

CONSTRAINED OPTIMIZATION

A general constrained optimization problem has the form

$$\begin{aligned} \min_x & f(x) \\ \text{s.t.} & g_i(x) \leq 0, \quad i=1, \dots, m \\ & h_i(x) = 0, \quad i=1, \dots, p \end{aligned}$$

where $x \in \mathbb{R}^d$. If x satisfies all the constraints,

(A) we say x is _____. Assume f is
defined on all feasible points.

Lagrangian Duality

The _____ function is

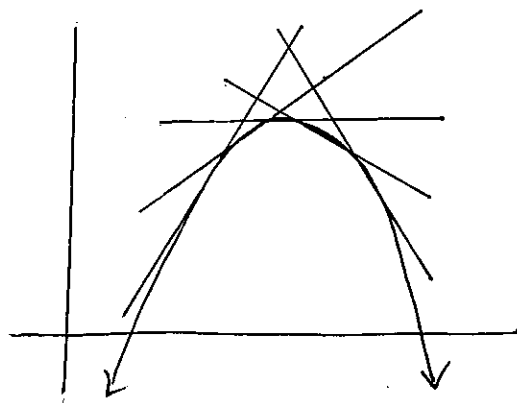
$$L(x, \lambda, \nu) := f(x) + \sum_{i=1}^m \lambda_i g_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

and $\lambda = [\lambda_1, \dots, \lambda_m]^T$ and $\nu = [\nu_1, \dots, \nu_p]^T$ are
called _____ or _____.

The (Lagrange) dual function is

$$L_D(\lambda, \nu) := \min_x L(x, \lambda, \nu)$$

Note L_D is concave, being the point-wise minimum of a family of affine functions



② Then we can define the optimization problem:

$$\max_{\lambda, \nu: \lambda_i \geq 0} L_D(\lambda, \nu)$$

Why would we constrain $\lambda_i \geq 0$?

The Primal

Similarly, we could define the function

$$L_P(x) := \max_{\lambda, \nu: \lambda_i \geq 0} L(x, \lambda, \nu)$$

and the primal optimization problem

$$\min_x L_P(x) = \min_x \max_{\lambda, \nu: \lambda_i \geq 0} L(x, \lambda, \nu)$$

It is like the dual but the min + max are swapped.

Notice that

$$L_p(x) = \begin{cases} f(x) & \text{if } x \text{ is feasible} \\ \infty & \text{else} \end{cases}$$

Thus the primal encompasses the original problem (i.e., they have the same solution), yet it is unconstrained.

Weak Duality

Claim $d^* := \max_{\lambda, \nu: \lambda_i \geq 0} \min_x L(x, \lambda, \nu)$

$$\leq \min_x \max_{\lambda, \nu: \lambda_i \geq 0} L(x, \lambda, \nu) =: p^*$$

Proof: Let \tilde{x} be feasible. Then for any λ, ν with $\lambda_i \geq 0$,

$$L(\tilde{x}, \lambda, \nu) = f(\tilde{x}) + \sum \lambda_i g_i(\tilde{x}) + \sum \nu_i h_i(\tilde{x}) \leq f(\tilde{x})$$

Hence

$$L_D(\lambda, \nu) = \min_x L(x, \lambda, \nu) \leq L(\tilde{x}, \lambda, \nu) \leq f(\tilde{x}).$$

This is true for any feasible \tilde{x} , so

$$L_D(\lambda, \nu) \leq \min_{\substack{\tilde{x} \\ \tilde{x} \text{ feasible}}} f(\tilde{x}) = p^*$$

Taking the max over $\lambda, \nu: \lambda_i \geq 0$, we have

$$d^* = \max_{\lambda, \nu: \lambda_i \geq 0} L_D(\lambda, \nu) \leq p^*$$

□

© The difference $p^* - d^*$ is called the _____.

Strong Duality

If $p^* = d^*$, we say strong duality holds.

Theorem | If the primal problem is _____ (f, g_i convex, h_i affine), and a constraint qualification holds, then $p^* = d^*$.

Examples | of constraint qualifications:

- All g_i are _____.
- $\exists x$ s.t. $h_i(x) = 0 \ \forall i$, $g_i(x) < 0 \ \forall i$
(strict feasibility)

KKT Conditions

Assume f, g_i, h_i are differentiable.

Necessity

Theorem | If $p^* = d^*$, x^* is primal optimal, and (λ^*, ν^*) is dual optimal, then the Karush-Kuhn-Tucker conditions hold:

○
KKT
conditions

$$\left[\begin{array}{l} (1) \nabla f(x^*) + \sum_{i=1}^m \lambda_i^* \nabla g_i(x^*) + \sum_{i=1}^p \nu_i^* \nabla h_i(x^*) = 0 \\ (2) g_i(x^*) \leq 0, \quad i=1, \dots, m \\ (3) h_i(x^*) = 0, \quad i=1, \dots, p \\ (4) \lambda_i^* \geq 0, \quad i=1, \dots, m \\ (5) \lambda_i g_i(x^*) = 0, \quad i=1, \dots, m \quad (\text{complementary slackness}) \end{array} \right.$$

Proof: (2)-(3) hold because x^* must be feasible, and (4) holds by def. of the dual problem. To prove (5),

$$\begin{aligned} f(x^*) &= L_D(\lambda^*, \nu^*) \quad [\text{by strong duality}] \\ &= \min_x \left(f(x) + \sum_{i=1}^m \lambda_i^* g_i(x) + \sum_{i=1}^p \nu_i^* h_i(x) \right) \\ &\leq f(x^*) + \sum_{i=1}^m \lambda_i^* g_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*) \\ &\leq f(x^*) \quad [\text{by (2)-(4)}] \end{aligned}$$

and therefore the two inequalities are equalities. Equality of the last two lines implies $\lambda_i^* g_i(x^*) = 0 \quad \forall i$, which is (5). Equality of the 2nd and 3rd lines implies x^* is a minimizer of $L(x, \lambda^*, \nu^*)$ w.r.t x .

○ Therefore $\nabla_x L(x^*, \lambda^*, \nu^*) = 0$, which is (1). \square

Sufficiency

Theorem | If the original problem is convex (i.e., f, g_i are convex functions, h_i are affine), and $\tilde{x}, \tilde{\lambda}, \tilde{\nu}$ satisfy the KKT conditions, then \tilde{x} is primal optimal, $(\tilde{\lambda}, \tilde{\nu})$ is dual optimal, and the duality gap is zero.

Proof: (2), (3) $\Rightarrow \tilde{x}$ is feasible

(4) $\Rightarrow L(x, \tilde{\lambda}, \tilde{\nu})$ is convex in x

(1) $\Rightarrow \tilde{x}$ is a minimizer of $L(x, \tilde{\lambda}, \tilde{\nu})$. Then

$$L_D(\tilde{\lambda}, \tilde{\nu}) = L(\tilde{x}, \tilde{\lambda}, \tilde{\nu})$$

$$= f(\tilde{x}) + \sum \tilde{\lambda}_i g_i(\tilde{x}) + \sum \tilde{\nu}_i h_i(\tilde{x})$$

$$= f(\tilde{x})$$

= 0 by (5)

□

Conclusion

If a constrained optimization problem is differentiable, and has convex objective function and constraint sets, then the KKT conditions are necessary and sufficient for primal/dual optimality (with zero duality gap). Thus, the KKT conditions can be used to solve such problems.

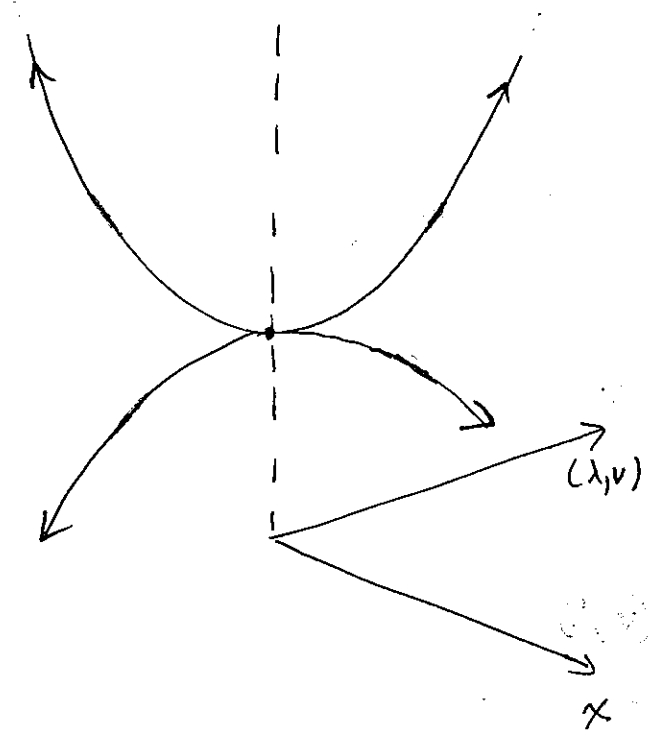
Saddle Point Property

If \tilde{x} and $(\tilde{\lambda}, \tilde{\nu})$ are primal/dual optimal w/ zero duality gap, they are a saddle point of L , i.e.

$$L(\tilde{x}, \lambda, \nu) \leq L(\tilde{x}, \tilde{\lambda}, \tilde{\nu}) \leq L(x, \tilde{\lambda}, \tilde{\nu})$$

for all $x, (\lambda, \nu)$.

Justification of this fact is left as an exercise.



Key

A. feasible; Lagrangian;

Lagrange multipliers or dual variables

B. dual; primal C. duality gap; convex; affine

Reference

Boyd & Vandenberghe, Convex Optimization, ch 5
(available online)