

- DFT:** $X_k = \sum_{n=0}^{N-1} x(n)e^{-j2\pi nk/N}$; $x(n) = \frac{1}{N} \sum_{k=0}^{N-1} X_k e^{j2\pi nk/N}$; length $\leq N$.
DTFT: $X_k = X(z)|_{z=e^{j2\pi k/N}} = X(e^{j\omega})|_{\omega=2\pi k/N}$ sampled on unit circle.
DTFS: $\mathcal{F}\{\text{periodic extension of } \{x(n)\}\} = \text{periodic extension of } \{X_k\}$.
except: factor of $1/N$ moved. **Note:** $k = 0, 1 \dots N - 1$ and $n = 0, 1 \dots N - 1$.
Text: Uses $X(k)$ not X_k . I *hate* that—too easy to confuse with $X(z)$!

- What:** Use DFT to compute DTFT at $\omega = \frac{2\pi k}{N}$: equispaced samples on u.c.
Why: Use for spectral analysis, and to recover $x(n)$ from DTFT. BUT:
Finite $x(n)$ has finite length $L \rightarrow$ using $N \geq L \rightarrow$ recover $x(n)$ from X_k .
length: But $N < L \rightarrow$ can only recover $\sum_k x(n + kN) \rightarrow$ *aliased* $x(n)$.

- EX:** $x(n) = \{1, 2, 3, 4, 5\} \rightarrow X(e^{j\omega}) = 1 + 2e^{-j\omega} + 3e^{-2j\omega} + 4e^{-j3\omega} + 5e^{-j4\omega}$
 Sample DTFT on unit circle at $\omega = 0, \frac{\pi}{2}, \pi, \frac{3\pi}{2}$ to get DFT:
DFT: $X_0 = 1 + 2 + 3 + 4 + 5 = 15$; $X_1 = 1 - 2j - 3 + 4j + 5 = 3 + 2j$;
N=4: $X_2 = 1 - 2 + 3 - 4 + 5 = 03$; $X_3 = 1 + 2j - 3 - 4j + 5 = 3 - 2j = X_1^*$.
IDFT: $x(0) = \frac{1}{4}[15 + (3 + 2j) + 3 + (3 - 2j)] = 6$ incorrect ($6=1+5$ aliased).
 $x(1) = \frac{1}{4}[15(1) + (3 + 2j)(j) + 3(-1) + (3 - 2j)(-j)] = 2$ correct.
 $x(2) = \frac{1}{4}[15(1) + (3 + 2j)(-1) + 3(1) + (3 - 2j)(-1)] = 3$ correct.
 $x(3) = \frac{1}{4}[15(1) + (3 + 2j)(-j) + 3(-1) + (3 - 2j)(j)] = 4$ correct.
Why? *Undersampled* $X(e^{j\omega})$ on unit circle \rightarrow aliasing (just like before).

- Inter-** Let $x(n)$ have length $L \leq N$. We're given $X(e^{j2\pi k/N}), k = 0, 1 \dots N-1$.
polate Then interpolate $X(e^{j\omega}) = \sum_{k=0}^{N-1} X(e^{j2\pi k/N})S(\omega - \frac{2\pi k}{N})$
DTFT: $S(\omega) = \frac{\sin(\omega N/2)}{N \sin(\omega/2)} e^{-j\omega(N-1)/2} = DTFT\{\frac{1}{N} \sum_{i=0}^{N-1} \delta(n - i)\}$.

- Zero:** $N > L \rightarrow DFT\{x(0) \dots x(L-1), 0 \dots 0(\text{zero-pad})\} \rightarrow$ above.
padding This smoothes DTFT (finer sampling in ω), BUT: no additional info.
EX: $DFT_N\{\sum_{i=0}^{L-1} \delta(n - i)\} = \frac{\sin(\pi k L/N)}{\sin(\pi k/N)} e^{-j\pi k(L-1)/N}$. Window blurs.

- Cyclic** $y(n) = h(n) \odot x(n) = \sum_{i=0}^{N-1} h(i)x((n-i))_N \Leftrightarrow Y_k = H_k X_k$.
convol: $(x(n))_N \Leftrightarrow$ N-point periodic extension of $x(n)$. "Cyclic" = "circular."
Ex: $\{1, 2, 3\} \odot \{4, 5, 6\}$: Take one complete cycle of each:
 $y(0) = \begin{matrix} 1,2,3,1,2,3 \\ 4,6,5,4,6,5 \end{matrix} = 1 \cdot 4 + 2 \cdot 6 + 3 \cdot 5 = 31$.
 $y(1) = \begin{matrix} 1,2,3,1,2,3 \\ 5,4,6,5,4,6 \end{matrix} = 1 \cdot 5 + 2 \cdot 4 + 3 \cdot 6 = 31$.
 $y(2) = \begin{matrix} 1,2,3,1,2,3 \\ 6,5,4,6,5,4 \end{matrix} = 1 \cdot 6 + 2 \cdot 5 + 3 \cdot 4 = 28$.
Check: $H_0 X_0 = (1 + 2 + 3)(4 + 5 + 6) = 90 = (31 + 31 + 28) = Y_0$ checks.

Goal: Compute DFT $X_k = \sum_{n=0}^{N-1} x(n)W_N^{nk}$ where $W_N = e^{-j2\pi/N}$.

Names: **FFT** is an *algorithm* for computing the **DFT**, which is a *transform*.

Why? Direct computation requires N^2 mults and $N(N-1)$ adds: too many!

Cooley- 1965 at IBM. Serious DSP dates from this algorithm.

Tukey: Divide up large DFT into smaller DFTs: $N = N_1N_2$.

coarse: $n = n_1 + N_1n_2$ where $n_1 = 0 \dots N_1 - 1$ and $n_2 = 0 \dots N_2 - 1$.

vernier: $k = k_2 + N_2k_1$ where $k_1 = 0 \dots N_1 - 1$ and $k_2 = 0 \dots N_2 - 1$.

indices $nk = (n_1 + N_1n_2)(k_2 + N_2k_1) = n_1k_2 + N_1n_2k_2 + N_2n_1k_1 + N_1N_2n_2k_1$.

exponent: $W_N^{nk} = W_N^{n_1k_2} W_{N_2}^{n_2k_2} W_{N_1}^{n_1k_1}$ using $W_{N_1N_2}^{N_1} = W_{N_2}$ and $W_N^{N_1N_2} = 1$.

DFT: $X_k = X_{k_2+N_2k_1} = \sum_{n_1=0}^{N_1-1} \sum_{n_2=0}^{N_2-1} x(n_1 + N_1n_2)W_N^{(n_1+N_1n_2)(k_2+N_2k_1)}$.

rewrite: $X_{k_2+N_2k_1} = \sum_{n_1=0}^{N_1-1} W_{N_1}^{n_1k_1} \left[W_N^{n_1k_2} \sum_{n_2=0}^{N_2-1} W_{N_2}^{n_2k_2} x(n_1 + N_1n_2) \right]$.

1. Compute N_1 (for each n_1) N_2 -point DFTs of $x(n_1 + N_1n_2)$ (in n_2).
2. Multiply result by *twiddle factors* $W_N^{n_1k_2}$. These are *twiddle mults*.
3. Compute N_2 (for each k_2) N_1 -point DFTs of the result. Now done!
4. (N_1N_2) -point $\rightarrow N_1(N_2\text{-point}) + N_2(N_1\text{-point}) + (N_1-1)(N_2-1)$ twiddle since $n_1 = 0$ or $k_2 = 0 \rightarrow W_N^{n_1k_2} = 1$. This is important below.

Visual: $(N_1 \times N_2)$ arrays: $x_{n_1,n_2} = x(n_1 + N_1n_2)$; $X_{k_1,k_2} = X(k_2 + N_2k_1)$.

1. Take N_2 -point DFT of each row (fixed n_1). Yields \hat{x}_{n_1,k_2} .
2. Multiply \hat{x}_{n_1,k_2} point-by-point by twiddle factor $W_N^{n_1k_2}$.
3. Take N_1 -point DFT of each column (fixed k_2). Yields X_{k_1,k_2} .

Radix-2 Cooley-Tukey Fast Fourier Transforms

$N = 2\frac{N}{2} \rightarrow$ *Decimation-in-time* FFT: $N_1 = 2$; $N_2 = N/2$; $n_1 = 0, 1$; $k_1 = 0, 1$.

$$X_{k_2} = 1 \sum_{n_2=0}^{N/2-1} W_{N/2}^{n_2k_2} x(2n_2) + W_N^{1k_2} \sum_{n_2=0}^{N/2-1} W_{N/2}^{n_2k_2} x(2n_2 + 1).$$

$$X_{k_2+N/2} = 1 \sum_{n_2=0}^{N/2-1} W_{N/2}^{n_2k_2} x(2n_2) - W_N^{1k_2} \sum_{n_2=0}^{N/2-1} W_{N/2}^{n_2k_2} x(2n_2 + 1).$$

N-point $\rightarrow 2(\frac{N}{2}\text{-point}) + \frac{N}{2}(2\text{-point}) + (\frac{N}{2} - 1)$ twiddle mults. Total: $\frac{N}{2} \log_2 N$ mults.

Note: $(N_1 - 1)(N_2 - 1) = (2 - 1)(\frac{N}{2} - 1) = \frac{N}{2} - 1 : \frac{1}{2}$ twiddle mults trivial!

$N = \frac{N}{2}2 \rightarrow$ *Decimation-in-freq.* FFT: $N_2 = 2$; $N_1 = N/2$; $n_2 = 0, 1$; $k_2 = 0, 1$.

$$X_{2k_1} = \sum_{n_1=0}^{N/2-1} W_{N/2}^{n_1k_1} [x(n_1) + x(n_1 + N/2)]1.$$

$$X_{2k_1+1} = \sum_{n_1=0}^{N/2-1} W_{N/2}^{n_1k_1} [x(n_1) - x(n_1 + N/2)]W_N^{1n_1}.$$

N-point $\rightarrow \frac{N}{2}(2\text{-point}) + 2(\frac{N}{2}\text{-point}) + (\frac{N}{2} - 1)$ twiddle mults. Total: $\frac{N}{2} \log_2 N$ mults.