Multi-stage replay with Crosscut

Jim Chow, Dominic Lucchetti, Tal Garfinkel, Geoffrey Lefebvre, Ryan Gardner, Joshua Mason, Sam Small VMware Peter M. Chen University of Michigan

Analyzing with records

- Detailed recordings of program execution
 - Program output, logs
 - syslogd
 - Ad-hoc records ("printf" debugging)
 - Backtraces
 - Core dumps
- Useful for analyzing programs

Correctness, performance, security

Debugging programs with records

• So useful, now commonplace to send records back to the developers:



Debugging programs with records

- Despite extensive collection of records, programs still hard to debug
 - Program output, logs
 - syslogd
 - Ad-hoc records ("printf" debugging)
 - Backtraces
 - Core dumps
- Don't know everything about a program
- Miss behavior: don't run all the time, or over all values---must decide *when* or *what* to record (balancing act)

Instead: record everything, all the time

- Easy to do with a VMM
- VMM: thin layer between software stack and HW
- All inputs into system must pass through the VMM
 - All possible decisions governed only by these inputs



• Easy to do with a VMM • VMM: thin layer between software stack and HW Application • All inputs into system must pass through the VMM **Operating system** All possible decisions governed only by these inputs Virtualization layer • If we save inputs and their timing have all non-determinism we need

• Easy to do with a VMM • VMM: thin layer between software stack and HW Application All inputs into system must pass through the VMM **Operating system** All possible decisions governed only by these inputs Virtualization layer • If we save inputs and their timing have all non-determinism we need • To replay: Start from clean state







- But have to replay to reclaim state:
 - Need a replayer.
 - Input not understandable
 - Need to feed all inputs to replayer to get any useful output.
 - Won't work with subset.



Goal: replay with discretion

• VM recordings are great.

Provide great coverage.

- But we want to replay with discretion:
 - Don't require/expose all inputs (no VM).
 - Don't require same abstractions (no VMM).
 - Don't sacrifice *coverage* for discretion.
- How are we going to do this?

VM replay is a state machine



- Unhappy one size fits all
- Convert inputs to a new machine

- Represents some computation in original

Generate abstractions

- Know we can generate new machine for any abstract state, computation in the log:
 - Program output, logs
 - syslogd
 - Ad-hoc records ("printf" debugging)
 - Backtraces
 - Core dumps

• VM recordings subsume all of the above.

Any data we need can be generated.

In other words:



In other words:









level state machine.

What are some useful state machines?



- Motivator:
 - Replay debugging for Dr. Watson-type crash



This machine is weird



- Motivator:
 - Replay for Dr. Watson-type crash dumps
 - Requires whole VM
 - Fine within an organization
 - Silly otherwise

A better machine



- Generate inputs for new abstract machine
 - Replays just a process
- Given inputs, cannot generate state for other processes
 - Better than original machine

How to represent process state



- Retain benefits of original system:
 - Record all input to a process
 - Including timing, races, shared memory accesses (*)
 - (*) at least for VMs that are virtual uniprocessors.
 - Agnostic to application, OS

Choose a representation independent of OS



- Represented with HW state.
 - Address space, registers, memory map, no OS state.
- Given representation, runs deterministic
 - Until needs input: syscall, signals, CPU counters, etc.
 - Input's only effect: modify represented state
 - Continually supply this state.

Figure out where state comes from

Normal system calls, traps, interrupts



- Only comes from 3 places
- Modified VMM
 - Hook sites of introduction
 - Control flow into/out of

process

System calls, traps, interrupts when slicing a process



Collect input without OS help

System calls, traps, interrupts when slicing a process



- OS injects state (directly on indirectly)
 - Do it with MMU: agnostic to OS.
 - Notice modifications to sliced process's address space when it's not running

How to detect writes



- OS injects state (directly on indirectly)
 - Do it with MMU: agnostic to OS.
 - Notice modifications to sliced process's address space when it's not running

A better machine



- Slices are standalone recordings:
 - Fully capable
 - Discrete
 - Not specific to OS/VMM
 - Why not replay at user-level?

High level logs are useful



- High level log: any user-level tool can be a replayer.
- User-level a big thing:
 - Lots of user-level tools for analyzing "live" programs.
 - Crosscut supports running slices in valgrind, Omniscient debugger, Chronicle. Standard tools, gdb and visual studio. High level log: Even Windows proc on valgrind.

Beyond processes

- Processes not only abstraction
 - Sometimes not useful at all.



Beyond processes

- Processes not only abstraction
 - Sometimes not useful at all.



We can solve this with yet another better machine



• Goal: create abstraction of script execution.

- Avoid details of underlying runtime.

Recording a script generically is not possible

- Recording a script generically is not possible.
 Requires intimate knowledge of the interpreter.
- Crosscut provides an API:
 - Given to guest software to create their own inputs to their own machines.
 - Guest doesn't really need an API just to create slices
 - What's wrong with *write()*?
 - Purpose:
 - Avoid overhead in the guest for creating the slice. How?
 - Do regular VM recording
 - Defer slicing to if/when it's needed.
 - Because ex-post-facto, original recording fast
 - Can be expensive---off the critical path.

What does the API do?

- Guest has code it wants to run.
- API provides notification to VMM of:
 - work (guest code) we'd like to do
 - at a location we want to do it
 - On replay, VMM executes the guest code in sandbox.
- Notification mechanism must be fast:
 - It is running during original recording.
 - Standard signalling mechanism: instruction to trigger faults
 - Triggering faults in original recording too slow
 - Regular hypercall mechanisms out: port I/O, #UD
 - Silent/fast during recording: memory references, MMU

How to use API to slice a script

Normal script transitions



What inputs do we want to save?

– How to represent them?

Modified Perl

Make deterministic recording of a script

What inputs do we want to save?

Normal script transitions



- Simplistic example: system calls.
- Should be sufficient
 - Calls into Perl runtime will result in calls to the system.
 If not, they are deterministic.
 - Should be able to replay calls to runtime from system call inputs.

How to use API to save syscalls

On occasion of slicing a script: intercepting inputs



- Modified Perl runtime:
 - Record system call inputs (deferred in sandbox).
- Isn't this a process replay log?
- The difference: not all system calls belong to script.
 - Interpreter initiates some on its own behalf.
 - Don't save those: want interpreter's behavior capable of change

Want the interpreter to run free

On occasion of slicing a script: intercepting inputs



- When replaying, interpreter runs freely
 - It can change it's behavior, because it's live.
 - When script requests input, fed from log
 - Behaves deterministically

What can you do with Perl log?



- Created a new machine: Perl interpreter + replay mechanism
- Internal Perl facilities can be used to debug the script: Perl debugger can set breakpoints, inspect variables.
- A better machine: let Perl interpreter do debugging, more natural than stapling Perl introspection into a VMM.

Goal: record program state

• VM recordings are great.

- Provide great coverage.

• But we want to replay with discretion:

– Don't require VM.

- Don't require VMM.
- Don't sacrifice *coverage* for discretion.
- Obtain discretion with slices.
 - Abstracts replay requirements.
 - Doesn't sacrifice coverage.

fin.