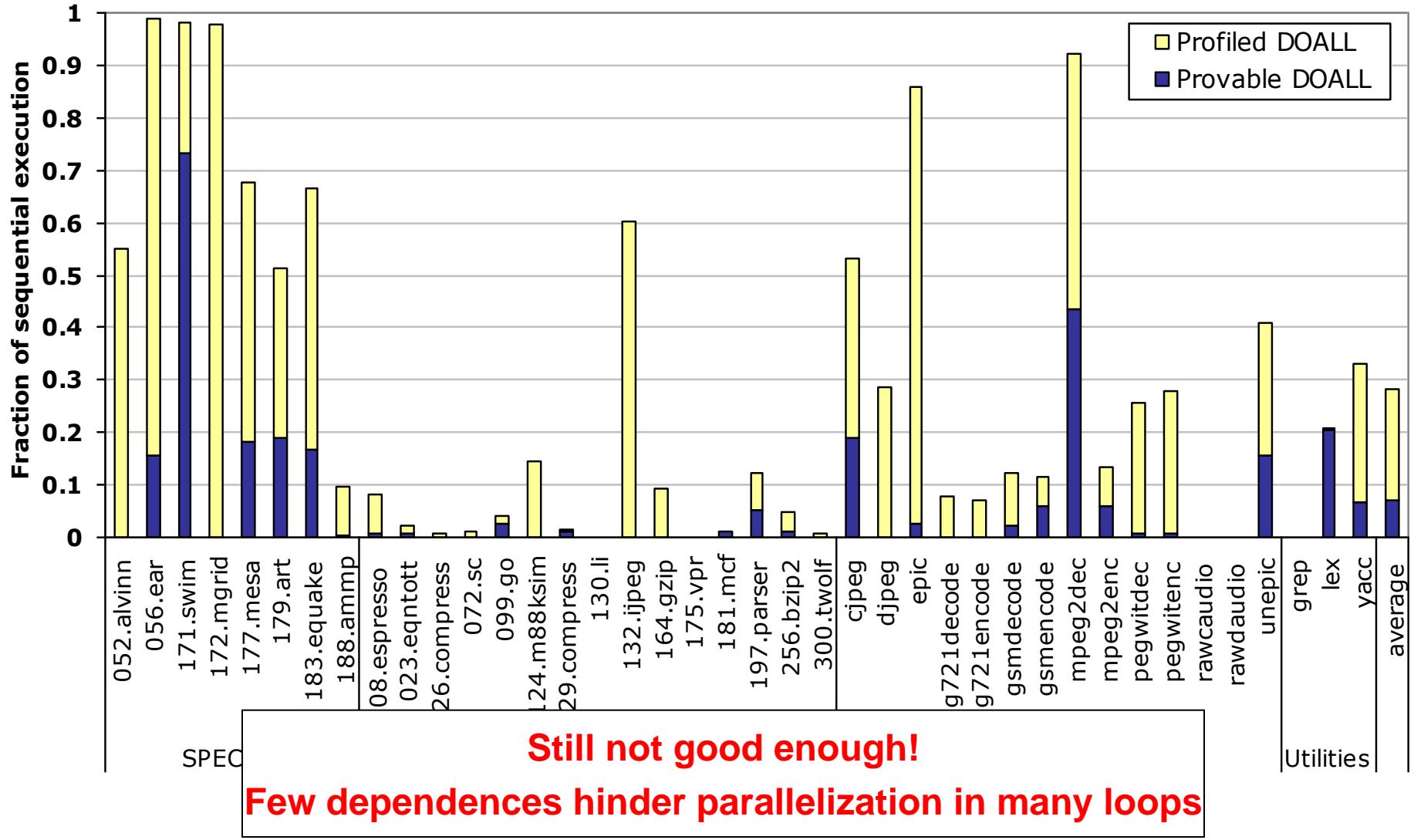
# We Know How to Break Some of These Dependences – Recall ILP Optimizations

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Apply accumulator variable expansion!



# DOALL Coverage – Provable and Profiled



# Speculative Loop Fission

```
1: while (node) {  
2:   work(node);  
3:   node = node->next;  
}
```

Sequential part

```
1: while (node) {  
4:   node_array[count++] = node;  
3:   node = node->next;  
}
```

If this were traversing an array, it would  
be a DOALL loop

Separate out data structure access  
and work

Parallelize work

Parallel part

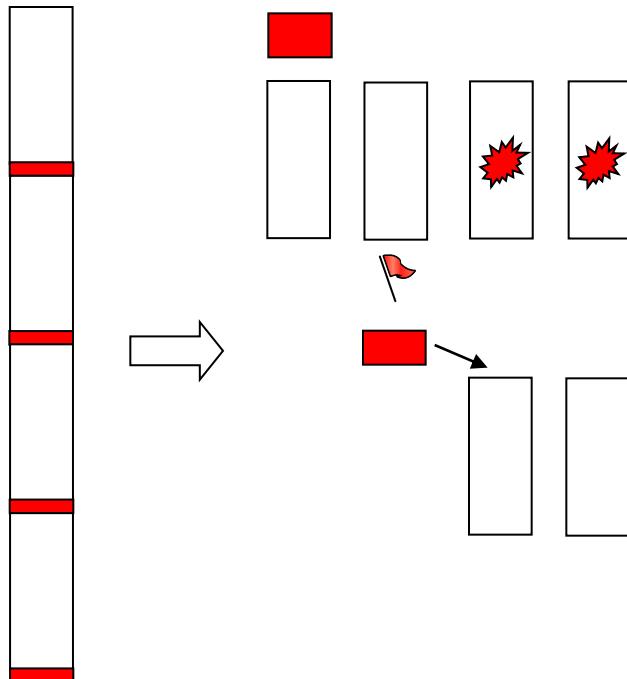
```
XBEGIN  
5: node = node_array[IS];  
i = 0;  
1':while (node && i++ < CS) {  
2:   work(node);  
3':   node = node->next;  
}  
RECV (THREADj-1)  
XCOMMIT  
SEND (THREADj+1)  
}
```

# Execution of Fissioned Loop

```
1: while (node) {  
2:   work(node);  
3:   node = node->next;  
}
```



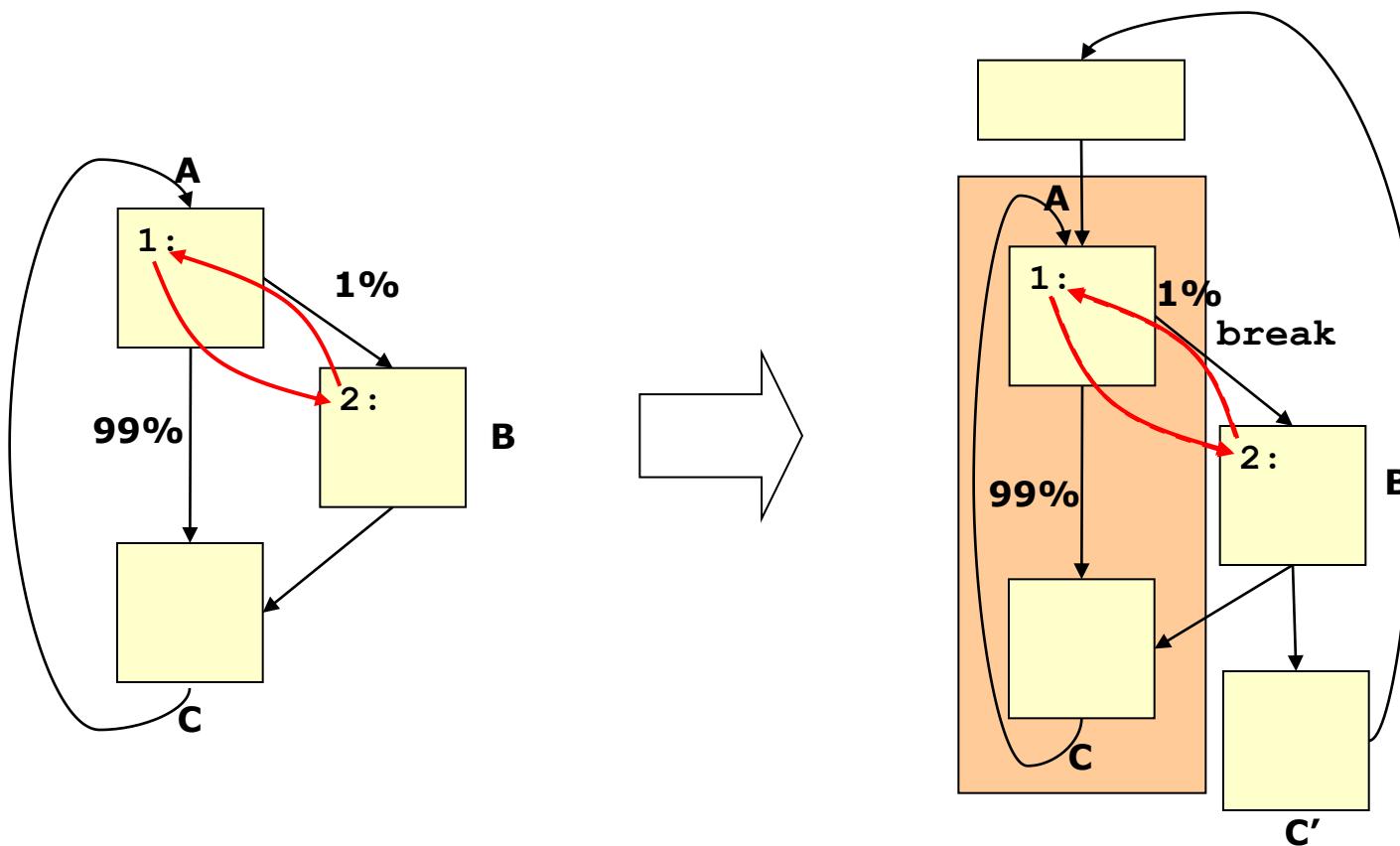
```
1: while (node) {  
4:   node_array[count++] = node;  
3:   node = node->next;  
}
```



```
XBEGIN  
5: node = node_array[IS];  
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2:   work(node);  
3':   node = node->next;  
}  
RECV (THREADj-1)  
XCOMMIT  
SEND (THREADj+1)  
}
```

# Infrequent Dependence Isolation

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# Infrequent Dependence Isolation - cont

```
for( j=0; j<=nstate; ++j ) {
    if( tystate[j] == 0 ) continue;
    if( tystate[j] == best ) continue;
    count = 0;
    cbest = tystate[j];
    for (k=j; k<=nstate; ++k)
        if (tystate[k]==cbest) ++count;
    if ( count > times) {
        best = cbest;
        times = count;
    }
}
```

1 %

Sample loop from yacc benchmark

```
j=0;
while (j<=nstate) {

    for( ; j<=nstate; ++j ) {
        if( tystate[j] == 0 ) continue;
        if( tystate[j] == best ) continue;
        count = 0;
        cbest = tystate[j];
        for (k=j; k<=nstate; ++k)
            if (tystate[k]==cbest) ++count;
        if ( count > times)
            break;
    }

    if (count > times) {
        best = cbest;
        times = count; j++;
    }
}
```

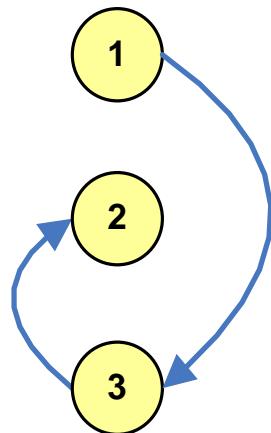
1 %

1 %

# Speculative Prematerialization

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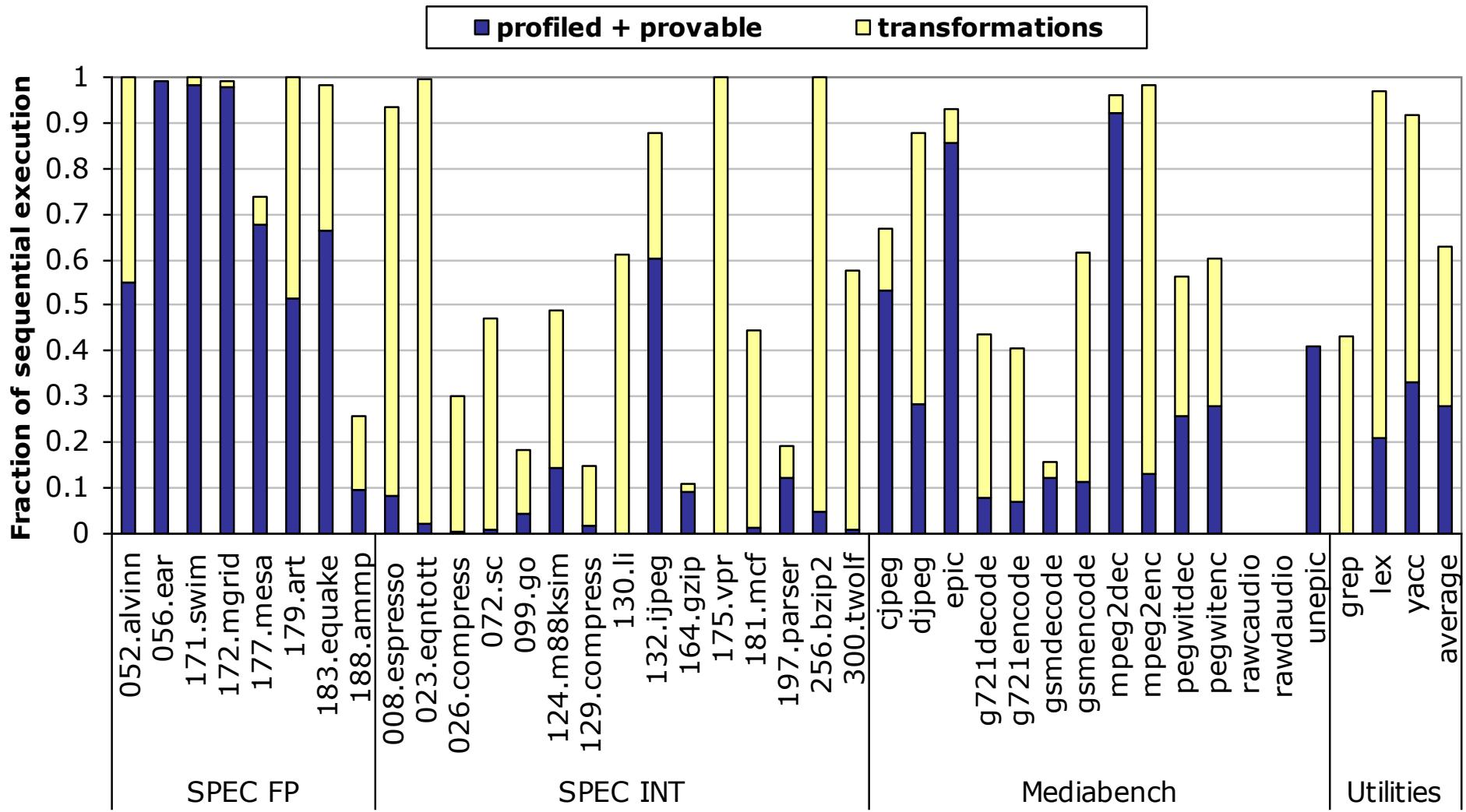
```
for (...) {  
1:   current = ... ;  
2:   work(last);  
3:   last = current;  
}
```



```
XBEGIN  
1': current =  
3': last =  
    for (...) {  
1:      current = ... ;  
2:      work(last);  
3:      last = current;  
    }  
XCOMMIT
```

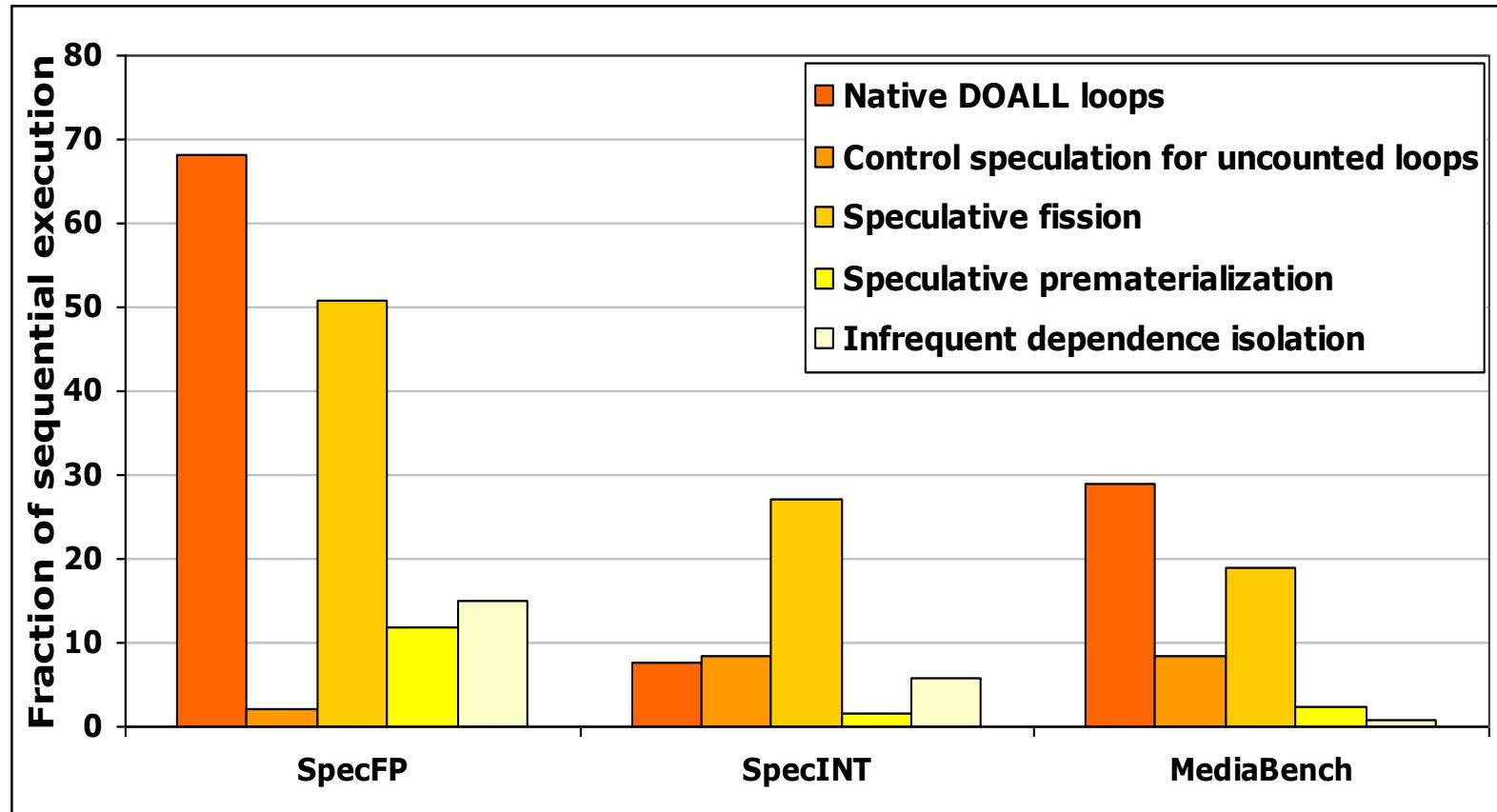
- ❖ Iterations are sequential due to distance 1 recurrence of the red variables
- ❖ Decouple iteration chunks by prematerializing *last* outside the loop

# DOALL Coverage – Profiled and Transformed



# Coverage Breakdown

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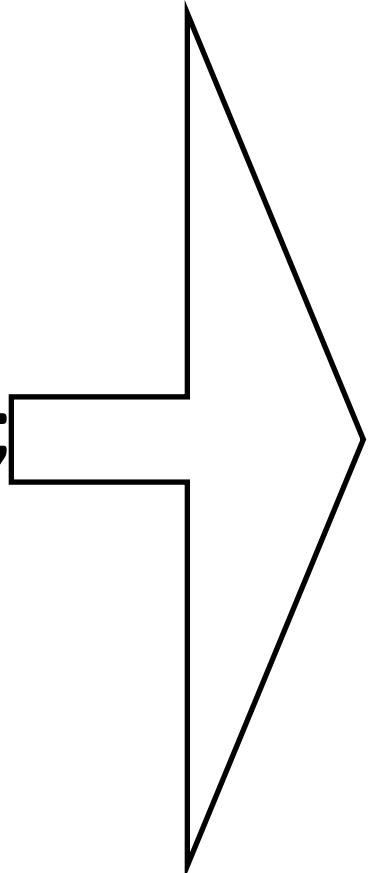


# Improving DSWP with Value Speculation

## Motivating Example

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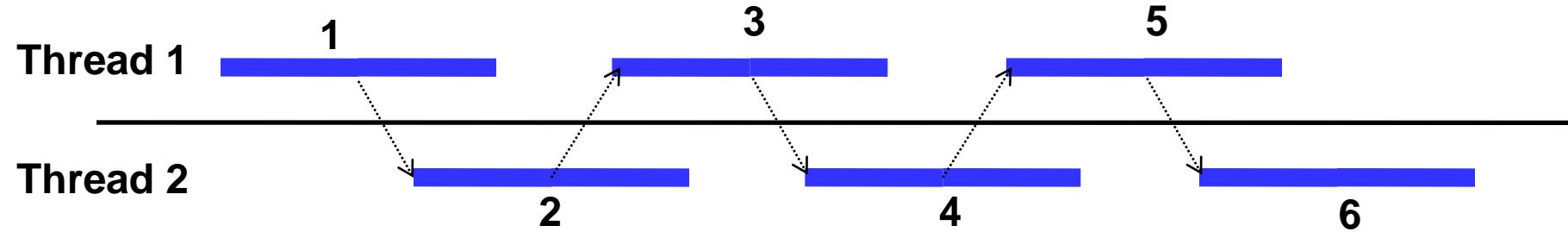
```
I = create_list();
while (...){
    process_list(I);
    modify_list(I);
}
```



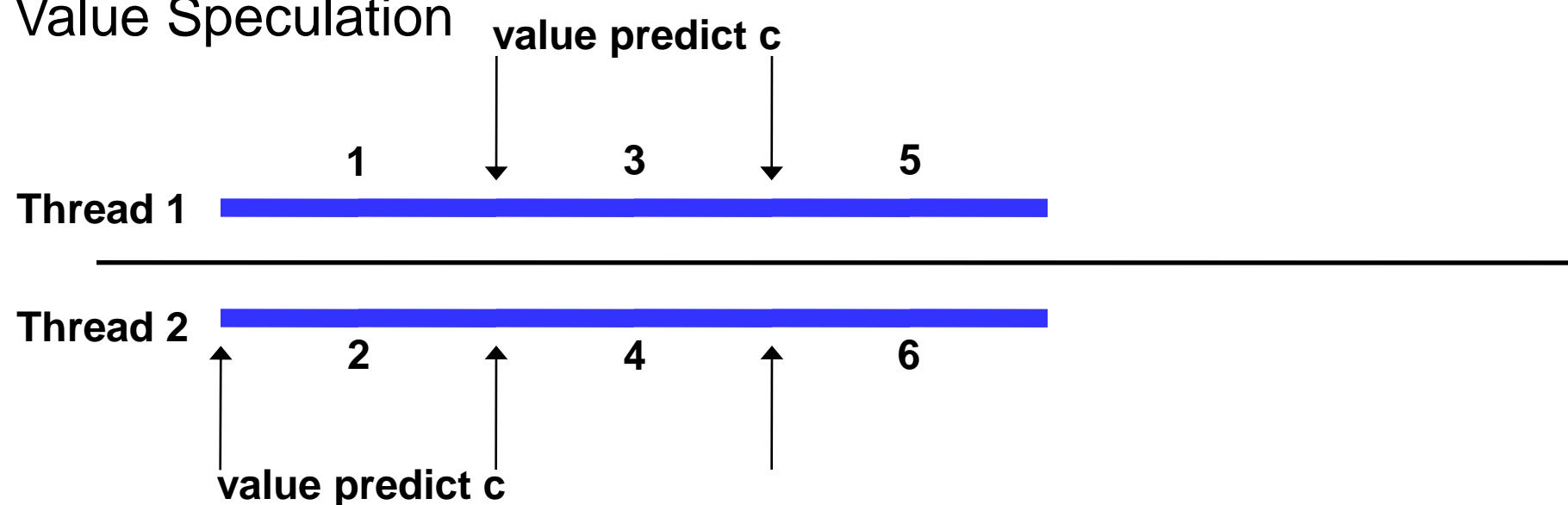
```
while (c != null){
    c1 = c;
    c = c->next_cl;
    int w = c1->pick_weight;
    if (w < wm){
        wm = w;
        cm = c1;
    }
}
```

# Need to Break Cross-thread Dependencies

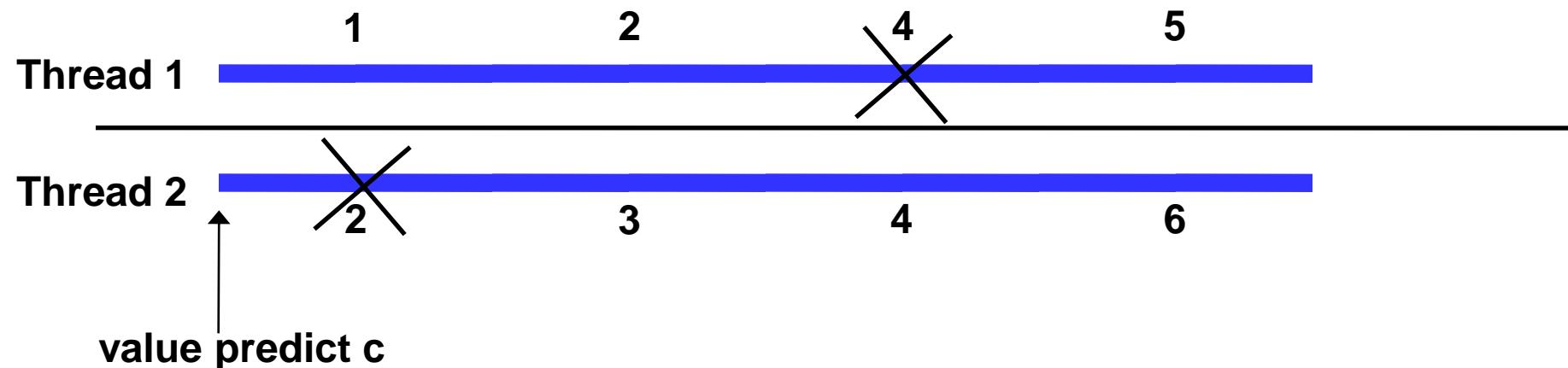
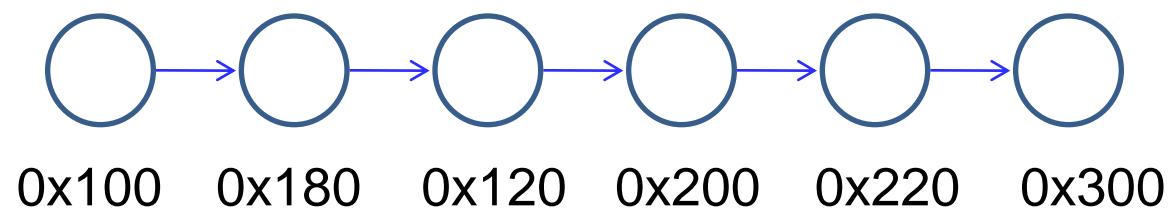
## Synchronization



## Value Speculation



# Value Speculation



# Value Speculation for TLP: Observation 1

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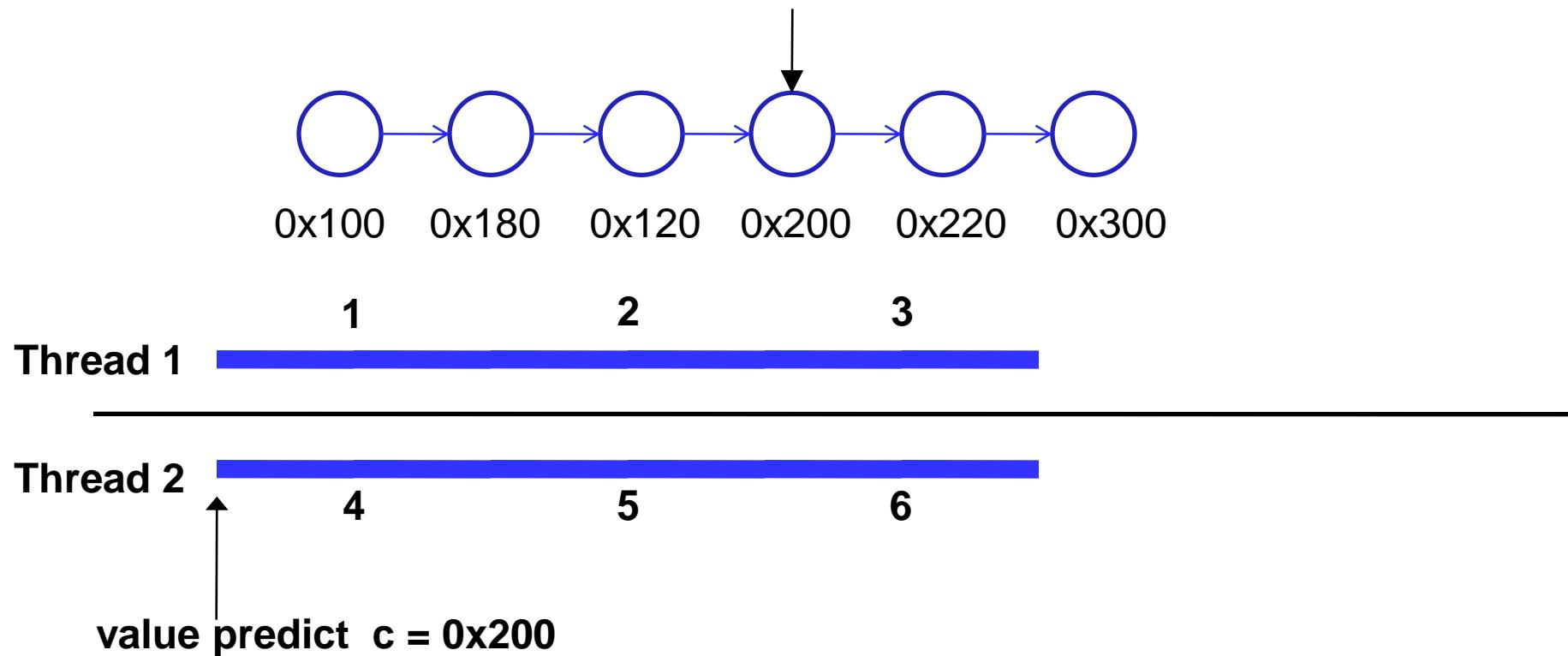
- ❖ Predicting a value in a *particular* iteration is harder than predicting *a value* of the operation within a *range* of iterations



# Value Speculation for TLP: Observation 2

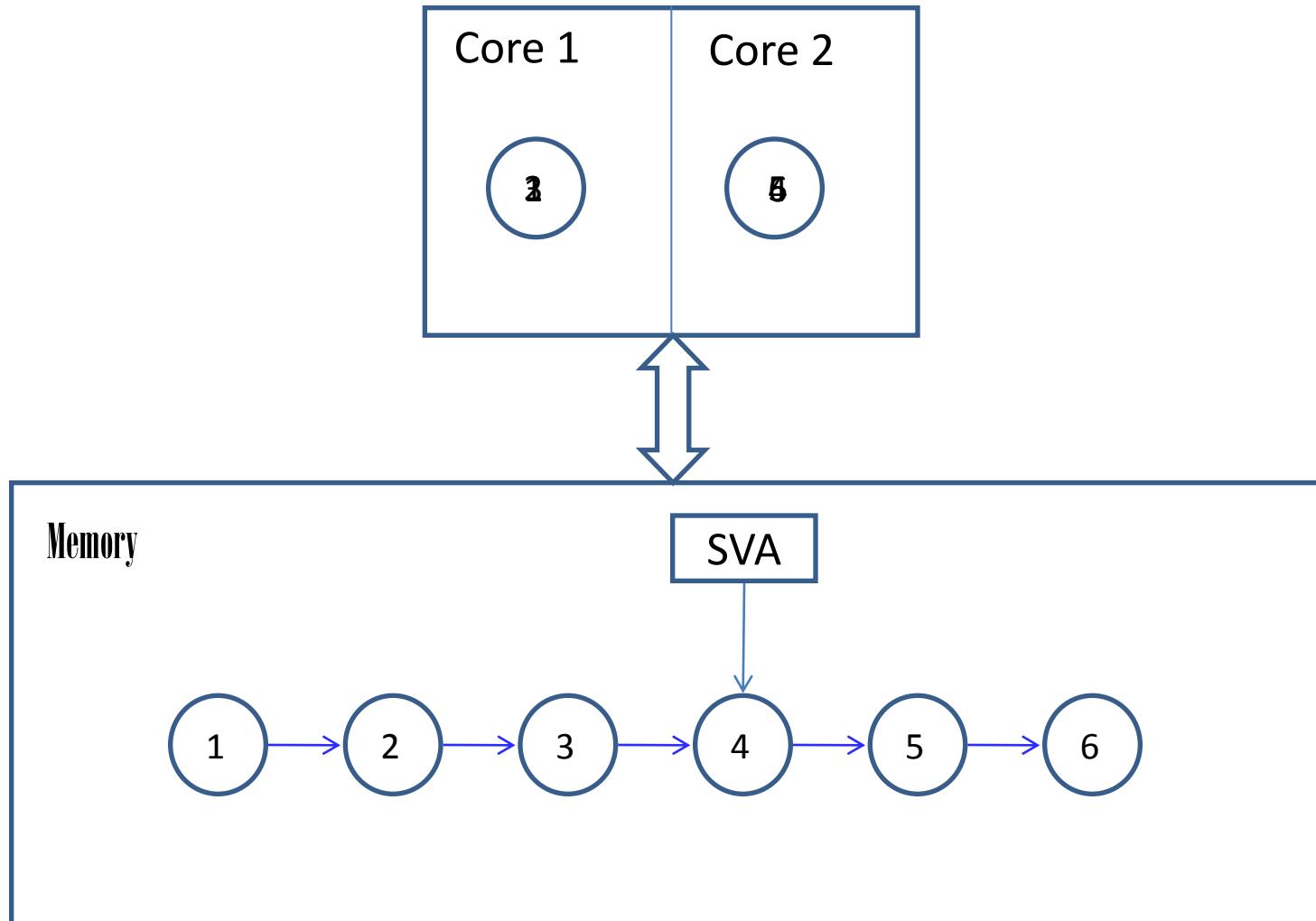
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- ❖ It is sufficient to predict the values in a few iterations to get TLP



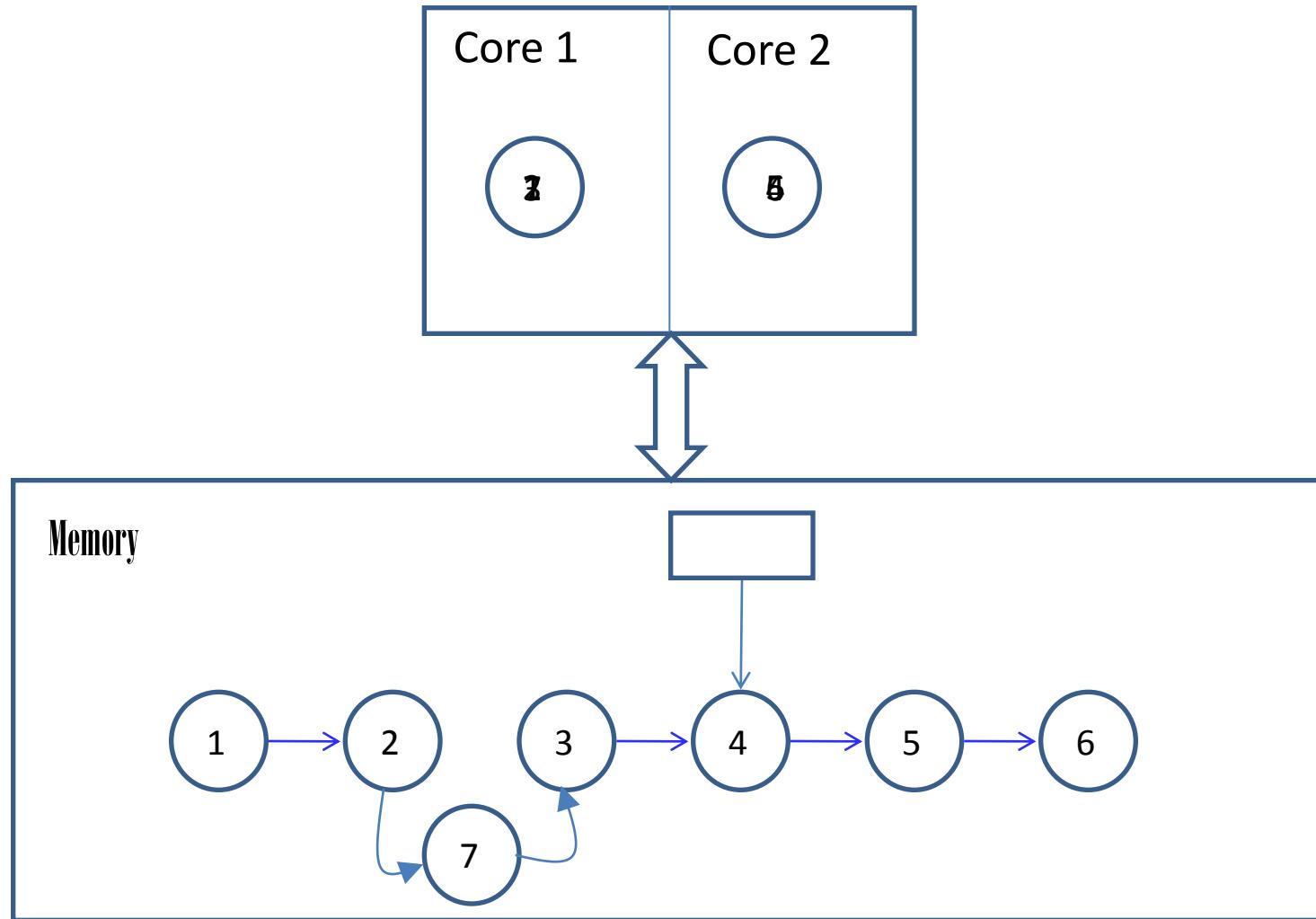
# Memoize and Speculative Parallel Execution

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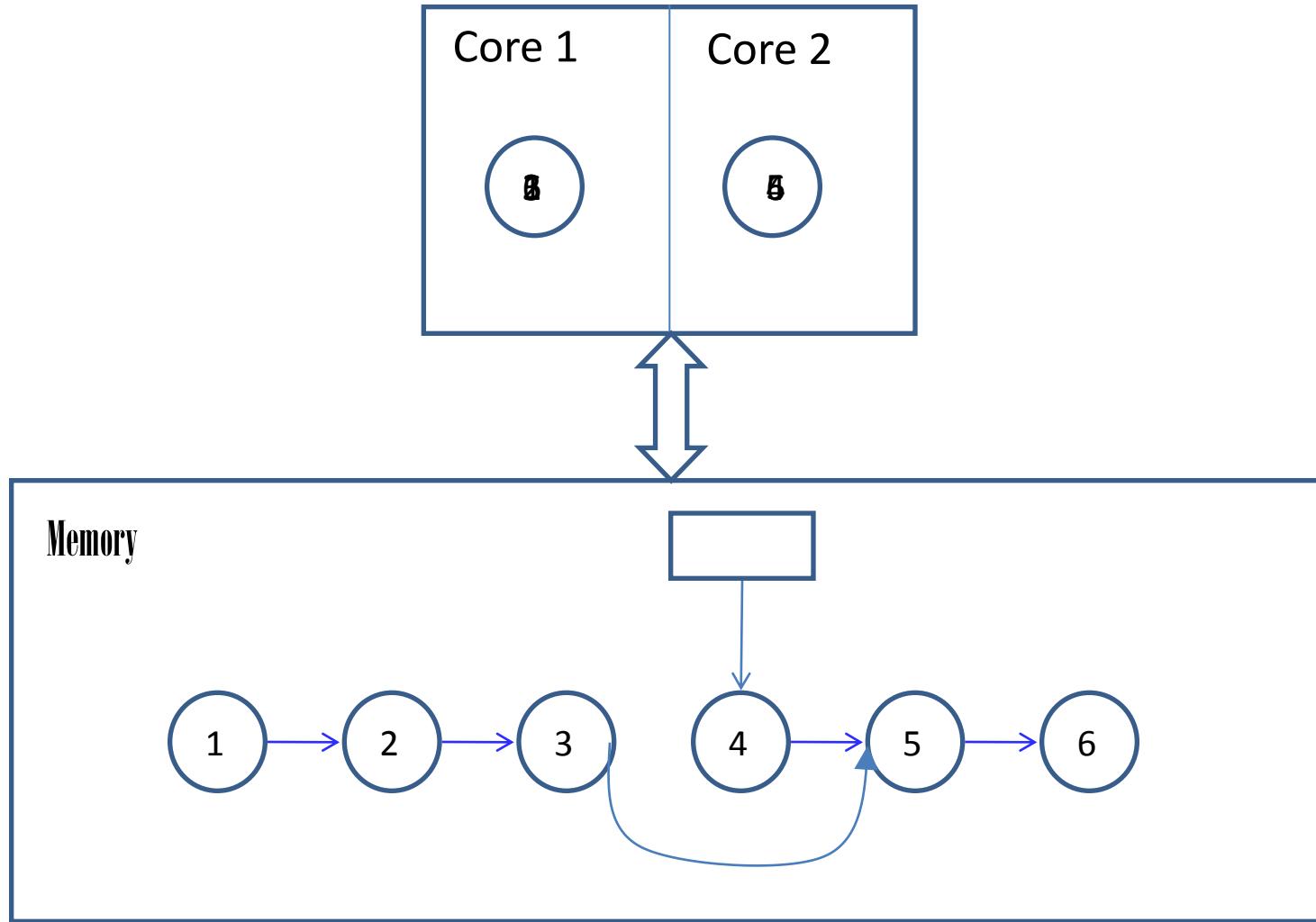
# Speculative Parallel Execution

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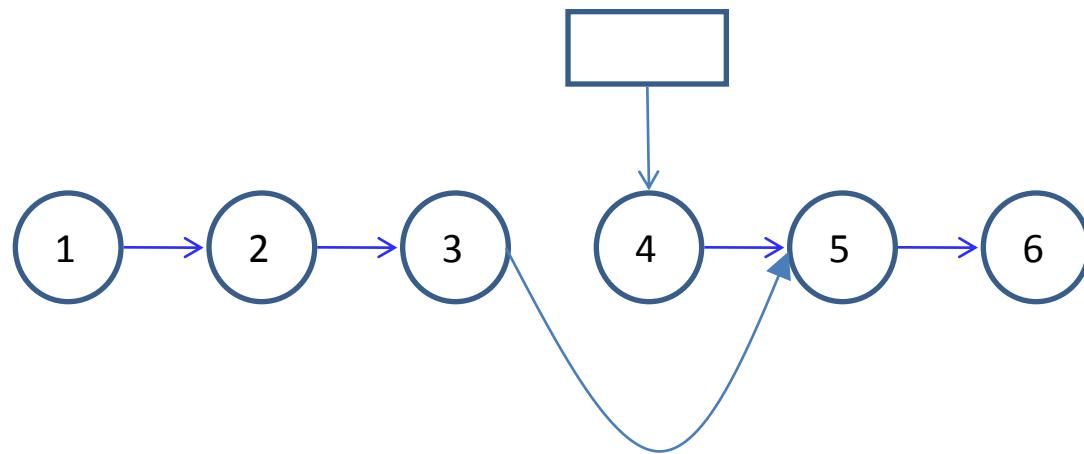
# Mis-speculation

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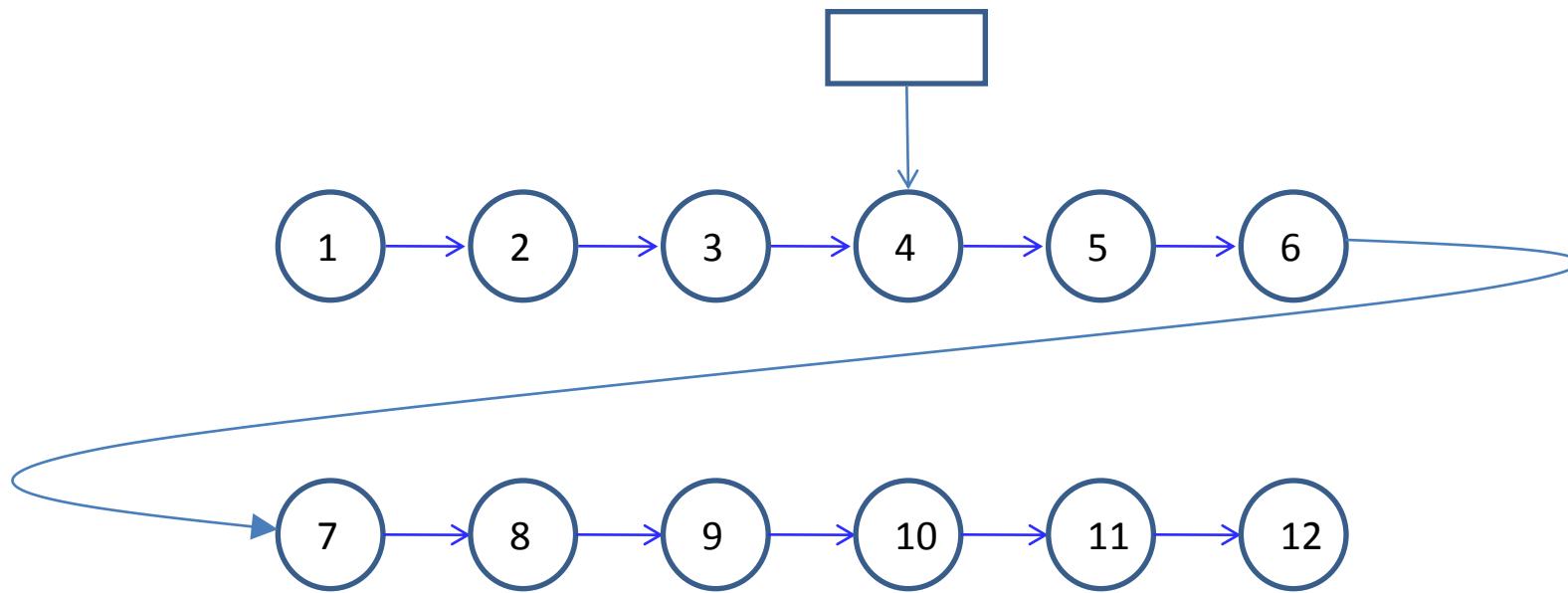
# Memoizing Once: Problems – Delete One of the Starting Nodes

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# Memoizing Once: Problems – The List Grows after the Initial Traversal

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# Discussion Point 1 – Fission vs Value Speculation

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- ❖ When is fission better?
- ❖ When is value speculation better?
- ❖ Is either more powerful than the other?
- ❖ Which would you use in your compiler and why?
- ❖ Would you be confident your code would run correctly if these transformations were done?

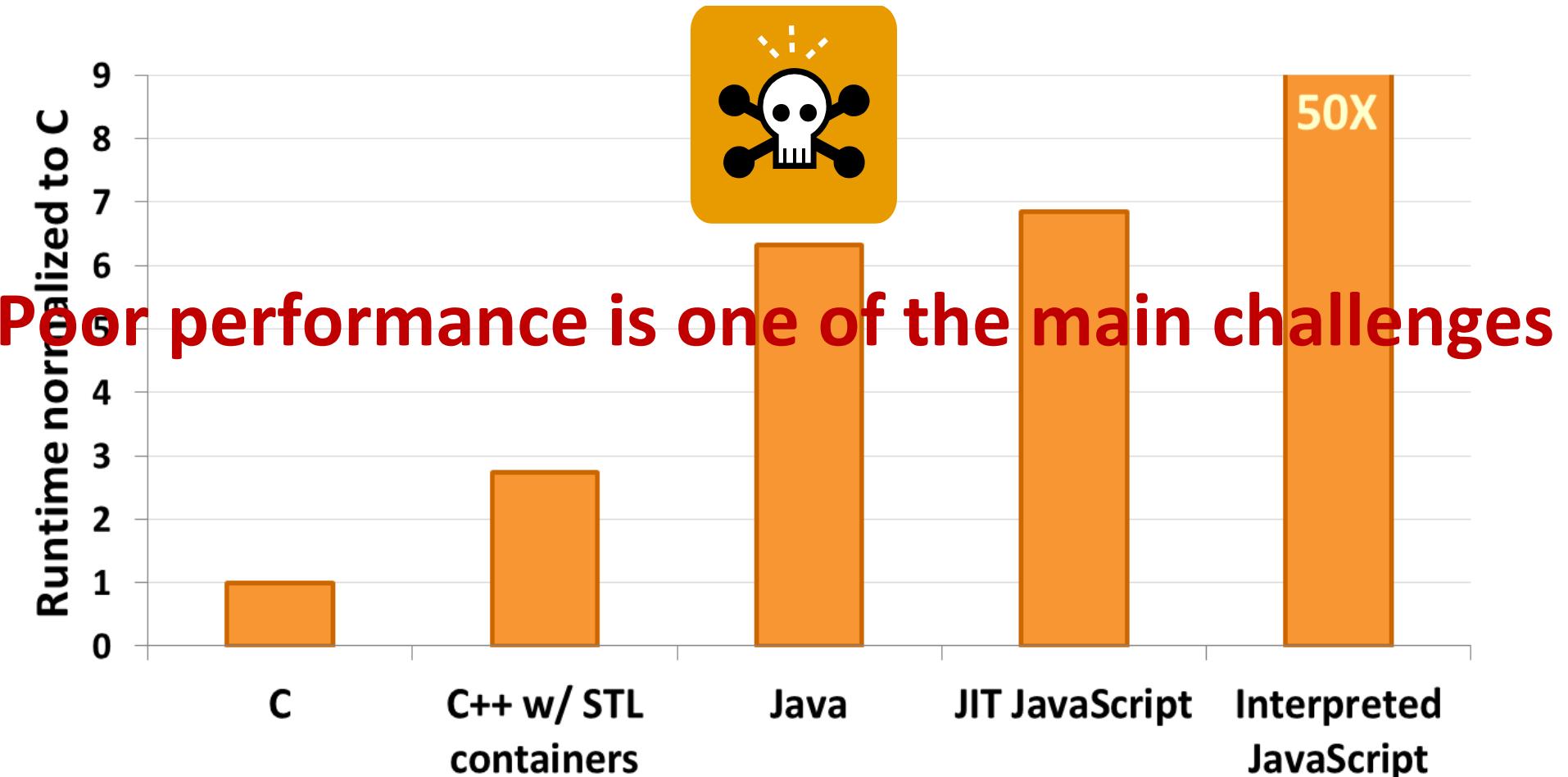
## Discussion Point 2 – Dependences

- ❖ Are these “tricks” valid for other common data structures?
  - » Are any generalizations needed?
- ❖ What other common data dependences are we missing in loops?

# Client-side Computation in JavaScript

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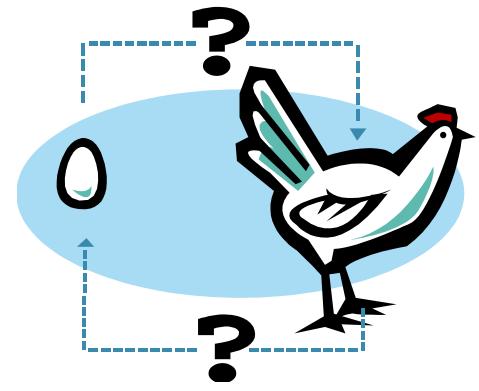
- ❖ Flexibility, ease of prototyping, and portability



# Client-side Applications

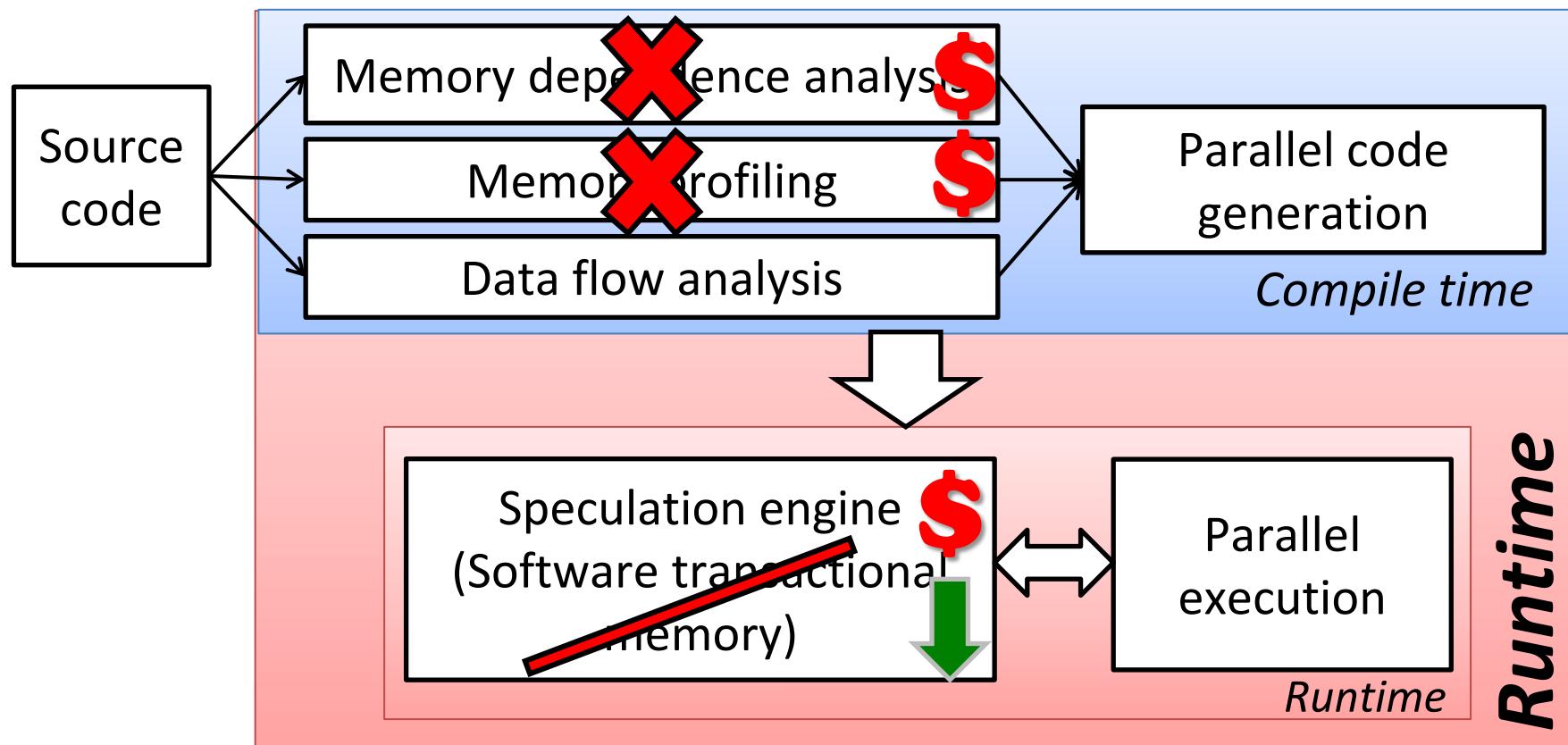
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- **Interaction-intensive:**
  - Largely composed of event handlers, triggered by user
  - Examples are *Gmail*, *Facebook*, etc.
- **Compute-intensive:**
  - Dominated by loops and hot functions
  - Online image editing such as Adobe's *Photoshop.com*, Google's *Picnik*
  - Lot more potential:
    - Online games
    - Video editing
    - Sound editing and voice recognition



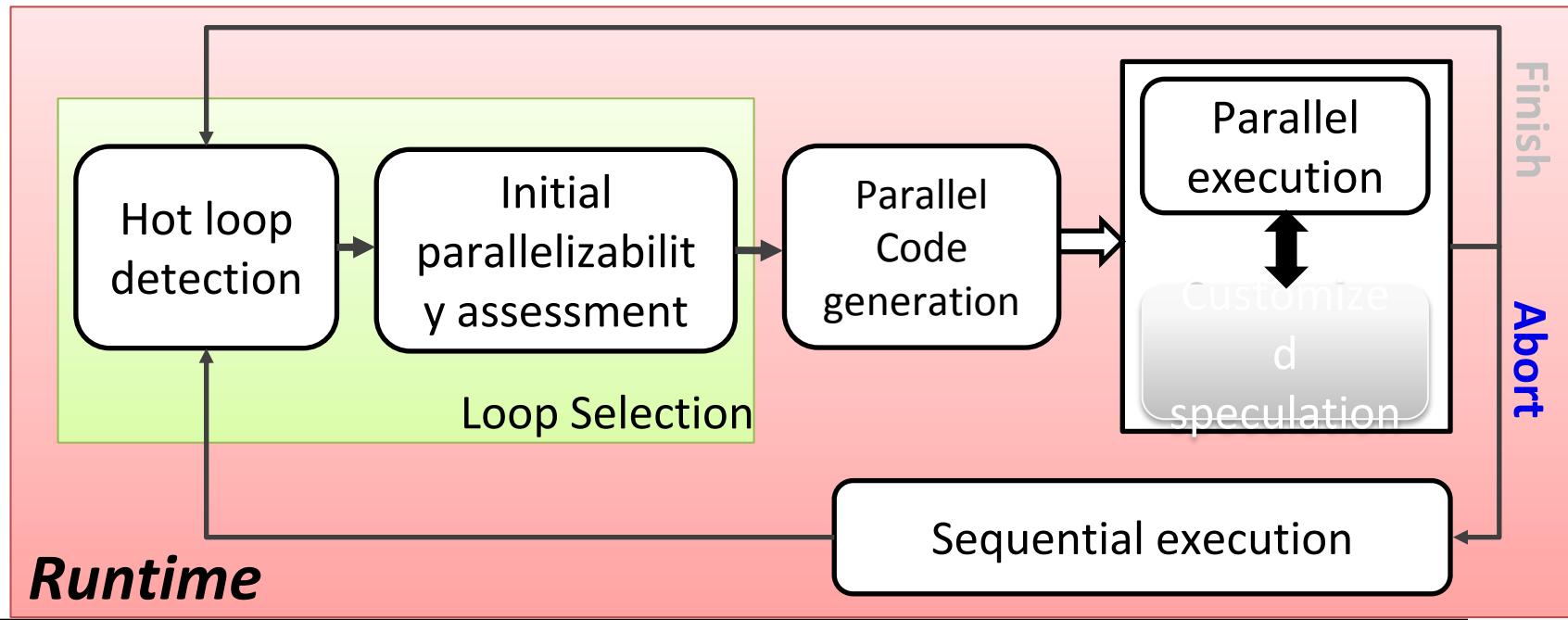
# JavaScript Parallelization

- ◆ A typical static parallelization flow



# ParaScript Approach

- ❖ Light-weight dynamic analysis & code generation for speculative DOALL loops
- ❖ Low-cost customized SW speculation with a single checkpoint



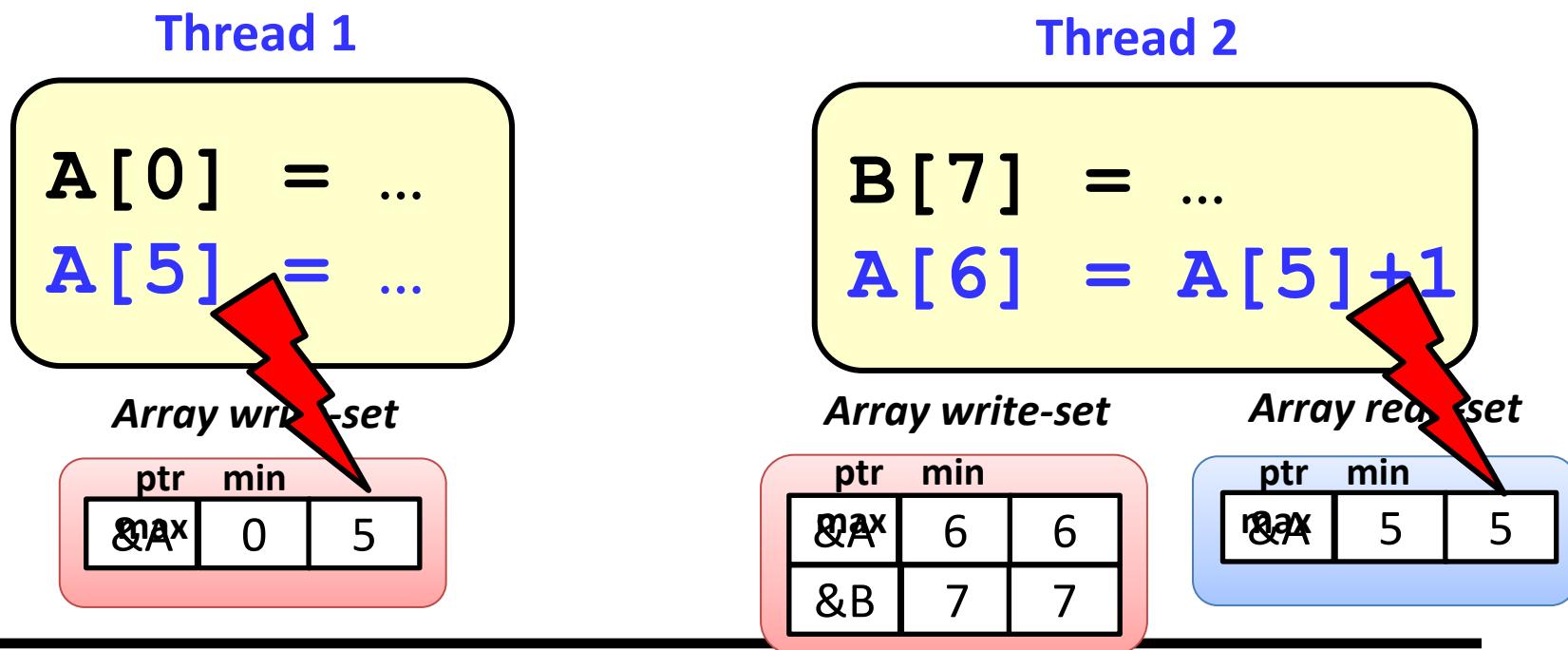
# Dependence Analysis

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Variable category	Analysis and runtime speculation method
Scalars & Objects	JIT compilation time data flow analysis
Scalar arrays	Runtime initial tests + range-based monitoring
Object arrays	Runtime reference-counting-based monitoring

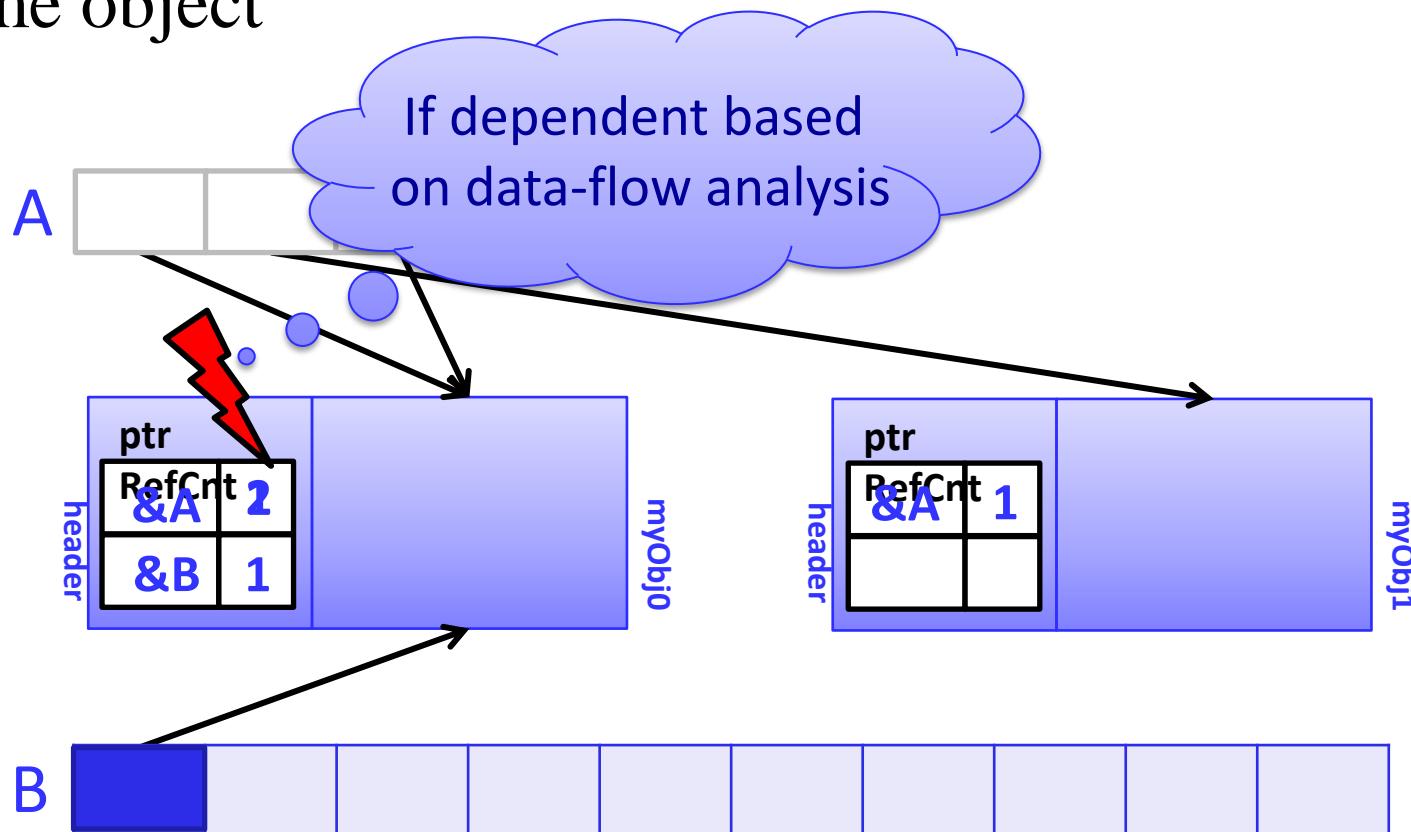
# Scalar Array Conflict Detection

- ❖ Initial assessment catches trivial conflicts
- ❖ Keep track of max and min accessed element indices
- ❖ Cross-check RD/WR sets after thread execution



# Object Array Conflict Detection

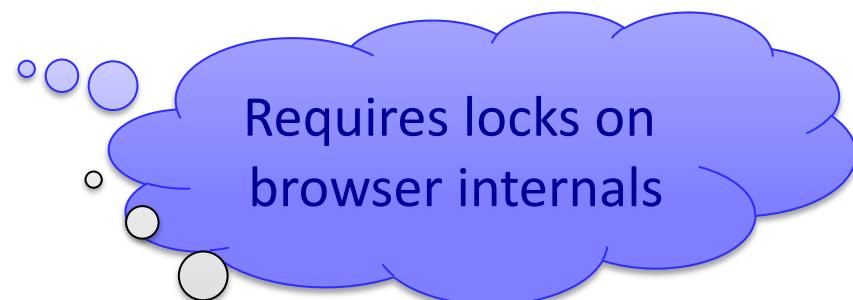
- ❖ More involved than scalar arrays
- ❖ Different indices of the same array may point to the same object



# Loop Selection

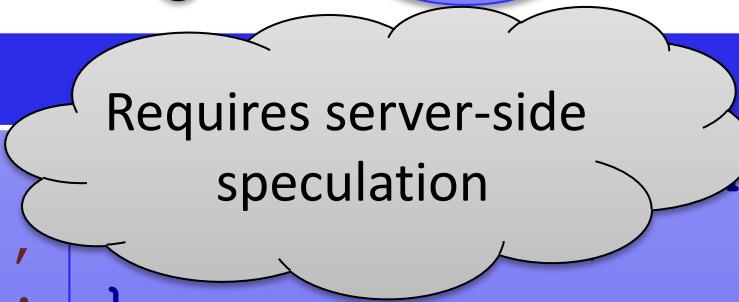
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- ❖ Focus on DOALL-counted (e.g. for loops)
- ❖ Avoid parallelizing loops with:
  - » Browser interactions
  - » HTTP request functions
  - » Runtime code insertion



## Source code

```
var addFunction =  
    new Function("a" , "b",  
                "return a+b;");  
  
eval("a = 7; b = 13;  
document.write(a+b);");
```



A grey cloud-shaped callout bubble with white outlines. Inside the bubble, the text "Requires server-side speculation" is written in black.

```
    }  
  
a = 7; b = 13;  
document.write(a+b);
```

# Checkpointing Mechanism

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- ❖ Go through all references, clone them, and ask GC not to touch the clones

Cloning process	
<b>Strings, numbers and Booleans</b>	<b>Copy the values</b>
<b>Custom objects</b>	<b>Deep-copy all properties, avoid recursion</b>
<b>Functions</b>	<ul style="list-style-type: none"><li>- Same as custom objects</li><li>- No need to clone source code</li></ul>
<b>Arrays</b>	<b>Clone all properties and elements</b>
<ul style="list-style-type: none"><li>• <b>Monitor overhead, back out if more expensive than a threshold</b></li></ul>	

# Checkpointing Optimizations

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- ❖ Selective variable cloning
  - » Only clone a variable if it is touched during speculative execution
- ❖ Array clone elimination
  - » Large arrays holding results of browser functions
  - » Instead of cloning the array, just call the function again for recovery
    - E.g. `getImageData` in the `canvas` HTML5 element

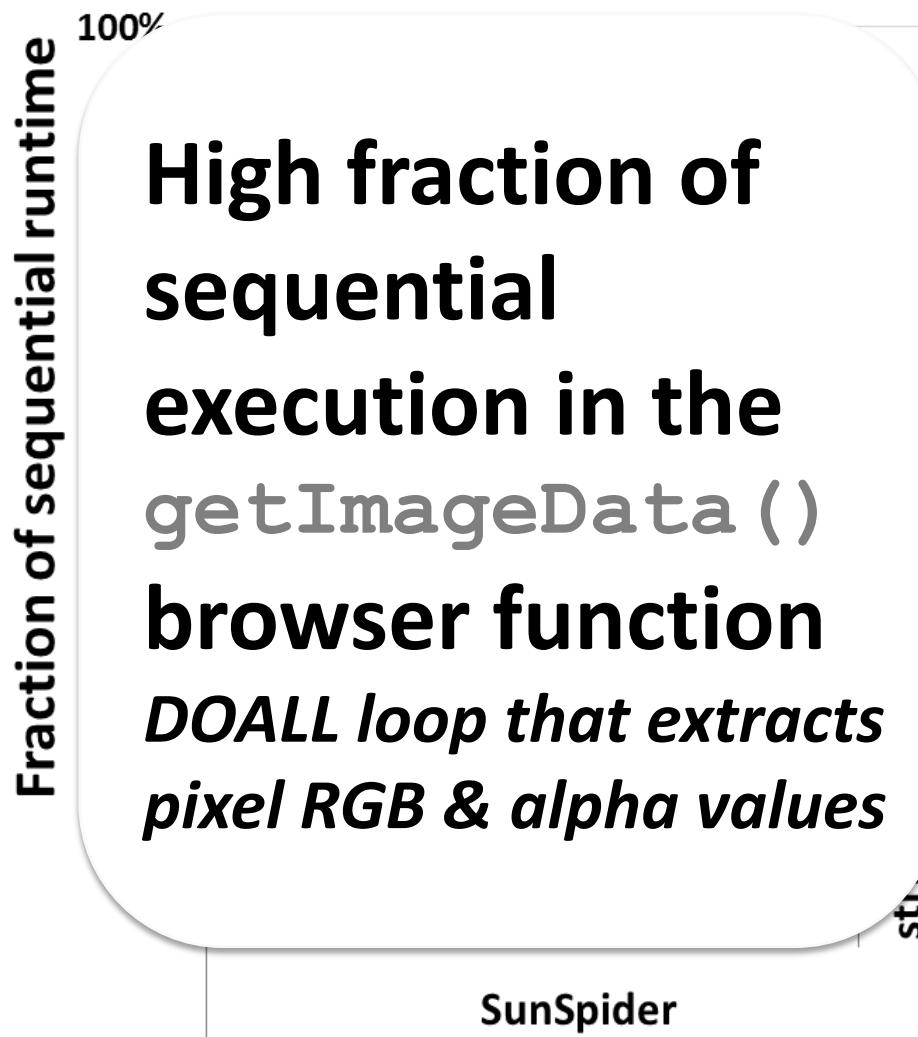
# Experimental Setup

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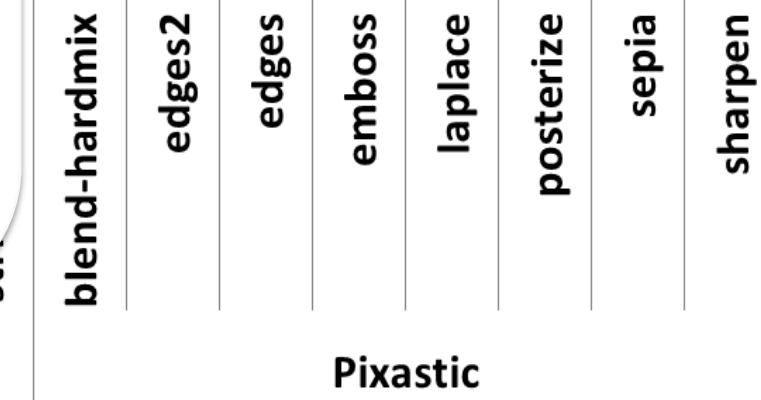
- ❖ Implemented in Firefox 3.7a1pre
- ❖ Subset of SunSpider benchmark suite
  - » Others identified as not parallelizable early on, causing 2% slow-down due to the initial analysis.
- ❖ A set of Pixastic Image Processing filters
- ❖ 8-processor system -- 2 Intel Xeon Quad-cores, running Ubuntu 9.10
- ❖ Ran each benchmark 10 times and took the average

# Parallelism Coverage

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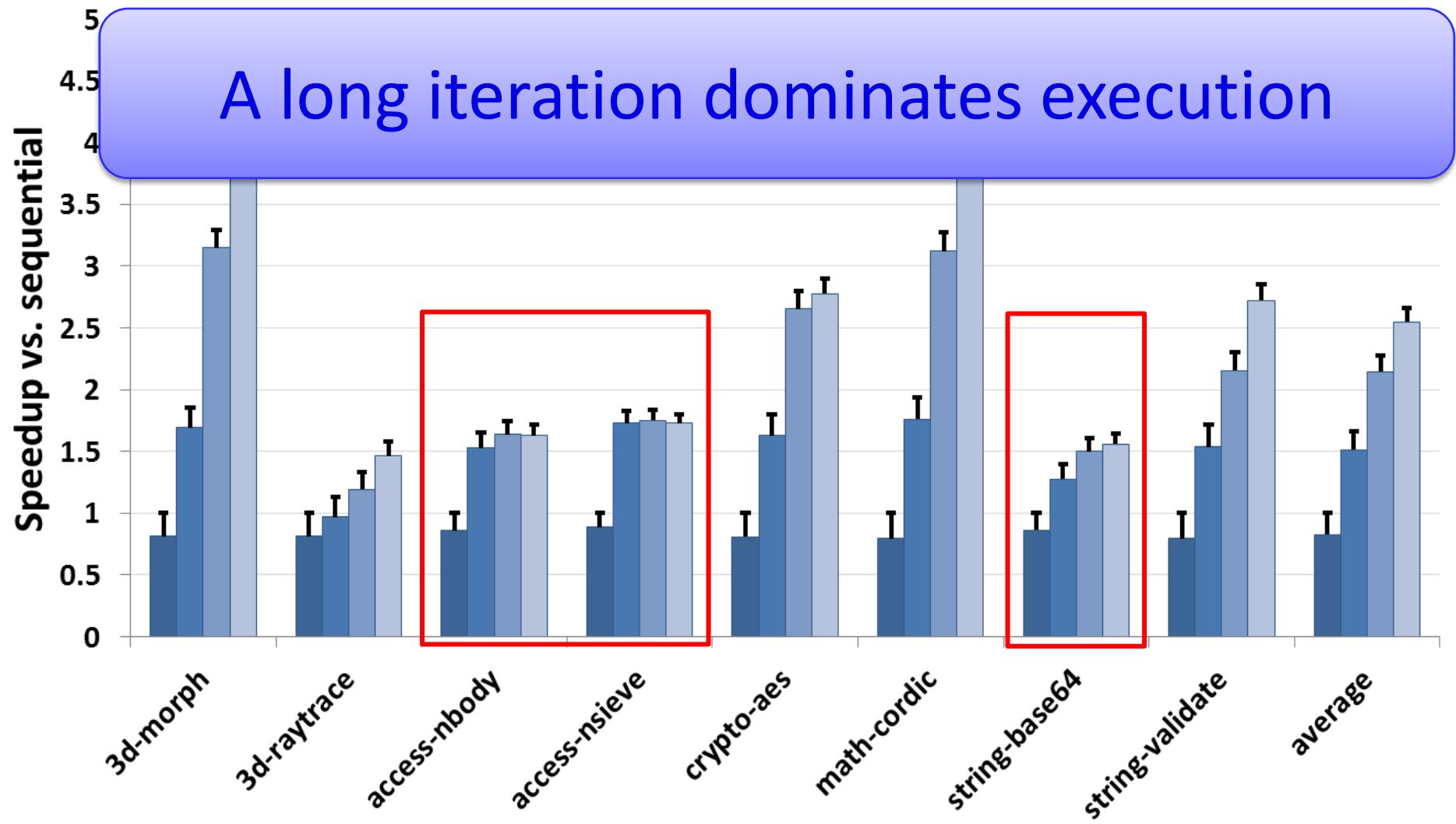


Parallelization Coverage



# SunSpider

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# Pixastic Image Processing

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