Major Combinational Automatic Test-Pattern Generation Algorithms

- Definitions
- D-Algorithm (Roth) -- 1966
 - D-cubes
 - Bridging faults
 - Logic gate function change faults
- PODEM (Goel) -- 1981
 - X-Path-Check
 - Backtracing
- Summary

Forward Implication



 Results in logic gate inputs that are significantly labeled so that output is uniquely determined
 AND gate forward implication table:

×	0	1	x	D	D
0	0	0	0	0	0
1	0	1	х	D	D
x	0	х	х	х	х
D	0	D	х	D	0
D	0	D	х	0	D

Backward Implication

Unique determination of all gate inputs when the gate output and some of the inputs are given



Implication Stack

- Push-down stack. Records:
 - Each signal set in circuit by ATPG
 - Whether alternate signal value already tried
 - Portion of binary search tree already searched





Objectives and Backtracing of ATPG Algorithm

- Objective desired signal value goal for ATPG
 Guides it away from infeasible/hard solutions
- Backtrace Determines which primary input and value to set to achieve objective
 - Use testability measures



Branch-and-Bound Search

- Efficiently searches binary search tree
- Branching At each tree level, selects which input variable to set to what value
- Bounding Avoids exploring large tree portions by artificially restricting search decision choices
 - Complete exploration is impractical
 - Uses heuristics

D-Algorithm -- Roth IBM (1966)

- Fundamental concepts invented:
- First complete ATPG algorithm
- D-Cube
- D-Calculus
- Implications forward and backward
- Implication stack
- Backtrack
- Test Search Space



Minimal set of logic signal assignments to show essential prime implicants of Karnaugh map



Gate	Inp	uts	Output	Gate	In	outs	Output
AND	A	B	d	NOR	d	e	Ē
1	0	Х	0	1	1	Х	0
2	X	0	0	2	Х	1	0
3	1	1	1	3	0	0	1

D-Cube

- Collapsed truth table entry to characterize logicUse Roth's 5-valued algebra
- Can change all D's to D's and D's to D's (do both)
- AND gate:



D-Cube Operation of D-Intersection

- y undefined (same as)
- = μ or λ requires inversion of D and D
- $= D-Intersection: 0 \cap 0 = 0 \cap X = X \cap 0 = 0$ 1 \circle 1 = 1 \circle X = X \circle 1 = 1

$$X \cap X = X$$

D-containment – Cube a contains Cube b if b is a subset of a



Primitive D-Cube of Failure

- Models circuit faults:
 - Stuck-at-0
 - Stuck-at-1
 - Bridging fault (short circuit)
 - Arbitrary change in logic function
- AND Output sa0: "1 1 D"

- Wire sa0:
- Propagation D-cube models conditions under which fault effect propagates through gate

" D"

Implication Procedure

- 1. Model fault with appropriate *primitive D-cube of failure* (PDF)
- 2. Select *propagation D-cubes* to propagate fault effect to a circuit output (*D-drive* procedure)
- 3. Select *singular cover* cubes to justify internal circuit signals (*Consistency* procedure)
- Put signal assignments in *test cube*
- Regrettably, cubes are selected very arbitrarily by D-ALG

Bridging Fault Circuit



Construction of Primitive D-Cubes of Failure

- 1. Make cube set α 1 when good machine output is 1 and set α 0 when good machine output is 0
- Make cube set β1 when failing machine output is 1 and β0 when it is 0
- 3. Change α 1 outputs to 0 and D-intersect each cube with every β 0. If intersection works, change output of cube to D
- 4. Change $\alpha 0$ outputs to 1 and D-intersect each cube with every $\beta 1$. If intersection works, change output of cube to D

Bridging Fault D-Cubes of Failure

Cube-set	а	b	a*	b*	Cube-set	а	b	a* b*
αΟ	0	Х	0	Х				
	X	0	X	0				
α1	1	Х	1	X	PDFs for	1	0	<u>1</u> D
	X	1	X	1	Bridging fault	0	1	D 1
BO	0	0	0	0				
<u>β</u> 1	X	1	1	1				
	1	Х	1	1				

D-Cube of Failure								
Cube-set	a b c	Cube-set	a b	С				
αΟ	0 X 0 X 0 0	PDFs for	0 1	D				
<u>α</u> 1 80	1 1 1	AND changing	1 0					
β1	1 X 1 X 1 1			Ð				

Europhic

Change

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D-Algorithm – Top Level

- 1. Number all circuit lines in increasing level order from PIs to POs;
- 2. Select a primitive D-cube of the fault to be the test cube;
 - Put logic outputs with inputs labeled as D (D) onto the *D*-frontier;
- 3. *D-drive* ();
- 4. Consistency ();
- 5. return ();

D-Algorithm – D-drive

while (untried fault effects on D-frontier)

- select next untried D-frontier gate for propagation;
- while (untried fault effect fanouts exist)
 - select next untried fault effect fanout; generate next untried propagation D-cube;
 - D-intersect selected cube with test cube;
 - if (intersection fails or is undefined) continue:
 - if (all propagation D-cubes tried & failed) break;
 - if (intersection succeeded)
 - add propagation D-cube to test cube -- recreate D-frontier; Find all forward & backward implications of assignment; save *D*-frontier, algorithm state, test cube, fanouts, fault; break;

else if (intersection fails & D and D in test cube) Backtrack (); else if (intersection fails) break;

if (all fault effects unpropagatable) Backtrack ();

D-Algorithm -- Consistency

- g = coordinates of test cube with 1's & 0's; if (g is only PIs) f
- for (each unjustified signal in g) Select highest # unjustified signal z in g, not a PI;
- if (inputs to gate z are both D and D) break;
 - while (untried singular covers of gate z) select next untried singular cover;

 - If (no more singular covers)
 - If (no more stack choices) f
 - else if (untried alternatives in Consistency) pop implication stack -- try alternate assignment;
 - else Backtrack ();
 - D-drive ();
 - If (singular cover D-intersects with z) delete z from g, add inputs to singular cover to g, find all forward and backward implications of new assignment, and break; If (intersection fails) mark singular cover as failed;

Backtrack

if (PO exists with fault effect) Consistency (); else pop prior implication stack setting to try alternate assignment;

if (no untried choices in implication stack) fault untestable & stop;

else return;

Circuit Example 7.1 and Truth Table



Singular Cover & D-Cubes



Singular cover – Used for justifying lines

Propagation D-cubes - Conditions under which difference between good/failing machines propagates

Steps for Fault d sa0

Step	Α	В	С	d	e	F	Cube type
1	1	1		D			PDF of AND gate
2				D	0	D	Prop. D-cube for NOR
3		1	1		0		Sing. Cover of NAND



























Inconsistent

- *d* = 0 and *m* = 1 cannot justify *r* = 1 (equivalence)
 - Backtrack
 - Remove **B** = 0 assignment











- Use X-PATH-CHECK to test whether D-frontier still there
- Objectives -- bring ATPG closer to propagating D (D) to PO
- Backtracing

Motivation

- IBM introduced semiconductor DRAM memory into its mainframes – late 1970's
- Memory had error correction and translation circuits improved reliability
 - D-ALG unable to test these circuits
 - Search too undirected
 - Large XOR-gate trees
 - Must set all external inputs to define output
 - Needed a better ATPG tool

PODEM High-Level Flow

- 1. Assign binary value to unassigned PI
- 2. Determine implications of all PIs
- 3. Test Generated? If so, done.
- 4. Test possible with more assigned PIs? If maybe, go to Step 1
- 5. Is there untried combination of values on assigned PIs? If not, exit: untestable fault
- 6. Set untried combination of values on assigned PIs using objectives and backtrace. Then, go to Step 2





































Backtrace (s, v_s) **Pseudo-Code**

v = v_s; while (s is a gate output)

if (s is NAND or INVERTER or NOR) $v = \overline{v}$; if (objective requires setting all inputs) select unassigned input *a* of *s* with hardest controllability to value *v*;

else

select unassigned input a of s with easiest controllability to value v;

S = a

return (s, v) /* Gate and value to be assigned */;

Objective Selection Code

if (gate g is unassigned) return (g, \overline{v}); select a gate *P* from the D-frontier; select an unassigned input / of P; if (gate *g* has controlling value) c = controlling input value of q; else if (0 value easier to get at input of XOR/EQUIV gate) *c* = 1; else c = 0; return (I, c);

PODEM Algorithm

while (no fault effect at POs) if (xpathcheck (D-frontier) (I, v) = Objective (fault, v_{fault}); $(p_i, v_{p_i}) = Backtrace (i, v_i);$ $imply (p_i, v_{p_i});$ if (PODEM (failt, $v_{fail}) == SUCCESS)$ return (SUCCESS); $(p_i, v_{p_i}) = Backtrack 0;$ $\begin{array}{l} (p, v_{p}) = Backtrack (; \\ (pl, v_{p}) = Backtrack (; \\ (pl, v_{p}); \\ (pl, v_{p}); \\ (poper (rault, v_{fault}) == SUCCESS) return \\ (SUCCESS); \end{array}$ Imply (pi, "X"); return (FAILURE); else if (implication stack exhausted) return (FAILURE); else *Backtrack* (); return (SUCCESS);

Summary

- D-ALG First complete ATPG algorithm
 - D-Cube
 - D-Calculus
 - Implications forward and backward
 - Implication stack
- Backup PODEM

 - Expand decision tree only around PIs
 - Use *X-PATH-CHECK* to see if *D-frontier* exists Objectives -- bring ATPG closer to getting D (D) to PO
 - Backtracing