

EECS 482
**Introduction to Operating
Systems**

Winter 2018

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Naming and directories

- How to specify file to be accessed?
 - File name, click on icon, or describe contents
- File name is usually hierarchical
 - E.g., /home/barisk/482/notes
 - Allows users to group related files into one folder
 - Allows easy searching, e.g., “ls /home/barisk/482”
- Must translate file name to disk block # of header
 - What data structure to use to store mapping?
 - A hash table?
 - » Doesn't naturally represent directories
 - Tree of directories

Directories

- **Directory: mapping information for a set of files**
 - Name of file → file header's disk block # for that file
 - Often a simple array of (name, file header's disk block #) entries
- **Directories are stored on disk**
- **Directories and files are largely equivalent**
 - Same storage structure
 - Directory entry can point to file or directory
- **Any differences between files and directories?**
 - User files are relevant to the user, OS files to the OS

Example:

/home/barisk/482/notes

1. Read file header for / (root directory)
 - ◆ Contains pointers to data blocks of / directory
2. Read data blocks of /
 - ◆ Contains list of the files and directories in /. For each entry, contains mapping from name -> header's disk block #
 - ◆ One of those entries is "home"
3. Read file header for /home
4. Read data blocks for /home
5. Read file header for /home/barisk
6. Read data blocks for /home/ barisk
7. Read file header for /home/ barisk /482
8. Read data blocks for /home/ barisk /482
9. Read file header for /home/ barisk /482/notes
10. Read first data block for /home/ barisk /482/notes

Eliminated by
caching file header
for current
working directory

Unified view of multiple storage devices

- Combine multiple storage devices into a file system
 - Each device contains own file system (starting with its root)
 - A directory entry can point to the root of a different device
- Example: loginlinux.engin.umich.edu
 - / (root)
 - bin (same device as /)
 - etc (same device as /)
 - tmp (separate storage device)
 - afs (network storage “device”)
- Directory now can map name to 1) file, 2) directory, or 3) device

File caching

- File systems store lots of data structures on disk
 - Data blocks
 - Directories
 - File headers (inodes) and indirect blocks
 - Free lists
- How to improve performance?
 - Data layout to minimize seeking overhead
 - Cache data in memory
- Should the file cache be in kernel's virtual address space or in physical memory?
 - Either is fine, but if virtual, pin it

File cache vs. Virtual memory

- Both use physical memory as a cache for disk
 - **Virtual memory**: Use disk for increased capacity
 - **File systems**: Use memory for faster performance
- File cache and virtual memory compete for physical memory
 - Local vs. global replacement
 - **Why have two mechanisms that both cache disk data in memory?**

Memory-mapped files

- Use the paging system to cache both virtual address space **and** disk
 - Map file into a virtual address space
 - Point the backing store for that part of the address space at the file's data blocks
- Example: How to load a program executable from disk to memory?

File cache design

- Normal design issues for caches
 - e.g., cache size, block size, replacement policy
- How to keep copies on disk and in memory consistent with each other?
- Two options:
 - Write through
 - Write back
- Pros and cons?
- What guarantees does your Project 3 give?

Project 3 due in a week

- When we said “defer work”...
- **Assert! Assert! Assert!**
 - E.g. that page table is consistent with software page table
- **Test multi-process cases**
 - E.g. `./myrandomtest.sh`
 - » Find a random `test_*` file and run it

Project 4

- Secure, multi-threaded network file server
 - Network programming, file systems, client-server systems, and security protocols
 - Experience writing significant concurrent program
- Start soon (3x LoC as projects 2 and 3)
 - Due Apr. 14th (final on April 23rd)

Multiple updates and reliability

- File system must ensure reliability/durability
 - Okay to lose data in address space
 - Data in file system must survive system crashes and power outages
- Challenge: Crashes in midst of multi-step updates
- Example: Transfer \$100 from Baris's account to Tia's account
 1. Deduct \$100 from Baris
 2. Add \$100 to Tia
- What happens on crash between steps 1 and 2?