

TOPICS FOR TODAY'S LECTURE

FAST FOURIER TRANSFORM (FFT)

1. Derivation using \mathcal{Z} xform
2. Derivation: General case
3. Radix-2 FFT ($N=2^n$)
 - a. Decimation-in-time
 - b. Decimation-in-freq.
4. Small radix-2 example

INTRODUCTION [1/2]

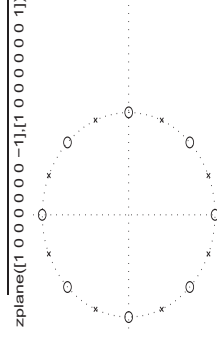
Goal: Compute DFT $X_k = \sum_{n=0}^{N-1} x[n]W_N^{nk}$
where: $W_N = e^{-j2\pi/N}$ (makes formulae neater).
FFT: *Algorithm* for computing;
DFT: A *transform* $x[n] \rightarrow X_k$.
Direct: N^2 Mults. and Adds (MADs).
History: Gauss; Cooley-Tukey (1965).
Note: DSP dates from this algorithm.

INTRODUCTION [2/2]

Idea: *Divide* large problem into smaller;
Then: *Conquer* by solving smaller ones.
EX: 1000-point problem: $(1000)^2=1$ million ops.
But: 100 10-point problems: $100(10)^2=10,000$ ops.
But: Great! But how can we do this?

z-TRANSFORM DERIVATION [1/2]

Recall: $X_k = X(z)|_{z=e^{j2\pi k/N}}$.
That is: Evaluate $X(z)$ at:



z-TRANSFORM DERIVATION [2/2]

Idea: $X_{2k} = X(z)|_{z=e^{j2\pi k/N}} = X(z)|_{z=e^{j2\pi k/(N/2)}}$.
and: $X_{2k+1} = X(z)|_{z=e^{j2\pi(2k+1)/N}} = X(z)|_{z=e^{j2\pi/N} e^{j2\pi k/(N/2)}}$.
That is: Even indices: $\frac{N}{2}$ -point DFT of $x[n] + x[n + \frac{N}{2}]$ (aliased)
and: Odd indices: $\frac{N}{2}$ -point DFT of $e^{j2\pi n/N} (x[n] - x[n + \frac{N}{2}])$ (modulate)
Why: *Aliased:* see HW. The *-sign:* $e^{j2\pi(n+(N/2)/N)} = -e^{j2\pi n/N}$.

GENERAL DERIVATION [1/5]

Goal: $(N_1 N_2)$ -point DFT \rightarrow
 N_2 (N_1) -point & N_1 (N_2) -point).
DEF: Coarse and vernier indices:
Coarse: $n = n_1 + N_1 n_2$ where $\begin{matrix} n_1=0\dots N_1-1 \\ n_2=0\dots N_2-1 \end{matrix}$.
Vernier: $k = k_2 + N_2 k_1$ where $\begin{matrix} k_1=0\dots N_1-1 \\ k_2=0\dots N_2-1 \end{matrix}$.
Why: Coarse and fine microscope controls.

GENERAL DERIVATION [2/5]

$$\begin{aligned} \text{nk} &= (n_1 + N_1 n_2)(k_2 + N_2 k_1) \\ &= n_1 k_2 + N_1 n_2 k_2 + N_2 n_1 k_1 + N_1 N_2 n_2 k_1. \end{aligned}$$

$$e^{-j2\pi(\cdot)}: W_N^{nk} = W_N^{n_1 k_2} W_{N_2}^{n_2 k_2} W_{N_1}^{n_1 k_1}$$

using: $W_{N_1 N_2}^{N_1} = W_{N_2}$ and $W_{N_1 N_2}^{N_1 N_2} = 1$.

GENERAL DERIVATION [3/5]

$$\text{DFT: } X_k = \hat{X}_{k_2 + N_2 k_1} = \sum_{n_1=0}^{N_1-1} \sum_{n_2=0}^{N_2-1} x[n_1 + N_1 n_2] W_N^{(n_1 + N_1 n_2)(k_2 + N_2 k_1)}$$

$$\text{Rewrite: } X_k = \hat{X}_{k_2 + N_2 k_1} = \underbrace{\sum_{n_1=0}^{N_1-1} W_{N_1}^{n_1 k_1}}_{N_1\text{-point DFT}} \left[\underbrace{W_N^{n_1 k_2} \sum_{n_2=0}^{N_2-1} W_{N_2}^{n_2 k_2} x[n_1 + N_1 n_2]}_{N_2\text{-point DFT for each } n_1} \right].$$

GENERAL DERIVATION [4/5]

1. Compute N_1 (for each n_1) N_2 -point DFTs of $x[n_1 + N_1 n_2]$.
2. Multiply result by *twiddle factors* $W_N^{n_1 k_2}$. Called *twiddle mults*.
3. Compute N_2 (for each k_2) N_1 -point DFTs of the result. Done!

$(N_1 N_2)$ -point $\rightarrow N_1(N_2\text{-point}) + N_2(N_1\text{-point}) + (N_1 - 1)(N_2 - 1)$ twiddle.
since: $n_1 = 0$ or $k_2 = 0 \rightarrow W_N^{n_1 k_2} = 1$. This is important below.

RADIX-2 COOLEY-TUKEY FFT

1. *Decimation-in-time:* $N = 2 \frac{N}{2}$.

Set: $N_1 = 2; N_2 = N/2; n_1 = 0, 1; k_1 = 0, 1$.

2. *Decimation-in-freq:* $N = \frac{N}{2} 2$.

Set: $N_1 = N/2; N_2 = 2; n_2 = 0, 1; k_2 = 0, 1$.

Then: $N\text{-point} \rightarrow 2(\frac{N}{2}\text{-point}) + \frac{N}{2}(2\text{-point})$
 $+ (\frac{N}{2} - 1)$ twiddle mults. **Total:** $\frac{N}{2} \log_2 N$ mults.

Note: $(N_1 - 1)(N_2 - 1) = (2 - 1)(\frac{N}{2} - 1) = \frac{N}{2} - 1$;
That is: Half the twiddle mults are trivial!

GENERAL DERIVATION [5/5]

Visual: $(N_1 \times N_2)$ arrays: $x_{n_1, n_2} = x[n_1 + N_1 n_2]$; $X_{k_1, k_2} = \hat{X}_{k_2 + N_2 k_1}$.

1. Take N_2 -point DFT of each row (fixed n_1). Yields \hat{x}_{n_1, k_2} .
2. Multiply \hat{x}_{n_1, k_2} point-by-point by the twiddle factor $W_N^{n_1 k_2}$.
3. Take N_1 -point DFT of each column (fixed k_2). Yields X_{k_1, k_2} .

Example: Later in lecture.

DEC.-IN-TIME RADIX-2 COOLEY-TUKEY FFT

$$X_{k_2} = 1 \sum_{n_2=0}^{N/2-1} \underbrace{W_{N/2}^{n_2 k_2} x[2n_2]}_{N/2\text{-point DFT}} + W_N^{1k_2} \sum_{n_2=0}^{N/2-1} \underbrace{W_{N/2}^{n_2 k_2} x[2n_2 + 1]}_{N/2\text{-point DFT}}.$$

$$X_{k_2 + N/2} = 1 \sum_{n_2=0}^{N/2-1} \underbrace{W_{N/2}^{n_2 k_2} x[2n_2]}_{N/2\text{-point DFT}} - W_N^{1k_2} \sum_{n_2=0}^{N/2-1} \underbrace{W_{N/2}^{n_2 k_2} x[2n_2 + 1]}_{N/2\text{-point DFT}}.$$

DEC-IN-FREQ RADIX-2 COOLEY-TUKEY FFT

$$X_{2k_1} = \sum_{n_1=0}^{N/2-1} \underbrace{W_{N/2}^{n_1 k_1} [x[n_1] + x[n_1 + N/2]]}_{\text{Half of 2-point DFT}} \cdot 1.$$

$$X_{2k_1+1} = \sum_{n_1=0}^{N/2-1} \underbrace{W_{N/2}^{n_1 k_1} [x[n_1] - x[n_1 + N/2]]}_{\text{Half of 2-point DFT}} W_N^{1n_1}.$$

COMPUTING \mathcal{F} USING DFT [1/3]

Goal: Compute numerically

$$\mathcal{F}: X(f) = \int x(t) e^{-j2\pi f t} dt.$$

$$\mathcal{F}^{-1}: x(t) = \int X(f) e^{j2\pi f t} df.$$

Note: $X(f) = X(\frac{\omega}{2\pi})$ & $df = \frac{d\omega}{2\pi}$.

Assume: $x(t)$ time-limited to $0 < t < T$.

Assume: $X(f)$ band-limited to $|f| < B/2$.

COMPUTING \mathcal{F} USING DFT [3/3]

Let: $X_k = X(k\Delta_f) \frac{1}{N} = X(k\Delta_f) B$

and: $\Delta_t \Delta_f = \frac{1}{N}$ and $N = BT$.

Then: $X_k \approx \sum_{n=0}^{N-1} x(n\Delta_t) e^{-j2\pi n k / N}$.

and: $x(n\Delta_t) \approx \frac{1}{N} \sum_{k=0}^{N-1} X_k e^{j2\pi n k / N}$.

Formulae: $BT = N$; $\Delta_t \Delta_f = 1/N$ $x[n] = X(n\Delta_t)$.

Formulae: $\Delta_t = 1/B$; $\Delta_f = 1/T$; $X_k = X(k\Delta_f) B$.

$$1. \begin{bmatrix} x[0] \\ x[2] \end{bmatrix} \begin{bmatrix} x[1] \\ x[3] \end{bmatrix} = \begin{bmatrix} a & b \\ c & d \end{bmatrix} = \begin{bmatrix} a & b \\ c & d \end{bmatrix} \rightarrow \begin{bmatrix} (a+c) \\ (a-c) \end{bmatrix} \begin{bmatrix} (b+d) \\ (b-d) \end{bmatrix}$$

$$2. \begin{bmatrix} (a+c) \\ (a-c) \end{bmatrix} \begin{bmatrix} (b+d) \\ (b-d) \end{bmatrix} \times \begin{bmatrix} 1 & 1 \\ 1 & -j \end{bmatrix} = \begin{bmatrix} (a+c) & (b+d) \\ (a-c) & -j(b-d) \end{bmatrix}$$

$$3. \begin{bmatrix} [(a+c) + (b+d)] \\ [(a-c) - j(b-d)] \end{bmatrix} \begin{bmatrix} [(a+c) - (b+d)] \\ [(a-c) + j(b-d)] \end{bmatrix} = \begin{bmatrix} X_0 & X_2 \\ X_1 & X_3 \end{bmatrix}$$

4. Read off $\{X_0, X_1, X_2, X_3\}$. Note index shuffling.

COMPUTING \mathcal{F} USING DFT [2/3]

Sample t : $t = n\Delta_t$ for $0 \leq n \leq N-1$.

Sample f : $f = k\Delta_f$ for $0 \leq k \leq N-1$.

Nyquist: $\Delta_t = 1/B$ and $\Delta_f = 1/T$.

Approx: $X(f) \approx \sum_{n=0}^{N-1} x(n\Delta_t) e^{-j2\pi f n \Delta_t} \Delta_t$.

Sample f : $X(k\Delta_f) \approx \sum_{n=0}^{N-1} x(n\Delta_t) e^{-j2\pi n k \Delta_t \Delta_f} \Delta_t$.

and: $x(n\Delta_t) \approx \sum_{k=0}^{N-1} X(k\Delta_f) e^{j2\pi n k \Delta_t \Delta_f} \Delta_f$.

COMPUTING \mathcal{F} USING DFT: EXAMPLE

EX: $\mathcal{F}\{e^{-t}\} = \frac{2}{\omega^2+1} = \frac{2}{4\pi^2 f^2+1}$.

Limits: $-6 < t < 6 \rightarrow T = 12$

and: $-8 < f < 8 \rightarrow B = 16$.

Values: $N = BT = (16)(12) = 192$;

and: $\Delta_t = \frac{1}{B} = \frac{1}{16}$; $\Delta_f = \frac{1}{T} = \frac{1}{12}$.

FFT: 192-point FFT of $x(t)$ sampled $t = n/16$

gives: 16X(f) sampled at $f = k/12$.

Result: Actual & computed coincide.